

Doubly autoparallel structure and curvature integrals

- An application to iteration-complexity analysis of convex optimization -

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統計多様体の幾何学とその周辺(14)

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Outline

- Introduction: Doubly autoparallel submanifolds
- Preliminaries
 - Dually flat structure on a symmetric cones
- **Characterization of DA submfd in sym. cones**
 - Several applications
- Conic linear program on convex cones Ω
 - Central trajectory
 - Geometric predictor-corrector method in Ω^*
- **Curvature integral and iteration-complexity**
- Application to primal-dual path following methods
- Concluding remark

Introduction

Doubly autoparallel submanifolds

Def. Statistical manifold: $(\mathcal{S}, g, \nabla, \nabla^*)$

$$Xg(Y, Z) = g(\nabla_X Y, Z) + g(Y, \nabla_X^* Z)$$

X, Y and Z : arbitrary vector fields on \mathcal{S}

★ g : Riemannian metric

★ (∇, ∇^*) : torsion-free affine connections

$$R^\nabla = 0, R^{\nabla^*} = 0 \Rightarrow \text{dually flat}$$

★ $\nabla^{(\alpha)} = \frac{1+\alpha}{2}\nabla + \frac{1-\alpha}{2}\nabla^*$: α -connections

- Def. Let (S, g, ∇, ∇^*) be a statistical manifold and M be its submanifold. We call M a **doubly autoparallel** submanifold in S when the followings hold:

- $\forall X, Y \in \mathcal{X}(M), \nabla_X Y \in \mathcal{X}(M)$
i.e. $H_M(X, Y) = 0$

- $\forall X, Y \in \mathcal{X}(M), \nabla_X^* Y \in \mathcal{X}(M)$
i.e. $H_M^*(X, Y) = 0$

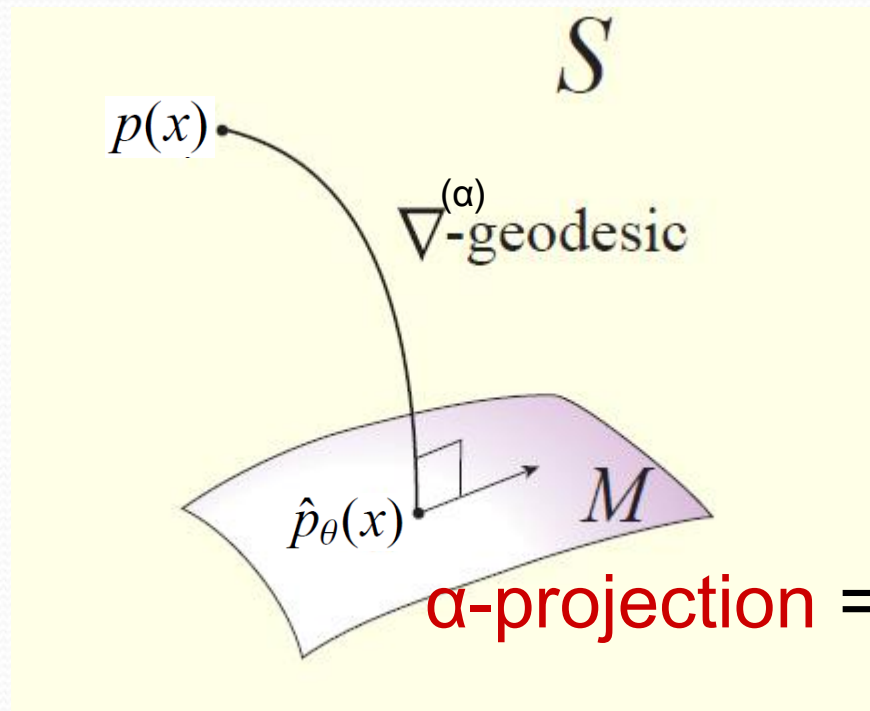
Important Properties

Proposition The following statements are equivalent:

- 1) A submanifold M is doubly autoparallel (DA)
- 2) M is autoparallel w.r.t. the α -connections
$$\nabla^{(\alpha)} = \{(1 + \alpha)\nabla + (1 - \alpha)\nabla^*\}/2$$
for **two different** α 's.
- 3) M is autoparallel w.r.t. **all** the α -connections.
- 4) **all** the α -geodesics connecting two points on M lay in M (if it is simply connected).
- 5) M is affinely constrained in both ∇ - and ∇^* -affine coordinates **if S is dually flat**.

Furthermore, for a parametric statistical model S

- If M is DA in S , then α -projections (q -MaxEnt) from p to M are unique for **all** α if they exist.



M is simultaneously an
exp. and mix. family

$$\alpha\text{-projection} = \arg \min D^{(\alpha)}(p, M)$$

Related topics and applications

Symmetric cones

- MLE for structured covariance matrices is tractable
(cast to convex program: **inversely linear structure**)
[Anderson 70, Malley 94]
- **Explicitly solvable Semi-Definite Programs** [O 99]
- Structure of α -power means on symmetric cones [O 04]

Related topics and applications

Probability simplex

- Statistical models Markov-isomorphic to the probability simplex [Nagaoka 17]
- Characterization and classification of DA submfds in prob. simplex via Hadamard algebra [O&Ishi 18]
- Learning theory [Mutus&Ay 03]

Miscellaneous

- The self-similar (*Barenblatt–Pattle*) solution for the porous medium equation [O&Wada 10]

General statistical manifolds

- Purely geometric study [Satoh *et al.* 21]

Preliminaries

[Faraut&Korani 94]

- Symmetric cone Ω in an Euclidian space E

- Homogeneous

$G(\Omega) = \{\tau \in GL(E) \mid \tau(\Omega) = \Omega\}$ acts transitively

- self-dual w.r.t. an inner product of E

$$\Omega = \Omega^*, \quad \Omega^* = \{y \in E \mid (x|y) > 0, \forall x \in \overline{\Omega} \setminus \{0\}\}$$

- Euclidean Jordan algebra $(V, *)$

- Commutative

- $x^2 * (x * y) = x * (x^2 * y)$, where $x^2 = x * x$

- Associative inner-product $(x * y|z) = (y|x * z)$

Prop. $\Omega = \text{int}\{x^2 \mid x \in V\}$ is a symmetric cone in V .

Ex. the set of real symmetric p. d. matrices $PD(n, \mathbf{R})$

$$V = \text{Sym}(n; \mathbf{R}), \quad X * Y = (XY + YX)/2$$

$$\tau_G(X) = GXG^T, \quad G \in GL(n, \mathbf{R})$$

$$(X|Y) = \text{tr}(XY), \quad \text{the unit: } I, \text{ the inverse: } X^{-1}$$

- $L(x) : V \rightarrow V, \quad L(x)y = x * y$
- $P(x, y) := L(x)L(y) + L(y)L(x) - L(x * y)$
- **Mutation:** $x \perp_a y := P(x, y)a$
isomorphic to $*$, the unit element: a^{-1}

Ex. $X \perp_A Y = (XAY + YAX)/2$

Preliminaries (Dually flat structure on Ω)

- Logarithmic characteristic function on Ω

$$\psi(x) := \log \int_{\Omega^*} e^{-\langle s, x \rangle} ds,$$

- positive definite Hessian on Ω
- $x^{-1} = -\text{grad } \psi(x), \quad (\text{grad } f(x)|u) = D_u f(x)$
- Ex. $\psi(x) = -\sum_{i=1}^n \log x^i$ on $\mathbf{R}_{++}^n, \quad \psi(x) = -\log \det X$ on $\text{PD}(n)$

- a coordinate system $(x^i) \quad x = \sum_{i=1}^n x^i e_i, \quad \{e_i\}_{i=1}^n : \text{a basis of } E$
- a dual coordinate system (s_i)

$$x^{-1} = \sum_{i=1}^n s_i e^i, \quad \{e^i\}_{i=1}^n : \text{a basis of } E \text{ with } (e^i | e_j) = -\delta_j^i$$

- D : the canonical **flat** affine connection on E
- $\{x^1, \dots, x^n\}$: affine coordinate system, i.e., $D \frac{\partial}{\partial x^i} \frac{\partial}{\partial x^j} = 0$
- g : Riemannian metric on Ω

$$g = Dd\psi = \sum_{i,j} \frac{\partial^2 \psi}{\partial x^i \partial x^j} dx^i dx^j.$$

- D' : the dual affine connection on Ω

$$Xg(Y, Z) = g(D_X Y, Z) + g(Y, D'_X Z)$$

(g, D, D') : dually flat structure on Ω

Pleriminaries and ex. on PD(n)

Dually flat structure on Ω [Uohashi&O 04]

- Potential: $-\log \det x$,

Ex. $-\log \det X$, ($X = \sum_{i=1}^N x^i E_i$.)

- Riemannian metric: $g_x \left(\frac{\partial}{\partial x^i}, \frac{\partial}{\partial x^j} \right) = (P(x)^{-1} e_i | e_j)$, $P(x) := P(x, x)$

$$g_X \left(\frac{\partial}{\partial x^i}, \frac{\partial}{\partial x^j} \right) = \text{tr}(X^{-1} E_i X^{-1} E_j)$$

- α -connections: $\left(\nabla_{\frac{\partial}{\partial x^i}}^{(\alpha)} \frac{\partial}{\partial x^j} \right)_x = (\alpha - 1)(e_i \perp_{x^{-1}} e_j)$

$$\left(\nabla_{\frac{\partial}{\partial x^i}}^{(\alpha)} \frac{\partial}{\partial x^j} \right)_X = \frac{\alpha - 1}{2} (E_i X^{-1} E_j + E_j X^{-1} E_i)$$

Characterization of DA submfds in Ω

Let W be a linear subspace in Jordan algebra $(V, *)$ and $p = q * q$ in Ω .

Thm. [OIT] The following 1)-3) are equivalent:

1) A Submanifold $M = (W + p) \cap \Omega$ is DA, where

$$W + p = \{w + p \mid w \in W\}$$

2) For **all** x in M , $u \perp_{x^{-1}} v \in W$, ($u, v \in W$)

3) The subspace $P(q)^{-1}W$ is a Jordan subalgebra.

Rem. (a) 3) is able to be checked at the **single** point p

(b) $M = \{(W' + p^{-1}) \cap \Omega\}^{-1}$ with $W' = P(p)^{-1}W$

The proof is based on 5) in the Proposition

(c) **Implication:** Classification of DA submflds in Ω reduces to that of Jordan subalgs of $(V, *)$.
 (For $V = \text{Sym}(n, \mathbf{R}) \rightarrow$ [Jacobson 87], [Malley 87])

- Ex. - Jordan subalgebras in $\text{Sym}(n, \mathbf{R})$
 - 1) fixed eigen vectors, 2) doubly symmetric, etc.

- Two bases $\{E_i\}_{i=1}^m$ and $\{F^i\}_{i=1}^m$ of $\text{Sym}(n, \mathbf{R})$

$$\mathcal{M} = \left\{ P \mid P = E_0 + \sum_{i=1}^m x^i E_i, \exists x = (x^i) \in \mathbf{R}^m \right\} \cap \text{PD}(n)$$

$$\mathcal{M} = \left\{ P \mid P^{-1} = F^0 + \sum_{i=1}^m s_i F^i, \exists s = (s_i) \in \mathbf{R}^m \right\} \cap \text{PD}(n).$$

Application(1) Means on Positive Operators

[Kubo & Ando 80]

- Def. (Axioms of means)

σ is a **mean** on self-conjugate positive operators

- i) $A \leq C, B \leq D \Rightarrow A\sigma B \leq C\sigma D$
- ii) $C(A\sigma B)C = (CAC)\sigma(CBC)$
- iii) $A_n \downarrow A, B_n \downarrow B \Rightarrow A_n\sigma B_n \downarrow A\sigma B$

where $A_n \downarrow A \stackrel{\text{def}}{\Leftrightarrow} (A_i \geq A_{i+1}, \forall i) \& (A_n \rightarrow A)$

- iv) $I\sigma I = I$

α -geodesics on PD(n)

- α -geodesic $P(s)$ boundary conds. : $P(0)=A, P(1)=B$

$$P^{(\alpha)}(s) = A^{1/2} \left\{ [(A^{-1/2}BA^{-1/2})^\alpha - I]s + I \right\}^{1/\alpha} A^{1/2}$$

$$\alpha = 1 \quad P(s) = A + s(B - A)$$

$$\alpha = 0 \quad \hat{P}(s) = A^{1/2} \exp(s \log A^{-1/2}BA^{-1/2}) A^{1/2}$$

$$\alpha = -1 \quad P^*(s) = \{A^{-1} + s(B^{-1} - A^{-1})\}^{-1}$$

$$AaB := P(1/2)$$

$$AgB := \hat{P}(1/2)$$

$$AhB := P^*(1/2)$$

$P^{(\alpha)}(1/2)$: a power mean

Means and α -geodesics on $PD(n)$ [0 04]

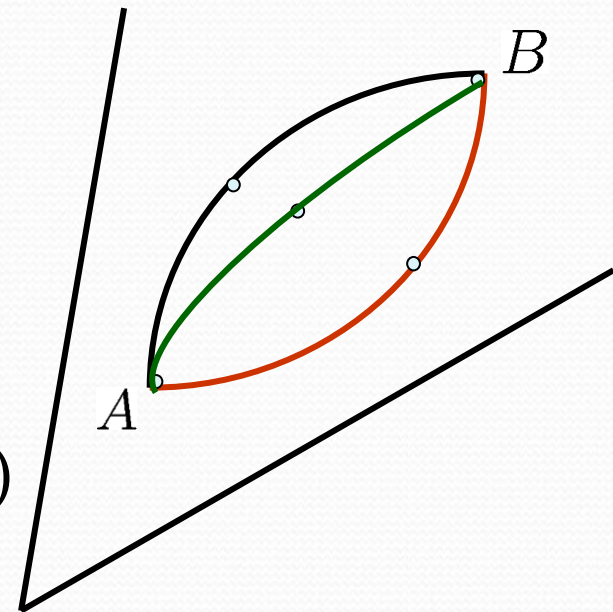
Thm. Points on α -geodesics for s in $[0,1]$ and α in $[-1,1]$ are 2-param. family of means, i.e.,

$$A\sigma_s^{(\alpha)}B = P^{(\alpha)}(s)$$

In particular, for fixed s in $[0, 1]$

$$P^{(\alpha)}(s) > P^{(\beta)}(s),$$

$$1 \geq \alpha > \beta \geq -1 \quad \text{AGH ineq. (s=1/2)}$$



Cor. A and B are in a DA submanifold M

$$\Rightarrow A\sigma_s^{(\alpha)}B \in M, s \in [0, 1], \alpha \in [-1, 1]$$

App.(2) MLE for structured covariance matrices

- Sample covariance S in $\text{PD}(n, \mathbf{R})$
- a zero-mean Gaussian p.d.f. with covariance mtx. Σ

$$p(x) = (2\pi)^{-n/2} (\det \Sigma)^{-1/2} \exp\left\{-\frac{1}{2}x^T \Sigma^{-1}x\right\}$$

- structured covariance mtx. (with linear constraints)

$$\Sigma \in \mathcal{M} = (E_0 + \mathcal{W}) \cap \text{PD}(n, \mathbf{R})$$

- Ex.

- Toeplitz matrices: $\{T = (t_{ij}) | t_{ij} = t_{ji} = y_{|i-j|}\}$
- zero-patterns : $\{\Sigma = (\sigma_{ij}) | \sigma_{ij} = \sigma_{ji} = 0, (i, j) \in \mathcal{E}\}$
- etc...

MLE for structured covariance matrices

- Negative logarithmic likelihood func (up to const.):

$$\ell(\Sigma) := -\log \det \Sigma^{-1} + \text{tr}(\Sigma^{-1}S) \rightarrow \min$$

- Rem Note that $-\log \det$ is a convex function.

- If \mathcal{M} is DA (inversely linear structure), then the **minimization** problem of $\ell(\Sigma)$ (**MLE**) s.t. $\Sigma \in \mathcal{M}$ is a **strictly convex** program.



Unique solution,

Numerically tractable (optimality eq. is linear)

App.(3) Convex program

Affine-scaling method and IG

- General convex program: Convex set $\mathcal{M} \subset \mathbf{R}^n$, $c \in \mathbf{R}^n$

$$\text{minimize } c^T x, \quad \text{s.t. } x \in \overline{\mathcal{M}}$$

- Ψ : a *good* convex barrier func. for \mathcal{M} ,

1) $\Psi(x) \rightarrow +\infty$ ($x \rightarrow \text{bd } \mathcal{M}$), 2) h : p.d. Hessian, 3) $+\alpha$

- Gradient flow for Riemannian mfd (\mathcal{M}, h)

$$\dot{x} = \frac{dx}{dt} = -h(x)^{-1}c, \quad x(0) \in \mathcal{M}$$

$x(t)$: *affine-scaling trajectory*

(numerically traced)



- Legendre transform \Rightarrow *linearized*

$$\dot{s} = -c, \quad s_i = \frac{\partial \Psi}{\partial x^i}, \quad i = 1, \dots, n, \quad \hat{s} := - \lim_{t \rightarrow +\infty} ct + s(0)_{,1}$$

- Opt. sol.: $\hat{x} = \text{grad}\Psi^*(\hat{s})$ (inverse Legendre trans.)
- **Red underlined**: needs the **explicit form of Ψ^***
(or solving the nonlinear eq.: $\hat{s} = \text{grad}\Psi(\hat{x})$)

Idea

- Ψ^* is known for a good barrier $\Psi \Rightarrow$ an explicit opt. sol. \hat{x}
- 1) Ω : sym. cones $\Rightarrow \psi(x) = -\log \det x$, $\psi^*(s) = -\log \det s$,
Legendre transform: $x \mapsto s = x^{-1}$
 - 2) \mathcal{M} realized by $\mathcal{M} = (a+W) \cap \Omega$ is DA in Ω
 \Rightarrow a) convexity of \mathcal{M} , b) linearized traj. belongs to \mathcal{M}

- Ex. SemiDefinite Program (SDP)

$$\underset{P}{\text{minimize}} (C|P), \text{ s.t. } P = E_0 + \sum_{i=1}^m x^i E_i \in \overline{\mathcal{M}} = \overline{(E_0 + \mathcal{W}) \cap \text{PD}(n)}$$

- If \mathcal{M} is DA in $\text{PD}(n)$ and $P \in \mathcal{M}$

- 1. Set $F^0 = P^{-1}, F^i = -P^{-1} E_i P^{-1}$, then

$$\mathcal{M} = \{P \mid P^{-1} = F^0 + \sum_{i=1}^m s_i F^i, \exists s = (s_i) \in \mathbf{R}^m\} \cap \text{PD}(n)$$

- 2. Solve $\tilde{C} \in \text{span}\{F^i\}_{i=1}^m$ meeting

$$\forall P \in \mathcal{M}, \quad (C|P) = (\tilde{C}|P) + \text{const.}$$

- 3. Spectral decomposition

$$\tilde{C} = (V_1 \ V_2) \begin{pmatrix} \Sigma_1 & O \\ O & O \end{pmatrix} \begin{pmatrix} V_1^T \\ V_2^T \end{pmatrix} = V_1 \Sigma_1 V_1^T$$

- 4. For $\forall P_0 \in \mathcal{M}$ with $S_0 = P_0^{-1}$, the opt. sol. is

$$\hat{P} = \lim_{t \rightarrow \infty} S(t)^{-1} = \lim_{t \rightarrow \infty} (-\tilde{C}t + S_0)^{-1} = P_0 - P_0 V_1 (V_1^T P_0 V_1)^{-1} V_1^T P_0$$

Rem. **Independent** of the objective function ($C|P$) and an initial value P_0

Interior point method (IP) for Conic linear program

Conic linear program -Notation-

- Vector space E of dimension n $E \ni x$
- The dual vector space E^* $E^* \ni s$
- $\langle s, x \rangle$: Paring
- Ω : proper open convex cone in E
- Ω^* : the dual cone of Ω
$$\Omega^* := \{s \in E^* \mid \langle s, x \rangle > 0, \forall x \in \overline{\Omega} \setminus \{0\}\}$$
- T^* : (Orthogonal) dual subspace of $T \subset E$
$$T^* = \{s \in E^* \mid \langle s, x \rangle = 0, \forall x \in T\}$$

Conic Linear Program

Given

$$c \in E^*, f \in E \text{ and } T \subset E$$

- Primal problem

(P) minimize $\langle c, x \rangle$, s.t. $x \in \overline{\mathcal{P}}$,
where $\mathcal{P} := (f + T) \cap \Omega$,

- Dual problem

(D) maximize $\langle s, f \rangle$, s.t. $s \in \overline{\mathcal{D}}$,
where $\mathcal{D} := (c + T^*) \cap \Omega^*$.

Typical Examples

- Linear program (**LP**):

$$E = E^* = \mathbf{R}^n, \quad \Omega = \Omega^* = \mathbf{R}_{++}^n$$

- Semidefinite program (**SDP**):

$E = E^*$: the set of real symmetric matrices

$\Omega = \Omega^*$: the set of positive definite matrices

- Second order cone (Lorentz cone) program (**SOCP**)
- Mixture of the aboves

ϑ -normal barrier on an open convex cone Ω

- Def. θ -normal barrier ψ on Ω ($\Leftarrow +\alpha$)

- A (smooth) convex function ψ satisfying, at each x in Ω ,

$$\psi(tx) = \psi(x) - \vartheta \log t,$$

$$|(D^2 d\psi)_x(X, X, X)| \leq 2((Dd\psi)_x(X, X))^{3/2}$$

for $\vartheta \geq 1$, $\forall t > 0$ and $\forall X \in T_x \Omega \cong E$

- $\psi(x) \rightarrow +\infty$ ($x \rightarrow \text{bd } \Omega$),

Rem. [Nesterov & Nemirovski 94] (1) Existence for all Ω (but not explicit forms), (2) the Hessian is p.d., (3) **self-concordance** (\Rightarrow the Newton method is efficient).

- Ex. $\psi(x) = -\sum_{i=1}^n \log x^i$ (LP), $\psi(x) = -\log \det X$ (SDP)

Dually flat structure on Ω (revisited)

- D : the canonical **flat** affine connection on E
- $\{x^1, \dots, x^n\}$: affine coordinate system, i.e.,
- g : Riemannian metric on Ω

$$D \frac{\partial}{\partial x^i} \frac{\partial}{\partial x^j} = 0$$

$$g = Dd\psi = \sum_{i,j} \frac{\partial^2 \psi}{\partial x^i \partial x^j} dx^i dx^j.$$

- D' : the dual affine connection on Ω
 $Xg(Y, Z) = g(D_X Y, Z) + g(Y, D'_X Z)$

(g, D, D') : dually flat structure on Ω

Remark

- $\{s_1, \dots, s_n\}$: dual coordinate system on E^* , s.t.

$$\langle s, x \rangle = \sum_i s_i(s) x^i(x)$$

- **Gradient map** $\iota : \Omega \rightarrow \Omega^*$ defined by

$$s_i \circ \iota = -\frac{\partial \psi}{\partial x^i}$$

induces **dually flat structure on Ω^*** from (g, D, D')

- (1) D^* : the canonical **flat** affine connection on E^*

$$D_{\iota_* X}^* \iota_* Y = \iota_* (D'_X Y) \quad (\iota^* D^* = D')$$

D^* -autoparallel in Ω^* \longleftrightarrow D' -autoparallel in Ω

Remark

(2) Riemannian metric $g^* := D^* d\psi^*$ on Ω^*

$$g = \iota^* g^*$$

(3) $\langle \iota_*(X), Y \rangle = -g_x(X, Y)$

Hessian norm : We denote the length of X in $T_x \Omega \cong E$ by

$$\|X\|_x := \|Z\|_s := \sqrt{g_x(X, X)} = \sqrt{g_s^*(Z, Z)},$$

where $s = \iota(x)$ and $Z = \iota_*(X)$.

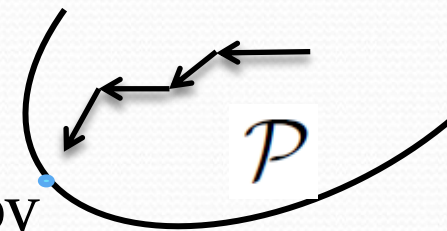
Curvature integral and iteration-complexity of IP

One of important computational performance indices for optimization algorithms is the **iteration-complexity**.

- Ω : **sym. cone** and $\mathcal{P} := (f + T) \cap \Omega$ is **DA**

\Rightarrow iteration-complexity=0 for (P)

- General case? Iter.-comp. is characterized by



- Curvature integrals along the **central trajectory** $\gamma_{\mathcal{P}}$

$$\int_{t_1}^{t_2} \|H_{\mathcal{P}}^*(\dot{\gamma}_{\mathcal{P}}(t), \dot{\gamma}_{\mathcal{P}}(t))\|_{\gamma_{\mathcal{P}}(t)}^{1/2} dt$$

- Similarly, for (D) curvature integrals along the dual c. t. $\gamma_{\mathcal{D}}$

$$\int_{t_1}^{t_2} \|H_{\mathcal{D}}(\dot{\gamma}_{\mathcal{D}}(t), \dot{\gamma}_{\mathcal{D}}(t))\|_{\gamma_{\mathcal{D}}(t)}^{1/2} dt$$

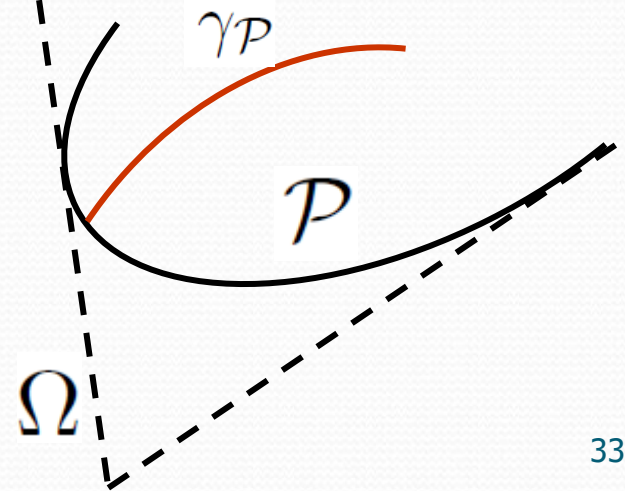
Central trajectory $\gamma_{\mathcal{P}}$

- **Primal problem:** minimize $\langle c, x \rangle$, s.t. $x \in \overline{\mathcal{P}}$,
where $\mathcal{P} := (f + T) \cap \Omega$,

- $x(t) := \gamma_{\mathcal{P}}(t)$: the unique minimizer of
minimize t $\langle c, x \rangle + \psi(x)$, s.t. $x \in \overline{\mathcal{P}}$.

for each $t > 0$

- $\gamma_{\mathcal{P}} := \{\gamma_{\mathcal{P}}(t) | t > 0\}$:
(Primal) **central trajectory**

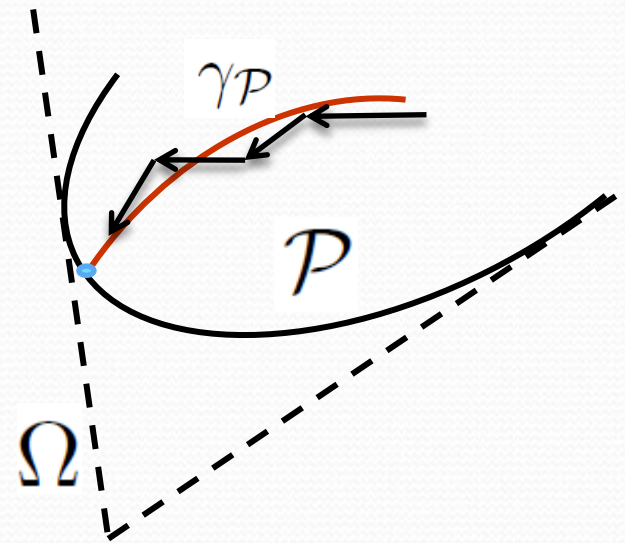


Central trajectory

- Homotopy path to the **opt. sol. of the primal problem**, i.e., $x(t)$ converges when $t \rightarrow \infty$.
- Numerically tracing $\gamma_{\mathcal{P}}$ is the standard and efficient way to solve the primal problem.

Path-following method

Idea: consider the problem in Ω^* and relate the complexity with the curvature



(1) Representation of feasible region

- A linear surj. operator $A : E \rightarrow \mathbf{R}^m$ s.t. $\text{Ker } A = T$

$$\mathcal{P} = \{x \in \Omega \mid Ax = b\},$$

$$\mathcal{D} = \{s \in \Omega^* \mid s = c - A^*y, y \in \mathbf{R}^m\}$$

where $A^* : \mathbf{R}^m \rightarrow E^*$ satisfying $y^T(Ax) = \langle A^*y, x \rangle$,

$$b := Af \in \mathbf{R}^m$$

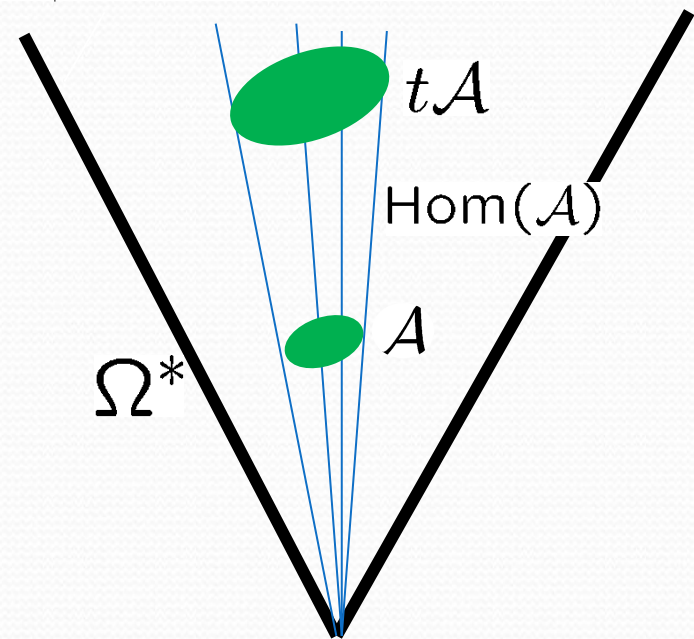
- $\dim \mathcal{P} = n - m$, $\dim \mathcal{D} = m$
- \mathcal{P} is D -autoparallel and \mathcal{D} is D^* -autoparallel

(2) Homogenization (conic hull)

- homogenization of \mathcal{D} in Ω^*

$$\text{Hom}(\mathcal{D}) := \bigcup_{t>0} t\mathcal{D}, \quad t\mathcal{D} := \{s \in \Omega^* \mid s = t\tilde{s}, \tilde{s} \in \mathcal{D}\}$$

- D^* -autoparallel
because \mathcal{D} is.
- $\dim \text{Hom}(\mathcal{D}) = m+1$



Homogenization

Lemma

The following relations hold in Ω^* :

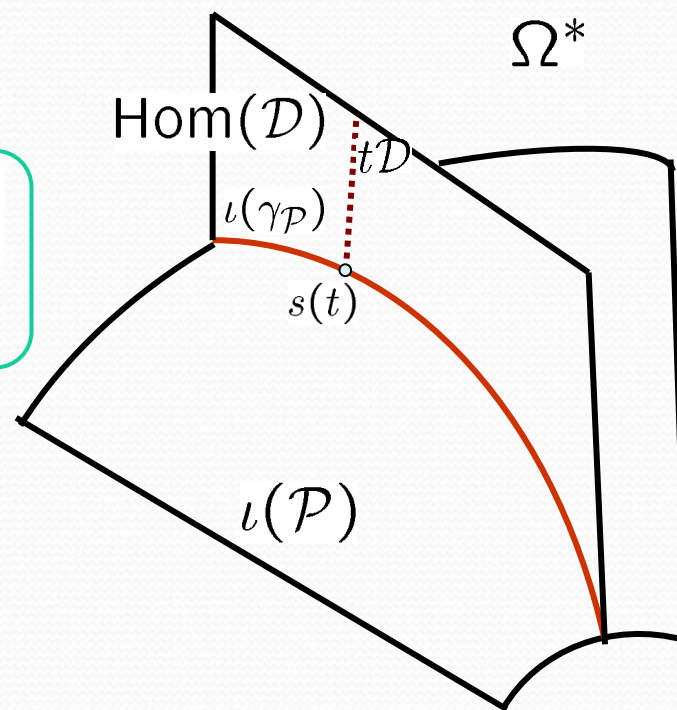
$$\iota(\gamma_{\mathcal{P}}) = \iota(\mathcal{P}) \cap \text{Hom}(\mathcal{D})$$

$$s(t) := \iota(x(t)) = \iota(\mathcal{P}) \cap t\mathcal{D}$$

$$\left[\begin{array}{l} L(x, y) := t\langle c, x \rangle + \psi(x) + y^T (b - Ax) \\ \partial L / \partial x = 0 \rightarrow s \in t\mathcal{D} \end{array} \right]$$

Remark

$\iota(\mathcal{P})$ and $t\mathcal{D}$ are orthogonal
w.r.t. g^* at $s(t)$ by definition.



3. Geometric predictor-corrector algorithm (tracing $\gamma_{\mathcal{P}}$ in $\text{Hom}(\mathcal{D})$)

Ideal case

- Predictor

From $s(t) \in \iota(\gamma_{\mathcal{P}})$

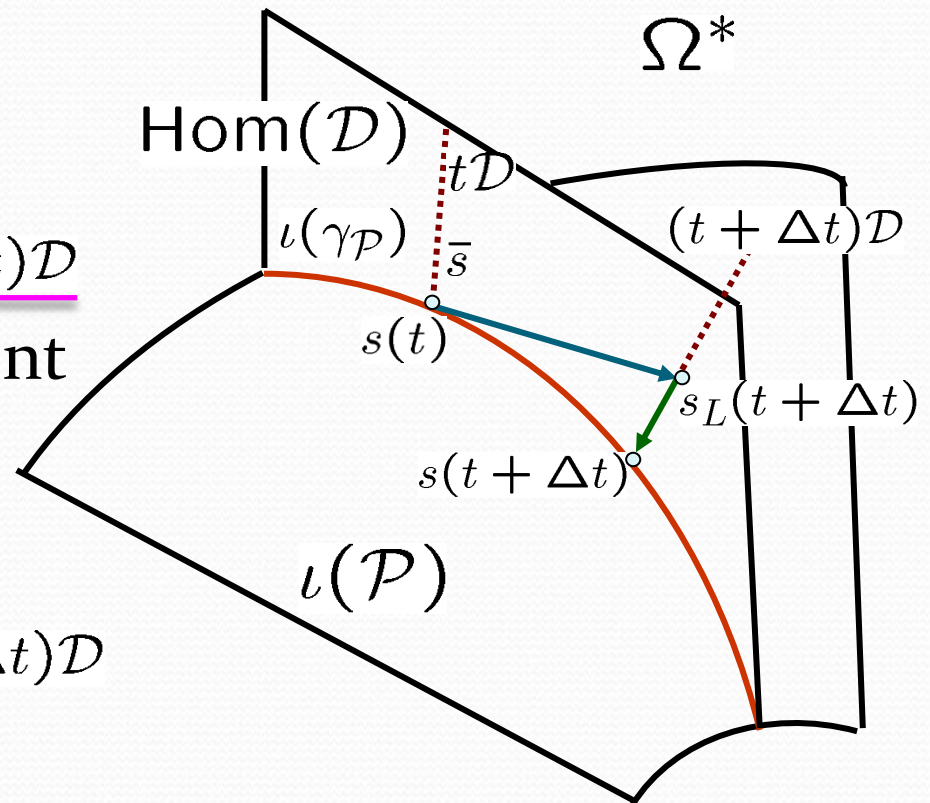
to $s_L(t + \Delta t) \in (t + \Delta t)\mathcal{D}$

with the direction tangent to $\iota(\gamma_{\mathcal{P}})$

- Corrector

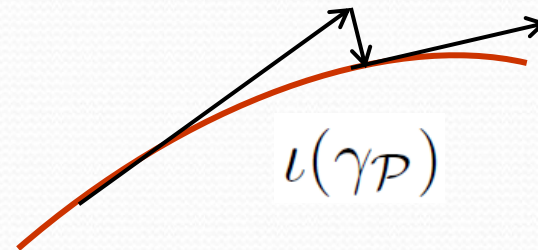
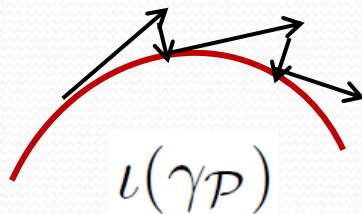
From $s_L(t + \Delta t) \in (t + \Delta t)\mathcal{D}$

to $s(t + \Delta t) \in \iota(\gamma_{\mathcal{P}})$



Intuitive observation

- $H_{\mathcal{P}}^*(\dot{\gamma}_{\mathcal{P}}(t), \dot{\gamma}_{\mathcal{P}}(t))$: the **Euler-Schouten embedding curvature (second fundamental form)** of $\iota(\gamma_{\mathcal{P}})$ with respect to D^*
- If $H_{\mathcal{P}}^*(\dot{\gamma}_{\mathcal{P}}(t), \dot{\gamma}_{\mathcal{P}}(t))$ is small at t , so is expected the iteration number !?

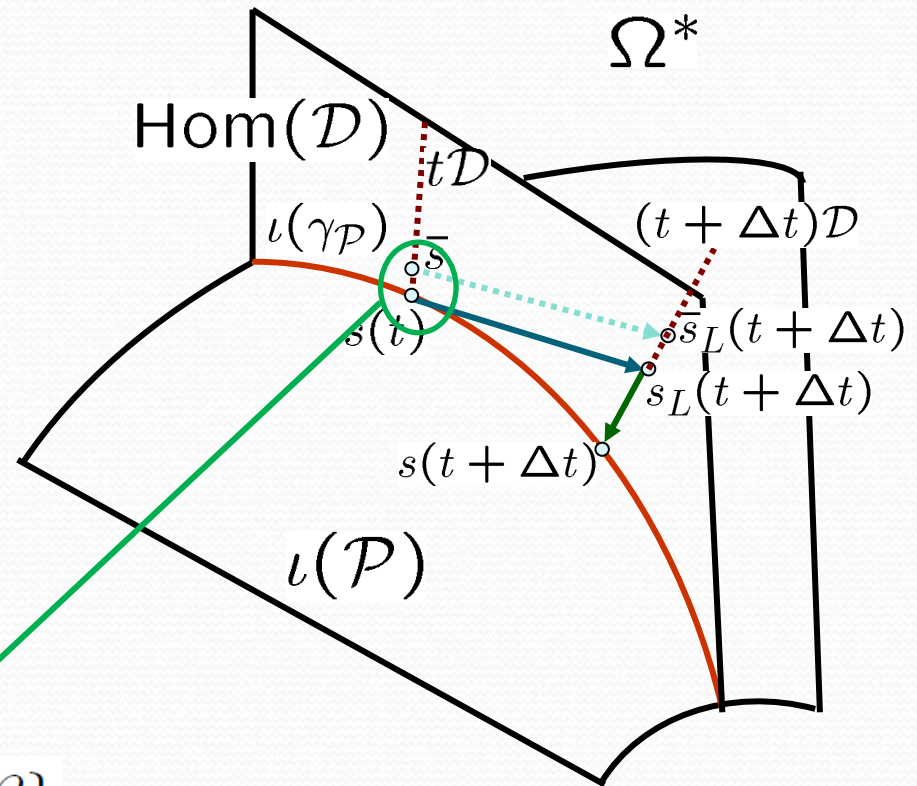


Actually,

$$\ddot{s} = D_{\dot{s}}^* \dot{s} = \iota_* (H_{\mathcal{P}}^*(\dot{\gamma}_{\mathcal{P}}, \dot{\gamma}_{\mathcal{P}}))$$

Remark: practical case

- Cannot expect that the corrector returns precisely on $\iota(\gamma_{\mathcal{P}})$
- Consider the point \bar{s} in the **neighborhood** of $s(t) \in \iota(\gamma_{\mathcal{P}})$ in the sense of Riemannian metric



$$\mathcal{N}_t(\beta) := \{s \in t\mathcal{D} \mid \delta(s) \leq \beta\}$$

Predictor

- The differential equation expressing $\iota(\gamma_{\mathcal{P}})$:

$$\dot{s} = (\text{id} - \underline{\Pi_s^\perp})c = \frac{1}{t}(\text{id} - \Pi_s^\perp)s$$

where Π_s^\perp is the orthogonal projection w.r.t. g^* from E^* to $T^* = \text{Range}A^*$ at s .

Note: $\Pi_s^\perp = 0 \Rightarrow$ ODE for the A-S traj. (up to sign)

- Hence, the predictor is defined by

$$\bar{s}_L(t + \Delta t) := \bar{s} + \Delta t(I - \Pi_{\bar{s}}^\perp)c \in (t + \Delta t)\mathcal{D}$$

Corrector

- Reduces to the following convex optimization on $t\bar{\mathcal{D}}$:

$$\text{minimize } F(s) := \langle s, f \rangle + \psi^*(s), \text{ s.t. } s \in t\bar{\mathcal{D}}$$

- **Newton direction** N for this opt. problem:

$$D^* dF(X, N) = -dF(X), \quad \forall X \in \mathcal{X}(t\mathcal{D})$$

- **Newton decrement**: measure of approximation of s

$$\delta(s) := \|N\|_s$$

- We define the corrector with a single **Newton step** by:

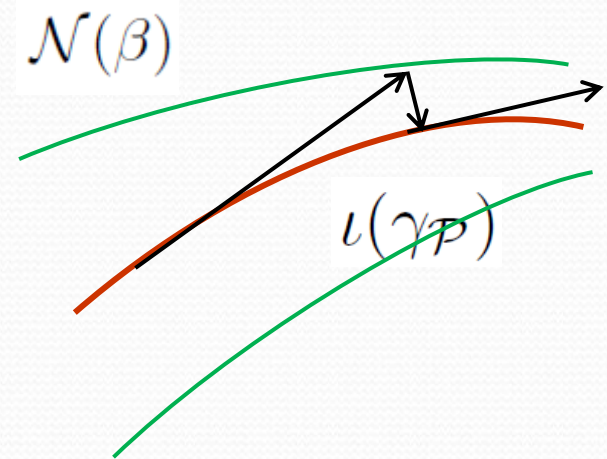
$$\bar{s}_L^+(t + \Delta t) := \bar{s}_L(t + \Delta t) + N_{\bar{s}_L}(t + \Delta t)$$

Tubular neighborhood

- The standard analysis technique in IP ensures the polynomiality of the complexity for this path-following strategy if all the generated points are near to $\iota(\gamma_{\mathcal{P}})$.
- Introduce the tubular neighborhood $\mathcal{N}(\beta)$ of $\iota(\gamma_{\mathcal{P}})$

$$\mathcal{N}(\beta) := \bigcup_{t \in (0, \infty)} \mathcal{N}_t(\beta),$$

where $\mathcal{N}_t(\beta) := \{s \in t\mathcal{D} \mid \delta(s) \leq \beta\}$.



4. Curvature integral and asymptotic iteration-complexity (Main result)

- **Assumption:** $\iota(\gamma_{\mathcal{P}})$ is **not** D^* -autoparallel, i.e.,
 $\beta \rightarrow 0$ implies that $\Delta t \rightarrow 0$ (If it is ?)

- **Theorem**

For $0 < t_1 < t_2$ and $s_1 \in \mathcal{N}(\beta) \cap t_1 \mathcal{D}$, let $\sharp(s_1, t_2, \beta)$ be **the iteration number** to find $s_2 \in \mathcal{N}(\beta) \cap t_2 \mathcal{D}$. Then,

$$\lim_{\beta \rightarrow 0} \frac{\sqrt{\beta} \times \sharp(s_1, t_2, \beta)}{I_{\mathcal{P}}(t_1, t_2)} = 1,$$

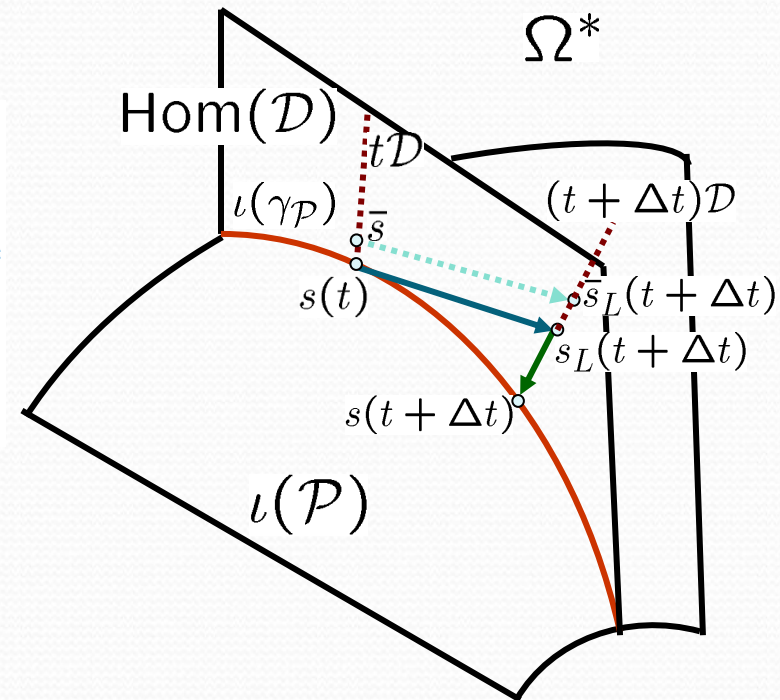
where

$$I_{\mathcal{P}}(t_1, t_2) := \frac{1}{\sqrt{2}} \int_{t_1}^{t_2} \|H_{\mathcal{P}}^*(\dot{\gamma}_{\mathcal{P}}(t), \dot{\gamma}_{\mathcal{P}}(t))\|_{\gamma_{\mathcal{P}}(t)}^{1/2} dt.$$

Outline of the proof

- Evaluate the Newton dec. of the predictor $\bar{s}_L(t + \Delta t)$ by $\|\ddot{s}(t)\|_{s(t)}$ (**For each iteration**)

$$\begin{aligned} & \delta(\bar{s}_L(t + \Delta t)) \\ &= \|s(t + \Delta t) - \bar{s}_L(t + \Delta t)\|_{\bar{s}_L(t + \Delta t)} + r_4 \\ &= \frac{(\Delta t)^2}{2} \|\ddot{s}(t)\|_{s(t)} + \delta(\bar{s}) + r_1 + r_2 + r_3, \end{aligned}$$



Outline of the proof

- Intermediate two relations for sufficiently small Δt and β .
(For each iteration)

- $$\begin{aligned} \sqrt{(1-\eta)\beta}(1-O(\sqrt{\beta})) &\leq \sqrt{w} - \sqrt{M_3}\delta(\bar{s}) \\ &\leq \frac{\Delta t}{\sqrt{2}} \|\ddot{s}(t)\|_{s(t)}^{1/2} + \sqrt{|r_1|} + \sqrt{M_3}(\Delta t)^2, \end{aligned}$$
- $$\begin{aligned} \frac{\Delta t}{\sqrt{2}} \|\ddot{s}(t)\|_{s(t)}^{1/2} - \sqrt{|r_1|} - \sqrt{M_3}(\Delta t)^2 \\ \leq \sqrt{w} + \sqrt{M_3}\delta(\bar{s}) \leq \sqrt{\beta}(1+O(\sqrt{\beta})) \end{aligned}$$

Outline of the proof

- Take summations of iterations

- $$\sqrt{(1-\eta)\beta} \sum_{k=1}^{\sharp(s_1, t_2, \beta)} (1 - O(\sqrt{\beta}))$$

$$\leq \frac{1}{\sqrt{2}} \int_{t_1}^{t_2} \|\ddot{s}(t)\|_{s(t)}^{1/2} dt + M' \sqrt{\Delta t_{\max}},$$

- $$\frac{1}{\sqrt{2}} \int_{t_1}^{t_2} \|\ddot{s}(t)\|_{s(t)}^{1/2} dt - M' \sqrt{\Delta t_{\max}}$$

$$\leq \sqrt{\beta} \sum_{k=1}^{\sharp(s_1, t_2, \beta)} (1 + O(\sqrt{\beta}))$$

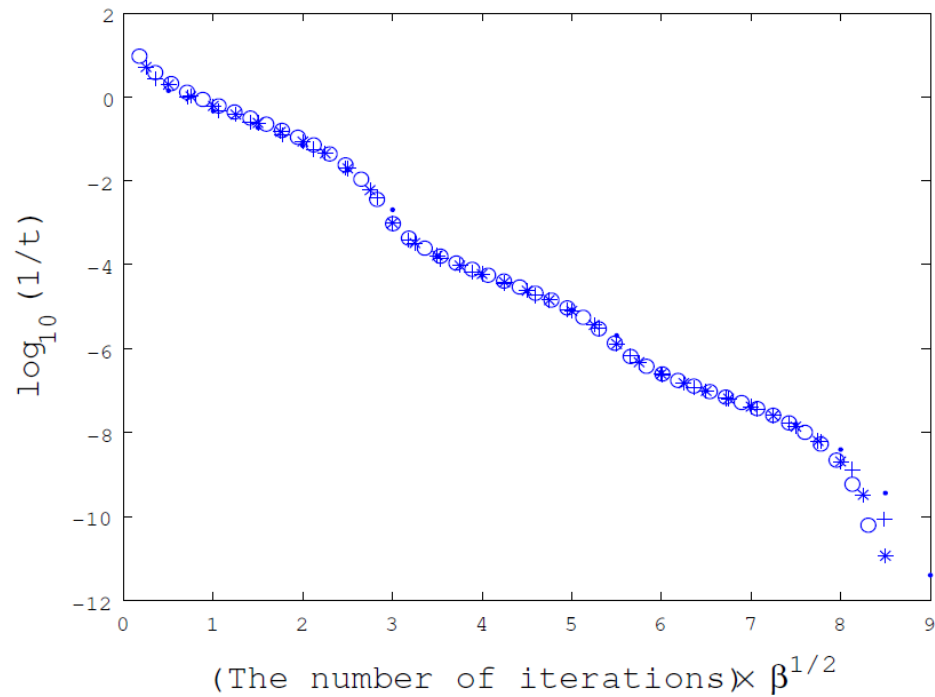
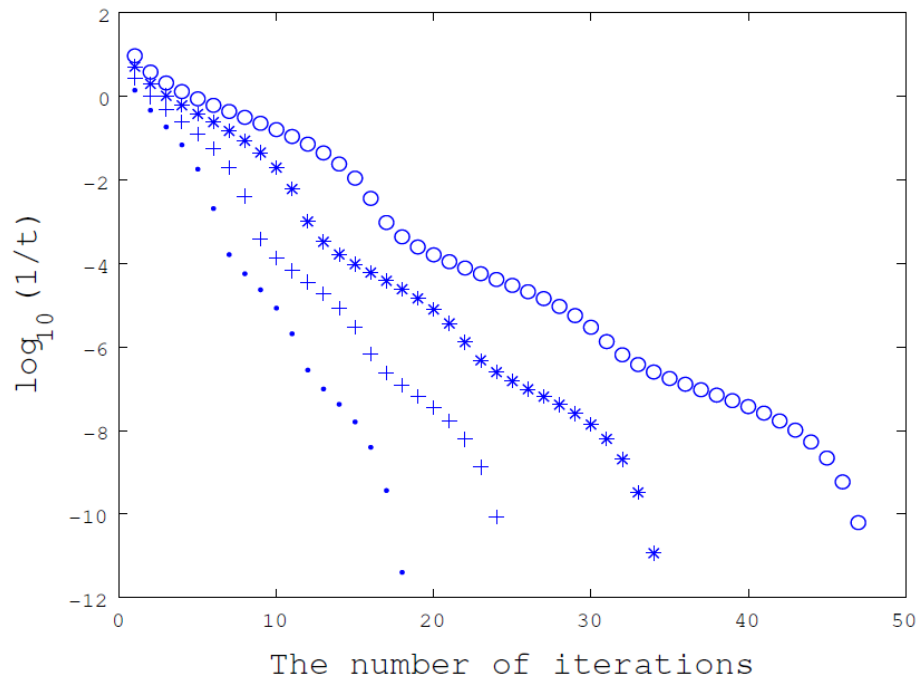
$$\ddot{s} = D_{\dot{s}}^* \dot{s} = \iota_* (H_{\mathcal{P}}^* (\dot{\gamma}_{\mathcal{P}}, \dot{\gamma}_{\mathcal{P}}))$$

Remark

- An asymptotic result for $\beta \rightarrow 0$ (and hence, $\Delta t \rightarrow 0$)
 - \mathcal{P} is DA $\Rightarrow \iota(\gamma_{\mathcal{P}})$ is DA (D^* -autoparallel) $\Rightarrow \Delta t \rightarrow \infty$
 \Rightarrow explicit sol.
- The same argument holds for the **dual** problem.
- The results are valid for general convex cones

Numerical experiment

- Curvature structure of CT for a certain LP



(\cdot : $\beta = 1/4$, $+$: $\beta = 1/8$, $*$: $\beta = 1/16$, \circ : $\beta = 1/32$)

Curved part is Straight and Straight part is Curved?(1)

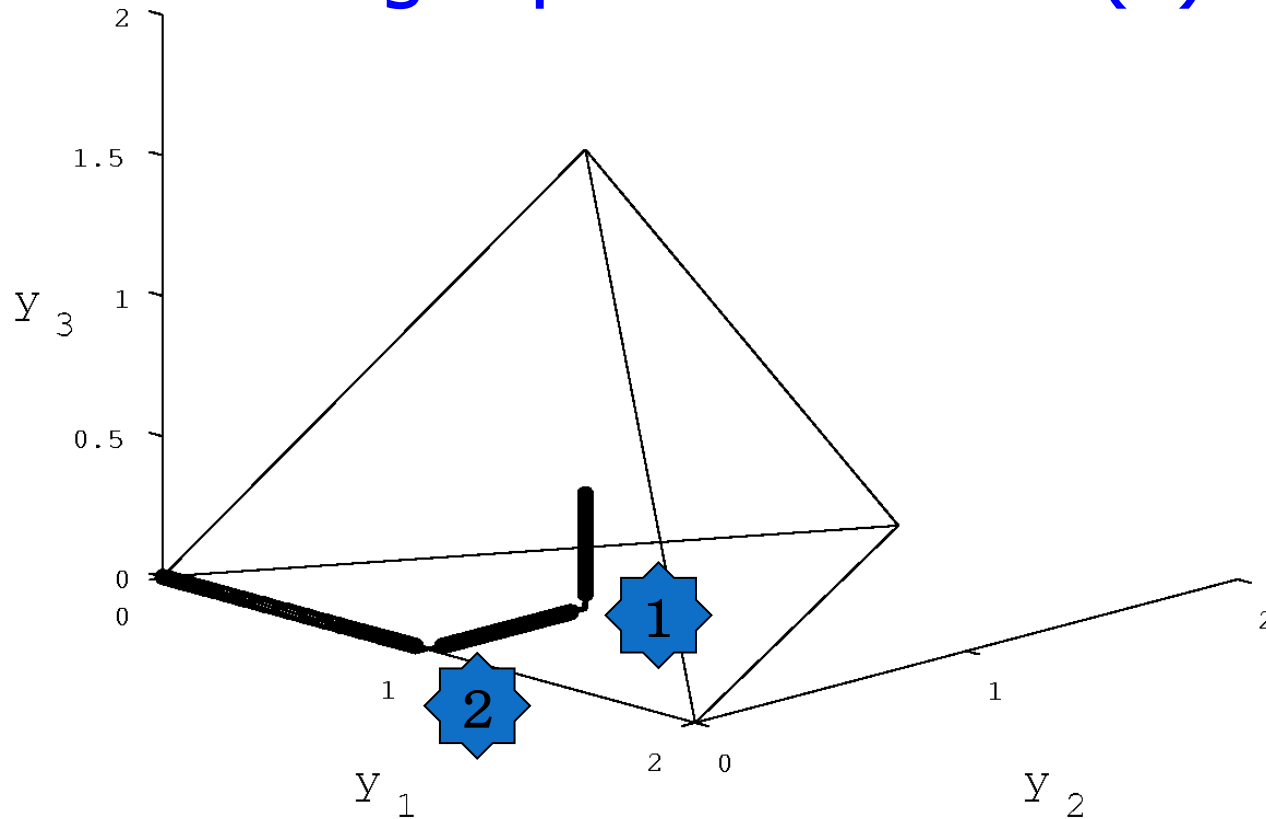


Figure 7: Exaple Figure

Proposition

It holds that $\|H_{\mathcal{P}}^*(\dot{\gamma}_{\mathcal{P}}, \dot{\gamma}_{\mathcal{P}})\|_{\gamma_{\mathcal{P}}(t)}^{1/2} \leq \frac{\sqrt{2\vartheta}}{t}$

ϑ : a constant determined by $\psi(x)$

- Remark

The above proposition gives the upper bound:

$$I_{\mathcal{P}}(t_1, t_2) \leq \sqrt{\vartheta} \log(t_2/t_1)$$

Further study for LP case

- Primal and Dual **Linear Program**:

$$\min c^T x$$

$$\text{s.t. } Ax = b, \quad x \geq 0, \quad A \in R^{m \times n}, \quad b \in R^n$$

$$\max b^T y$$

$$\text{s.t. } s = c - A^T y, \quad s \geq 0,$$

Application to

Primal-dual path-following (PDPF) method

- current **main-stream** IP (cheap in each iteration)
- The following quantity has been known to play an important and similar role in complexity analysis of **PDPF** method:

$$I_{PD}(t_1, t_2) = \int_{t_1}^{t_2} h_{PD}(t)^{1/2} dt$$

where $h_{PD}(t)$ is given by

$$h_{PD}(t) := \frac{1}{t^2} ((I_n - Q(t))e) \underline{*} (Q(t)e).$$

e : the unit element of Jordan product *

$Q(t)$: a certain projection matrix

Proposition

It holds that

$$h_{PD}(t)^2 = \left(\frac{1}{2} \|H_{\mathcal{P}}^*(\dot{\gamma}_{\mathcal{P}}(t), \dot{\gamma}_{\mathcal{P}}(t))\|_{\gamma_{\mathcal{P}}(t)} \right)^2 + \left(\frac{1}{2} \|H_{\mathcal{D}}(\dot{\gamma}_{\mathcal{D}}(t), \dot{\gamma}_{\mathcal{D}}(t))\|_{\gamma_{\mathcal{D}}(t)} \right)^2$$

Remark :

- **geometric implication** of the quantity of $I_{PD}(t_1, t_2)$

- inequalities
$$\max\{I_{\mathcal{P}}(t_1, t_2), I_{\mathcal{D}}(t_1, t_2)\} \leq I_{PD}(t_1, t_2) \leq I_{\mathcal{P}}(t_1, t_2) + I_{\mathcal{D}}(t_1, t_2).$$

Concluding Remark

- Tractable characterization of **DA** submfd's in symmetric cones Ω
- Application to conic linear programs
 - Explicit sol. when the feasible region M is DA in Ω .
 - M is DA \Rightarrow **AS (CT) traj. is DA** (D^* -autoparallel) $\Rightarrow \Delta t \rightarrow \infty$
 \Rightarrow explicit sol.
- **Extension**: # of iterations and **curvature integral** of CT
 - Asymptotic analysis ($\beta \rightarrow 0$)
 - Complemented by numerical experiment for finite β
 - Geometric structure of CT has a influence on complexity of the IP algorithm

- Relation among iteration-complexities of **P. D. and PD algorithm.**
- DA submanifolds in a certain submfd in Jordan algebras [OIT]
- Future work: **Geometrical** study for general stat. mfd.
 - Various geometrical concepts for mutually dual connections and their characterizations (Furuhata *et al.*)
 - Classifications
 - Families of continuous probability densities
 - Applications (Ex. Study of ODE's on manifolds?)

Thank you for your attention

References:

[OIT] A. Ohara, H. Ishi and T. Tsuchiya,

Doubly autoparallel structure and curvature integrals

- An application to iteration-complexity analysis of convex optimization -, Information Geometry, to appear.