# Rational function semifields of tropical curves are finitely generated over the tropical semifield 

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#### Abstract

概 要 We prove that the rational function semifield of a tropical curve is finitely generated as a semifield over the tropical semifield $\boldsymbol{T}:=(\boldsymbol{R} \cup\{-\infty\}$ ，max,+$)$ by giving a specific finite generating set．


## 1 Introduction

This manuscript is on tropical geometry．Tropical geometry is an algebraic geometry over the tropical semifield $\boldsymbol{T}:=(\boldsymbol{R} \cup\{-\infty\}$ ，max，+ ），which is developing from the bigining of this century．In tropical geometry，we study tropical varieties．Tropical varieties are polyhedral complexes obtained from algebraic varieties by a limit operation called tropicalization．Tropical curves，one－dimensional tropical varieties，correspond to algebraic curves．In this manuscript， an（abstract）tropical curve means a metric graph，i．e．，the underlying metric space of the pair （ $G, l$ ），where $G$ is an unweighted，undirected，finite，connected，nonempty multigraph that may have loops（in what follows，we call such a multigraph a graph simply），$l: E(G) \rightarrow \boldsymbol{R}$ is a length function，and $E(G)$ stands for the set of edges of $G$ ．For a tropical curve $\Gamma$ ，a map $f: \Gamma \rightarrow \boldsymbol{R} \cup\{-\infty\}$ is a rational function on $\Gamma$ if $f$ is a continuos piecewise affine function with integer slopes，or a constant function of $-\infty$ ．The set $\operatorname{Rat}(\Gamma)$ of all rational functions on $\Gamma$ becomes a semifield over $\boldsymbol{T}$ with the pointwise maximum operation as its addition $\oplus$ and the pointwise usual addition as its multiplication $\odot$ ．We call $\operatorname{Rat}(\Gamma)$ the rational function semifield of $\Gamma$ ．The following theorem is our main theorem：

Theorem 1．Let $\Gamma$ be a tropical curve．Then，the rational function semifield $\operatorname{Rat}(\Gamma)$ is finitely generated as a semifield over $\boldsymbol{T}$ ．

This is a tropical analogue of the fact that the function field of an algebraic curve over $\boldsymbol{C}$ is generated by two elements over $\boldsymbol{C}$ ．To prove this theorem，the following lemma is our key：

Lemma 2 （［1，Lemma 2］）．Let $\Gamma$ be a tropical curve．Then， $\operatorname{Rat}(\Gamma)$ is generated by all chip firing moves and all constant functions as a group with $\odot$ as its binary operation．

Here，a chip firing move $\mathrm{CF}\left(\Gamma_{1}, l\right)$ is the rational function defined by the pair of a subgraph $\Gamma_{1}$ and a number $l \in \boldsymbol{R}_{>0}$ as follows： $\mathrm{CF}\left(\Gamma_{1}, l\right)(x):=-\min \left\{\operatorname{dist}\left(\Gamma_{1}, x\right), l\right\}$ ，where a subgraph is a closed subset with a finite number of connected components and dist $\left(\Gamma_{1}, x\right)$ denotes the distance
between $\Gamma_{1}$ and $x$. By this lemma, it is enough to find a finite set of rational functions which generates all chip firing moves as a semifield over $\boldsymbol{T}$.

Note that this manuscript is based on [2]. In [2], a tropical curve may have edges of length $\infty$. Even a tropical curve has an edge of length $\infty$, the assertion of Theorem 1 holds.

## 2 Proof of Theorem 1

Let $\Gamma$ be the tropical curve obtained from the pair $(G, l)$ of a graph $G$ and a length function $l$. Then, the pair $(G, l)$ is called a model for $\Gamma$. There are many possible models for $\Gamma$. We frequently identify a vertex (resp. an edge) of $G$ with the corresponding point (resp. the corresponding closed subset) of $\Gamma$. The canonical model ( $G_{\circ}, l_{\circ}$ ) for $\Gamma$ is defined as follows: if $\Gamma$ is not homeomorphic to a circle, then let $V\left(G_{\circ}\right):=\{x \in \Gamma \mid \operatorname{val}(x) \neq 2\}$; otherwise, let $V\left(G_{\circ}\right)$ be consist of arbitrary one point of $\Gamma$, where $\operatorname{val}(x)$ denotes the minimum number of the connected components of $U \backslash\{x\}$ with any neighborhood $U$ of $x$. Let $G_{\circ}$ be the graph which has $V\left(G_{\circ}\right)$ as its set of vertices and the set of all closures of connected components of $\Gamma \backslash V\left(G_{\circ}\right)$ as its set of edges, and $l_{\circ}$ the length function which maps each edge of $G_{\circ}$ to its length. Fix a direction on edges of $G_{\circ}$. Let each $e \in E\left(G_{\circ}\right)$ be identified with the interval $\left[0, l_{\circ}(e)\right]$ with this direction. For each edge $e \in E\left(G_{\circ}\right)$, let $x_{e}=\frac{l_{0}(e)}{4}$, $y_{e}=\frac{l_{\circ}(e)}{2}$, and $z_{e}=\frac{3 l_{\circ}(e)}{4}$. We define rational functions

$$
f_{e}:=\operatorname{CF}\left(\left\{y_{e}\right\}, \frac{l_{\circ}(e)}{2}\right), g_{e}:=\operatorname{CF}\left(\left\{x_{e}\right\}, \frac{l_{\circ}(e)}{4}\right), h_{e}:=\operatorname{CF}\left(\left\{z_{e}\right\}, \frac{l_{\circ}(e)}{4}\right) .
$$

Let $d$ be the diameter of $\Gamma$, i.e., $d=\sup \{\operatorname{dist}(x, y) \mid x, y \in \Gamma\}=\max \{\operatorname{dist}(x, y) \mid x, y \in \Gamma\}$. Let $R$ be the semifield generated by $f_{e}, g_{e}, h_{e}$ for any $e \in E\left(G_{\circ}\right)$ and $\operatorname{CF}(\{v\}, d)(=-\operatorname{dist}(v, \cdot))$ for any $v \in V\left(G_{\circ}\right)$ over $\boldsymbol{T}$. We will show that this semifield $R$ coincides with $\operatorname{Rat}(\Gamma)$.

In Algorithm $1, \overline{e_{i} \backslash S}$ denotes the closure of $e_{i} \backslash S$, and if $S$ consists of only one point $x$, then we write $l_{x}$ instead of $l_{\{x\}}$.
Remark 3. Let $\Gamma$ be a tropical curve and $S_{1}$ a proper connceted subgraph of $\Gamma$. Let $l \leq l_{S_{1}}$ and $S_{2}:=\left\{x \in \Gamma \mid \operatorname{dist}\left(S_{1}, x\right) \leq l\right\}$. With $a:=\min \left\{k \in Z_{>0} \mid l / k \leq l_{S_{2}}\right\}, m:=\min \{k \in$ $\left.\boldsymbol{Z}_{>0} \mid l_{S_{2}} / k \leq l / a\right\}$ and any $l^{\prime}>0$, by the definition of chip firing moves, we have

$$
\begin{gathered}
\operatorname{CF}\left(S_{2}, \frac{l}{a}\right)=\operatorname{CF}\left(S_{1}, \frac{l}{a}\right) \odot \bigodot_{k=1}^{a} \bigodot_{\substack{x^{\prime} \in \Gamma \cdot \\
\operatorname{dist}\left(S_{1}, x^{\prime}\right)=\frac{k l}{a}}}\left\{\operatorname{CF}\left(\left\{x^{\prime}\right\}, \frac{l}{a}\right) \odot \frac{l}{a}\right\}, \\
\mathrm{CF}\left(\left\{x \in \Gamma \left\lvert\, \operatorname{dist}\left(S_{2}, x\right) \leq \frac{k l_{S_{2}}}{m}\right.\right\}, \frac{l_{S_{2}}}{m}\right)= \\
\operatorname{CF}\left(\left\{x \in \Gamma \left\lvert\, \operatorname{dist}\left(S_{2}, x\right) \leq \frac{(k-1) l_{S_{2}}}{m}\right.\right\}, \frac{l_{S_{2}}}{m}\right) \\
\odot \bigodot_{\substack{x^{\prime} \in \Gamma: \\
\operatorname{dist}\left(S_{2}, x^{\prime}\right)=\frac{k l S_{2}}{m}}}\left\{\operatorname{CF}\left(\left\{x^{\prime}\right\}, \frac{l_{S_{2}}}{m}\right) \odot \frac{l_{S_{2}}}{m}\right\}, \\
\mathrm{CF}\left(S_{2}, l_{S_{2}}\right)=\left\{\operatorname{CF}\left(S_{2}, \frac{l}{a}\right) \oplus\left(-\frac{l_{S_{2}}}{m}\right)\right\} \odot \bigodot_{k=1}^{m-1} \mathrm{CF}\left(\left\{x \in \Gamma \left\lvert\, \operatorname{dist}\left(S_{2}, x\right) \leq \frac{k l_{S_{2}}}{m}\right.\right\}, \frac{l_{S_{2}}}{m}\right),
\end{gathered}
$$

and

$$
\mathrm{CF}\left(S_{1}, l+l^{\prime}\right)=\mathrm{CF}\left(S_{1}, l\right) \odot \mathrm{CF}\left(S_{2}, l^{\prime}\right) .
$$

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Algorithm 1
Input: \(\Gamma\) : a tropical curve
    \(E\left(G_{\circ}\right)=\left\{e_{1}, \ldots, e_{n}\right\}\) : a labeling of edges of the canonical model for \(\Gamma\)
    \(S\) : a proper connected subgraph of \(\Gamma\)
Output: \(l_{S}\)
    \(i \leftarrow 1\)
    while \(i \leq n\) do
        if \(e_{i} \cap S=\varnothing\) then
            \(l_{i} \leftarrow(\) the diameter of \(\Gamma), i \leftarrow i+1\)
        else \(\left\{S \supset e_{i}\right\}\)
            \(l_{i} \leftarrow\) (the diameter of \(\left.\Gamma\right), i \leftarrow i+1\)
        else \(\left\{S \supset \partial e_{i}\right\}\)
            \(l_{i} \leftarrow\left(\right.\) the length of \(\left.\overline{e_{i} \backslash S}\right) / 2, i \leftarrow i+1\)
        else \(\left\{S \subset e_{i}^{\circ}\right\}\)
            \(l_{i} \leftarrow \min \left\{\operatorname{dist}(S, x) \mid x\right.\) is one of the endpoints of \(\left.e_{i}\right\}, i \leftarrow i+1\)
        else
            \(l_{i} \leftarrow\left(\right.\) the length of \(\left.\overline{e_{i} \backslash S}\right), i \leftarrow i+1\)
        end if
    end while
    \(l_{S} \leftarrow \min \left\{l_{1}, \ldots, l_{n}\right\}\)
    return \(l_{S}\)
```

Lemma 4. Let $e$ be an edge of $G_{\circ}$ and $v, w$ the endpoint(s) of $e$ (possibly $v=w$ ). Let $x$ be in $e \backslash\{v, w\}$. Then, $\operatorname{CF}\left(\{x\}, l_{x}\right) \in R$.

Proof. If $x$ is the midpoint of $e$, then $\operatorname{CF}\left(\{x\}, l_{x}\right)=f_{e} \in R$. Suppose that $x$ is not the midpoint of $e$. Assume that $0<l_{x} \leq \frac{l_{0}(e)}{4}$ and $g_{e}(x)=-\frac{l_{0}(x)}{4}$. Then

$$
\begin{aligned}
\mathrm{CF}\left(\{x\}, l_{x}\right)= & \left\{\left(\frac{l_{\circ}(e)}{2}-l_{x}\right) \odot f_{e} \oplus\left(l_{x}-\frac{l_{\circ}(e)}{2}\right) \odot f_{e}^{\odot(-1)}\right\}^{\odot(-1)} \\
& \odot\left(-\frac{l_{\circ}(e)}{4}\right) \odot g_{e}^{\odot(-1)} \oplus\left(-l_{x}\right) \in R .
\end{aligned}
$$

Similarly, if $0<l_{x} \leq \frac{l_{0}(e)}{4}$ and $h_{e}(x)=-\frac{l_{o}(e)}{4}$, then $\mathrm{CF}\left(\{x\}, l_{x}\right) \in R$.
When $\frac{l_{0}(e)}{4}<l_{x} \leq \frac{l_{o}(e)}{3}$ and $g_{e}(x)=-\frac{l_{o}(e)}{4}$, we have

$$
\begin{aligned}
\mathrm{CF}\left(\{x\}, l_{x}\right)= & \left\{\left(\frac{l_{\circ}(e)}{2}-l_{x}\right) \odot f_{e} \oplus\left(l_{x}-\frac{l_{\circ}(e)}{2}\right) \odot f_{e}^{\odot(-1)}\right\}^{\odot(-1)} \\
& \odot\left\{\left(-\frac{l_{\circ}(e)}{4}\right) \odot g_{e}^{\odot(-1)}\right\}^{\odot 2} \oplus\left(-l_{x}\right) \in R .
\end{aligned}
$$

Similarly, if $\frac{l_{o}(e)}{4}<l_{x} \leq \frac{l_{o}(e)}{3}$ and $h_{e}(x)=-\frac{l_{o}(e)}{4}$, then $\mathrm{CF}\left(\{x\}, l_{x}\right) \in R$.

When $\frac{l_{\circ}(e)}{3}<l_{x}<\frac{l_{\circ}(e)}{2}$ and $g_{e}(x)=-\frac{l_{\circ}(e)}{4}$, we have

$$
\begin{aligned}
\mathrm{CF}\left(\{x\}, l_{\circ}(e)-2 l_{x}\right)= & \left\{\left(\frac{l_{\circ}(e)}{2}-l_{x}\right) \odot f_{e} \oplus\left(l_{x}-\frac{l_{\circ}(e)}{2}\right) \odot f_{e}^{\odot(-1)}\right\}^{\odot(-1)} \\
& \odot\left\{\left(-\frac{l_{\circ}(e)}{4}\right) \odot g_{e}^{\odot(-1)}\right\}^{\odot 2} \oplus\left(2 l_{x}-l_{\circ}(e)\right) \in R .
\end{aligned}
$$

Similarly, if $\frac{l_{\circ}(e)}{3}<l_{x}<\frac{l_{\circ}(e)}{2}$ and $h_{e}(x)=-\frac{l_{\circ}(e)}{4}$, then $\mathrm{CF}\left(\{x\}, l_{\circ}(e)-2 l_{x}\right) \in R$.
Let $x$ be in the fifth case. Since

$$
\begin{aligned}
& \mathrm{CF}\left(\left\{x_{1} \in \Gamma \left\lvert\, \operatorname{dist}\left(x, x_{1}\right) \leq \frac{l_{\circ}(e)}{2}-l_{x}\right.\right\}, \frac{l_{\circ}(e)}{2}-l_{x}\right) \\
= & \operatorname{CF}\left(\{x\}, \frac{l_{\circ}(e)}{2}-l_{x}\right) \odot \bigodot_{\substack{x_{1} \in e: \\
\operatorname{dist}\left(x, x_{1}\right)=\frac{l_{\circ}(e)}{2}-l_{x}}}\left\{\operatorname{CF}\left(\left\{x_{1}\right\}, \frac{l_{\circ}(e)}{2}-l_{x}\right) \odot\left(\frac{l_{\circ}(e)}{2}-l_{x}\right)\right\} \in R,
\end{aligned}
$$

with inputs $l=\frac{l_{\circ}(e)}{2}-l_{x}, S_{1}=\left\{x_{1} \in \Gamma \mid \operatorname{dist}\left(x, x_{1}\right) \leq l_{\circ}(e) / 2-l_{x}\right\}$ in Remark 3, we have $\mathrm{CF}\left(\{x\}, l_{x}\right) \in R$.

When $x$ is in the sixth case, by the same argument, we have $\mathrm{CF}\left(\{x\}, l_{x}\right) \in R$.
Note that $l_{x}$ coincides with $\min (\operatorname{dist}(x, v)$, $\operatorname{dist}(x, w))$ in the setting of Lemma 4.
By Remark 3 and Lemma 4, we prove the following three lemmas.
Lemma 5. For any $x \in \Gamma$ and any positive real number $l, \operatorname{CF}(\{x\}, l)$ is in $R$.
Proof. For any $x \in \Gamma$ and $l>0$, by the definition of chip firing moves, we have $\operatorname{CF}(\{x\}, l)=$ $\mathrm{CF}(\{x\}, d) \oplus(-l)$. Hence it is sufficient to check that $\mathrm{CF}(\{x\}, d) \in R$. If $x \in V\left(G_{\circ}\right)$, then $\mathrm{CF}(\{x\}, d) \in R$.

Suppose that there exists an edge $e \in E\left(G_{\circ}\right)$ which does not have $x$ as an endpoint. Considering Remark 3 with $l=l_{x}, S_{1}=\{x\}$, by Lemma 4, we have

$$
\mathrm{CF}\left(S_{2}, l_{S_{2}}\right) \in R
$$

and

$$
\mathrm{CF}\left(S_{1}, l+l_{S_{2}}\right)=\mathrm{CF}\left(S_{1}, l\right) \odot \mathrm{CF}\left(S_{2}, l_{S_{2}}\right) \in R
$$

Since $S_{2}$ contains a lot of whole edges of $G_{\circ}$ more than $S_{1}$ and the set of edges of $G_{\circ}$ is finite, by repeating inputs of $l=l_{S_{2}}, S_{1}=S_{2}$ in Remark 3, we have $\operatorname{CF}(\{x\}, d) \in R$.

Lemma 6. For any proper connected subgraph $\Gamma_{1}$ and any positive real number $l, \operatorname{CF}\left(\Gamma_{1}, l\right)$ is in $R$.

Proof. By Lemma 5, if $\Gamma_{1}$ consists of only one point, then we have the conclusion. Assume that $\Gamma_{1}$ does not consist of only one point.

Suppose that $\Gamma_{1}$ contains no whole edges of $G_{\circ}$ and that there exists an edge $e \in E\left(G_{\circ}\right)$ containing $\Gamma_{1}$. Let $x_{1}$ and $x_{2}$ be the endpoints of $\Gamma_{1}$. Let $x$ be the midpoint of $\Gamma_{1}$. By Lemma 5 , for any positive real number $l$, we have

$$
\mathrm{CF}\left(\Gamma_{1}, l\right)=\left[\left\{\mathrm{CF}\left(\{x\}, l+\operatorname{dist}\left(x_{1}, x\right)\right) \odot \operatorname{dist}\left(x_{1}, x\right)\right\}^{\odot(-1)} \oplus 0\right]^{\odot(-1)} \in R
$$

Suppose $\Gamma_{1}$ contains $p$ edges．Let $\partial \Gamma_{1} \cup\left(V\left(G_{\circ}\right) \cap \Gamma_{1}\right)=\left\{x_{1}, \ldots, x_{q}\right\}$ ．We may assume that $x_{1}, \ldots, x_{q}$ are distinct．Let $\Gamma_{11}, \ldots, \Gamma_{1 s}$ be connected components of $\Gamma_{1} \backslash\left\{x_{1}, \ldots, x_{q}\right\}$ ．For a suffi－ ciently small positive real number $\varepsilon$ ，let $\Gamma_{1 i}^{\prime}$ be the connected subgraph $\left\{x \in \Gamma_{1 i} \mid\right.$ for any $j, \operatorname{dist}\left(x, x_{j}\right) \geq$ $\varepsilon\}$ of $\Gamma$ ．Then，we have

$$
\mathrm{CF}\left(\Gamma_{1}, \varepsilon\right)=\left\{\bigoplus_{k=1}^{q} \mathrm{CF}\left(\left\{x_{k}\right\}, \varepsilon\right)\right\} \odot \bigodot_{k=1}^{s}\left(\varepsilon \odot \mathrm{CF}\left(\Gamma_{1 k}^{\prime}, \varepsilon\right)\right) .
$$

The last divisor is in the first case，and thus it is in $R$ ．By inputting $l=\varepsilon, S_{1}=\Gamma_{1}$ and by repeating inputs $l=l_{S_{2}}, S_{1}=S_{2}$ in Remark 3，we have $\operatorname{CF}\left(\Gamma_{1}, d\right) \in R$ ．From this，for any $l>0$ ， we have $\mathrm{CF}\left(\Gamma_{1}, l\right)=\mathrm{CF}\left(\Gamma_{1}, d\right) \oplus(-l) \in R$ ．

Lemma 7．For any proper subgraph $\Gamma_{1}$ and any positive real number $l, \mathrm{CF}\left(\Gamma_{1}, l\right)$ is in $R$ ．
Proof．Let $\Gamma_{1}$ be a proper subgraph of $\Gamma$ ．Let $s$ be the number of connected components of $\Gamma_{1}$ ．If $s=1$ ，then the conclusion follows Lemma 6．Assume $s \geq 2$ ．Let $\Gamma_{1}^{\prime}, \ldots, \Gamma_{s}^{\prime}$ be all the distinct connected components of $\Gamma_{1}$ ．For $l^{\prime}>0$ ，let $\Gamma_{k}^{\prime}\left(l^{\prime}\right):=\left\{x \in \Gamma \mid \operatorname{dist}\left(\Gamma_{k}^{\prime}, x\right) \leq l^{\prime}\right\}$ ．If $l^{\prime}$ is sufficiently small，then the intersection of $\Gamma_{1}^{\prime}\left(l^{\prime}\right), \ldots, \Gamma_{s}^{\prime}\left(l^{\prime}\right)$ is empty．Let $l_{1}^{\prime}$ be the minimum value of $l^{\prime}$ such that this intersection is nonempty．By induction on $s, \mathrm{CF}\left(\bigcup_{k=1}^{s} \Gamma_{k}^{\prime}\left(l_{1}^{\prime}\right), d\right) \in R$ ． On the other hand，

$$
\mathrm{CF}\left(\Gamma_{1}, l_{1}^{\prime}\right)=\bigoplus_{k=1}^{s} \mathrm{CF}\left(\Gamma_{k}^{\prime}, l_{1}^{\prime}\right) \in R .
$$

Hence

$$
\mathrm{CF}\left(\Gamma_{1}, d\right)=\mathrm{CF}\left(\Gamma_{1}, l_{1}^{\prime}\right) \odot \mathrm{CF}\left(\bigcup_{k=1}^{s} \Gamma_{k}^{\prime}\left(l_{1}^{\prime}\right), d\right) \in R .
$$

In conclusion，for any $l>0$ ，we have

$$
\mathrm{CF}\left(\Gamma_{1}, l\right)=\mathrm{CF}\left(\Gamma_{1}, d\right) \oplus(-l) \in R .
$$

From Lemmas 2，5，6，7，we have Theorem 1.

## 参考文献

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