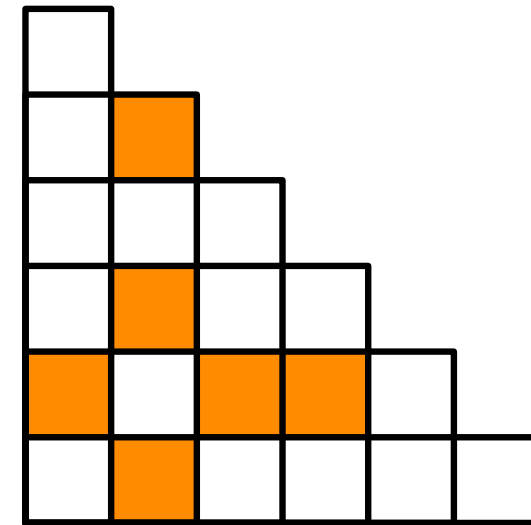
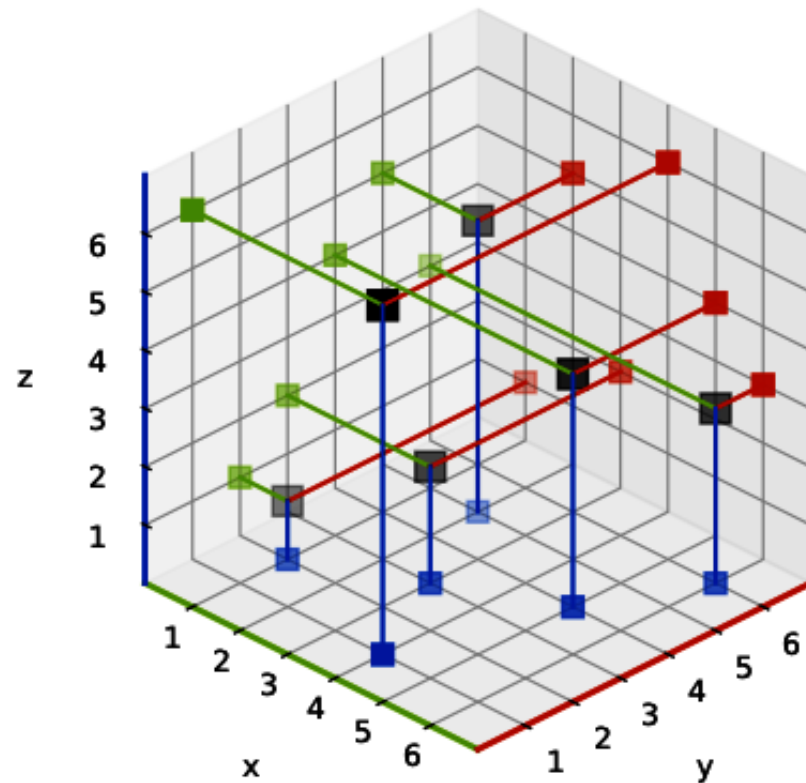


Pattern avoiding 3D-permutations and triangle bases

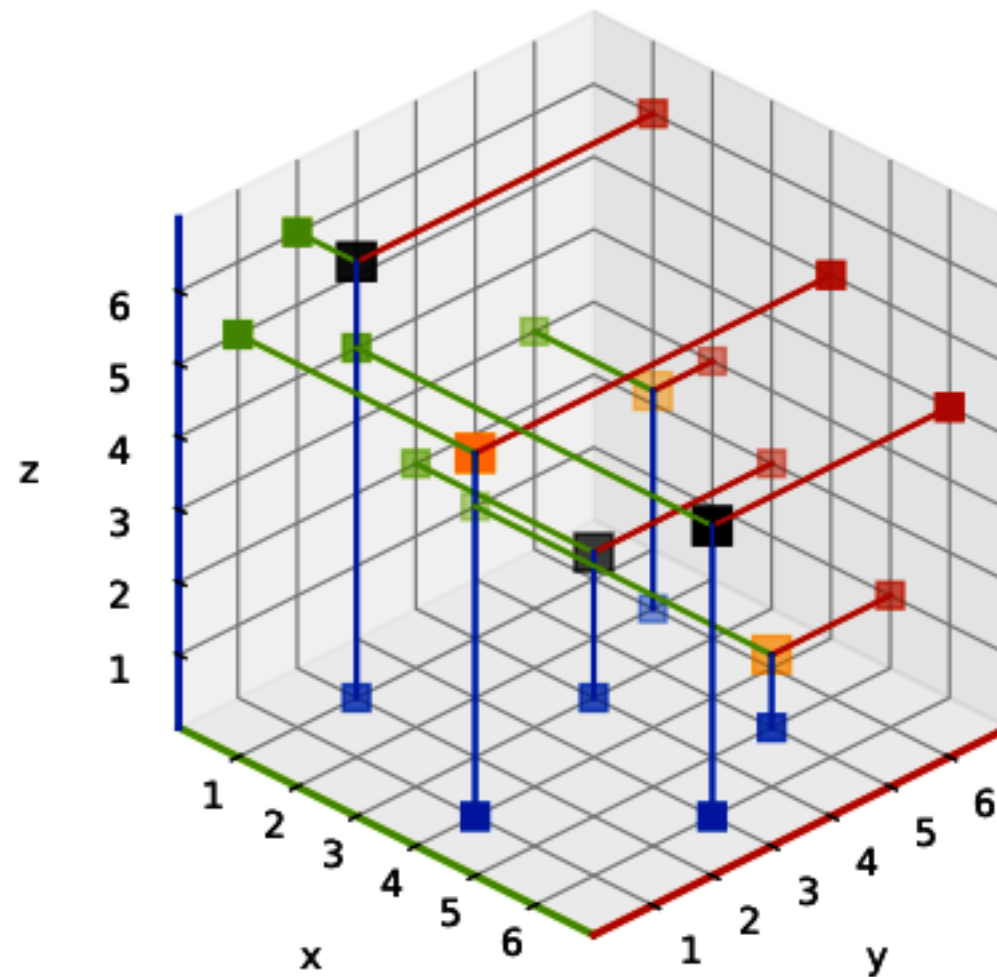
Juliette Schabanel

LaBRI, Université de Bordeaux, France



I- The objects

a) Pattern avoiding 3-permutations

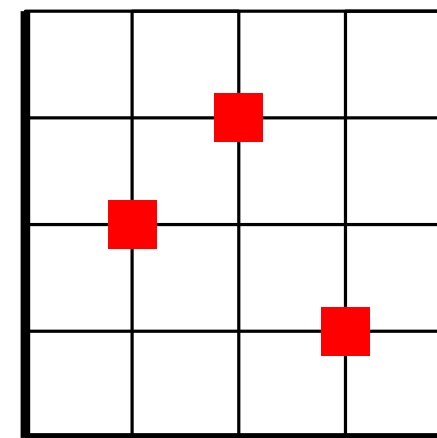
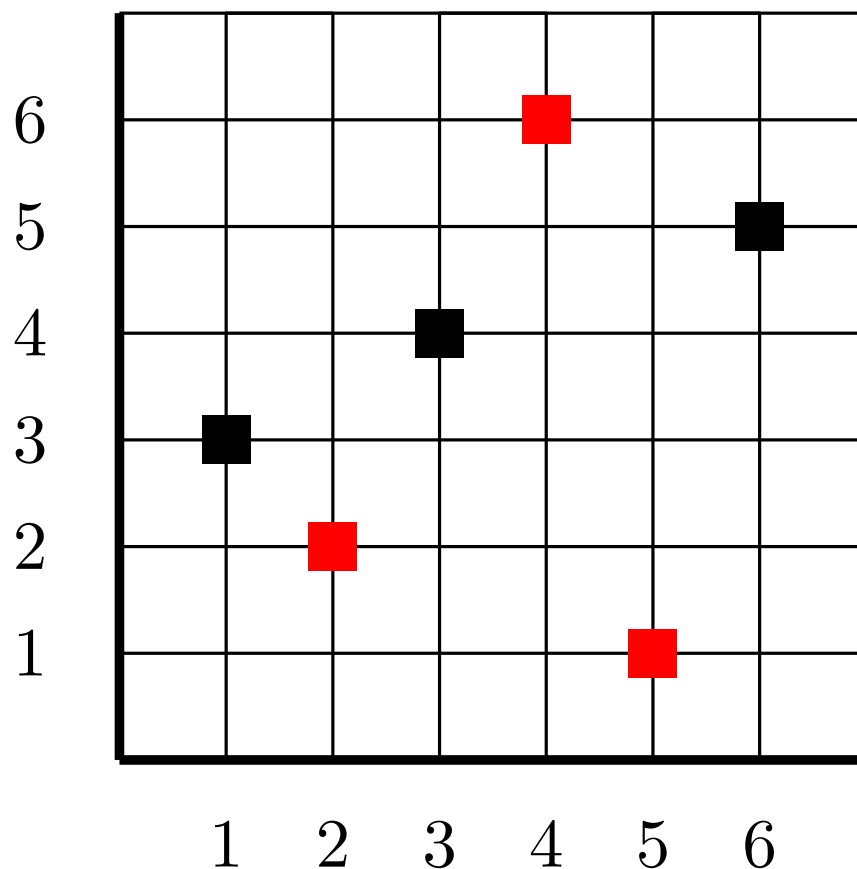


Pattern avoidance in permutations

The **diagram** of a permutation $\sigma \in \mathfrak{S}_n$ is the set of points $P_\sigma = \{(i, \sigma(i)) \mid 1 \leq i \leq n\}$. It has exactly one point per row and per column.

A permutation $\sigma \in \mathfrak{S}_n$ **contains** a pattern $\pi \in \mathfrak{S}_k$ if there is a set of indices I such that $\sigma|_I \simeq \pi$. Otherwise, it **avoids** it.

$Av_n(\pi_1, \dots, \pi_k) = \{\sigma \in \mathfrak{S}_n \text{ avoiding all the } \pi_i\text{'s}\}$. [Knuth, Kitaev, Mansour, ...]



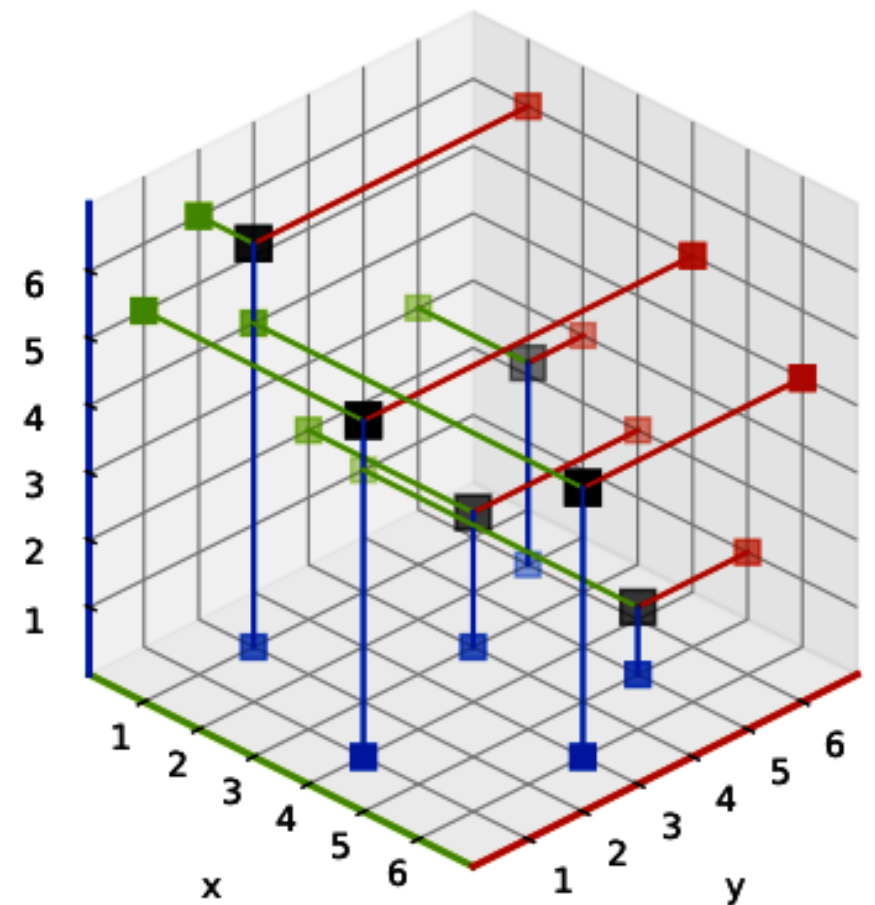
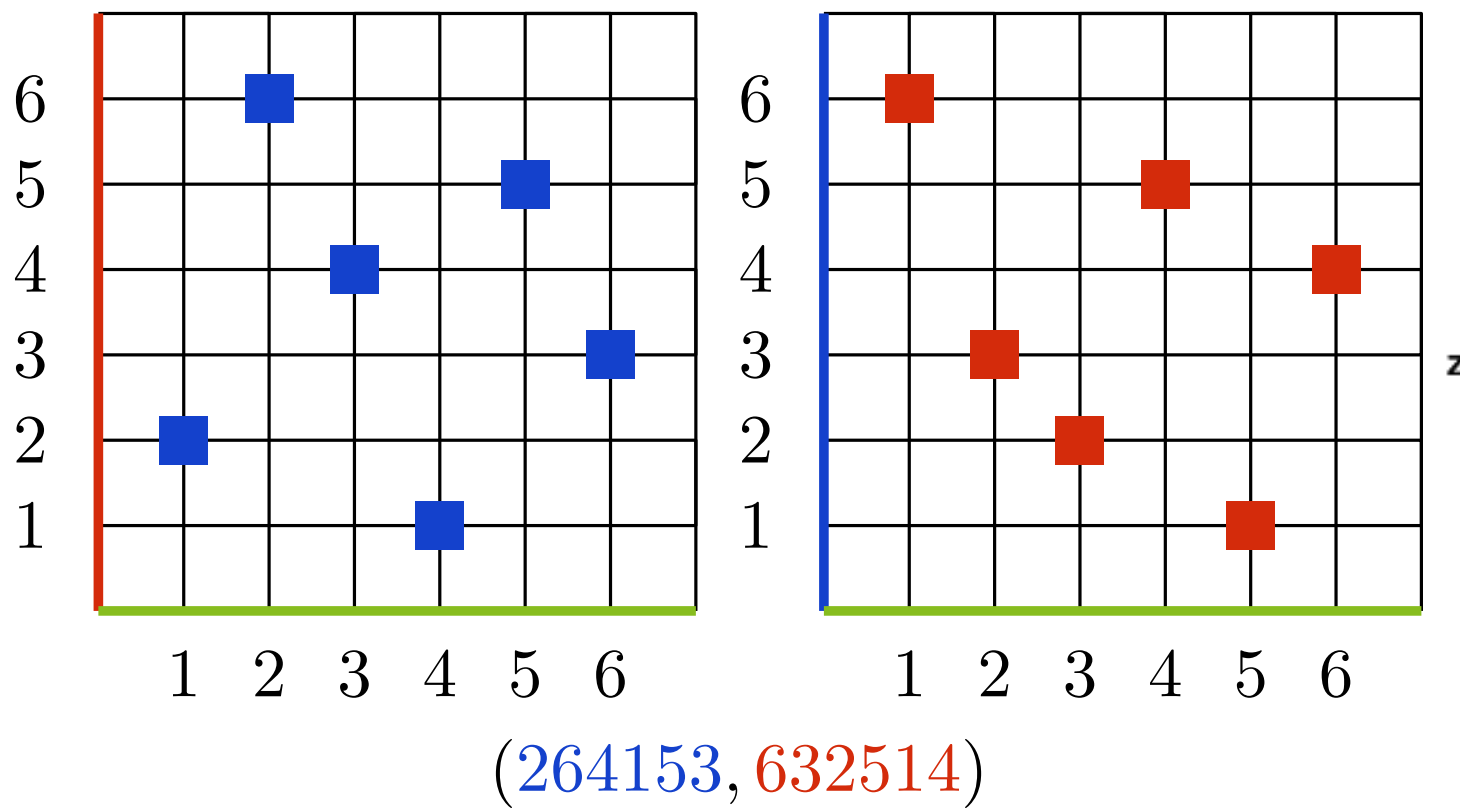
$\sigma = 324615$ contains the pattern $\pi = 231$.

Pattern avoidance in 3-permutations

A **3-diagram** has exactly one point per plane of the grid.

It is coded by a **3-permutation** $(\sigma, \tau) \in \mathfrak{S}_n^2$:

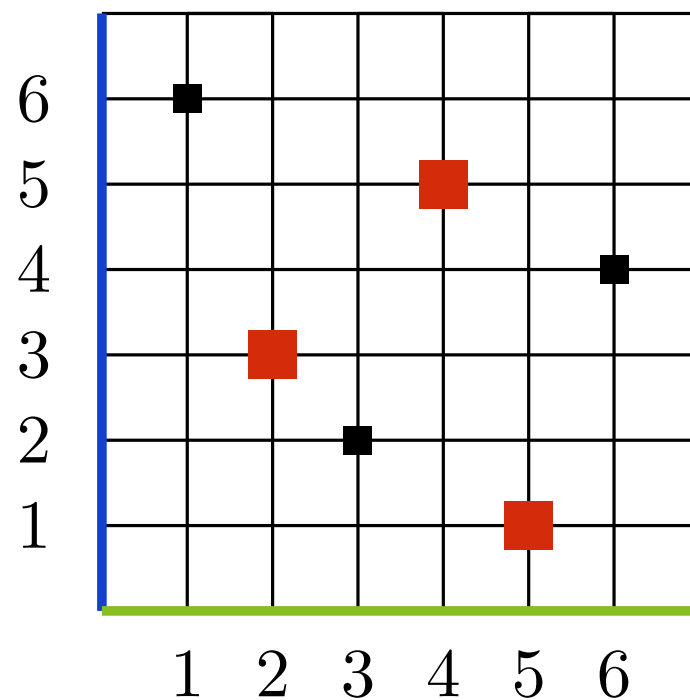
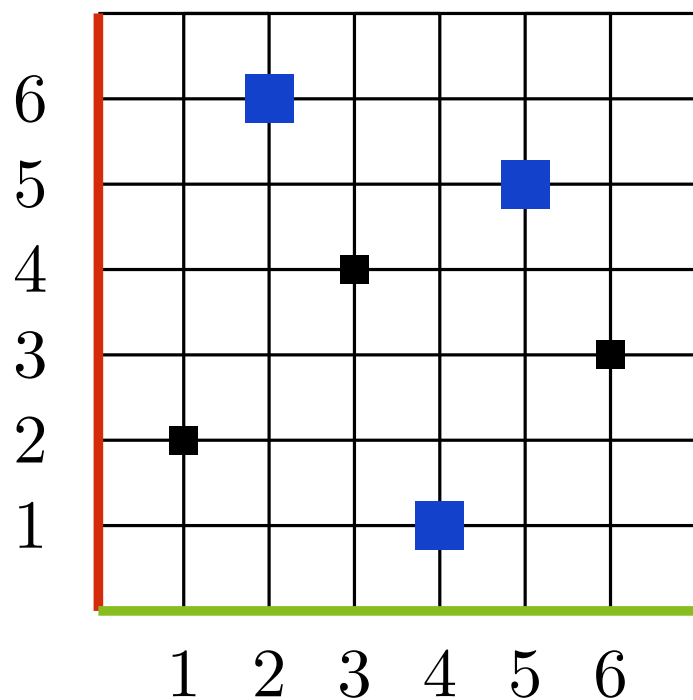
$$P_{(\sigma, \tau)} = \{(i, \sigma(i), \tau(i)) \mid 1 \leq i \leq n\}.$$



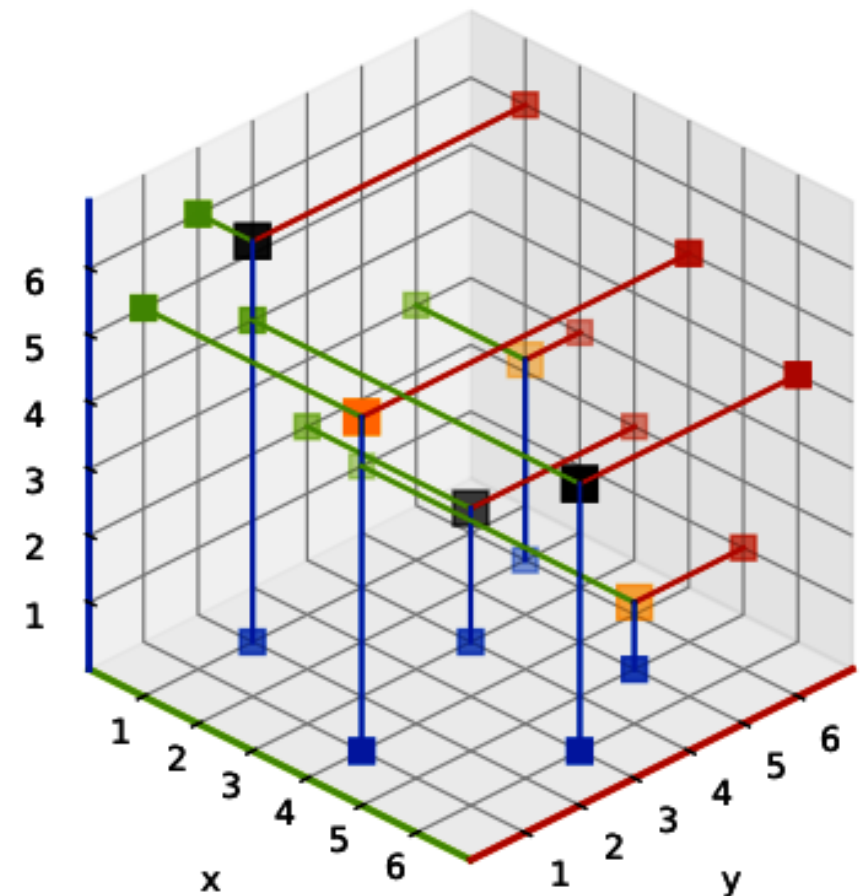
Pattern avoidance in 3-permutations

A **3-diagram** has exactly one point per plane of the grid.

A 3-permutation $(\sigma, \tau) \in \mathfrak{S}_n^2$ **contains** a pattern $(\pi_1, \pi_2) \in \mathfrak{S}_k^2$ if there is a set of indices $I \subset \llbracket 1, n \rrbracket$ such that $\sigma|_I \simeq \pi_1$ and $\tau|_I \simeq \pi_2$. Otherwise it **avoids** it.



$(\textcolor{blue}{264153}, \textcolor{red}{632514})$ contains the pattern $(\textcolor{blue}{312}, \textcolor{red}{231})$.



Pattern avoidance classes

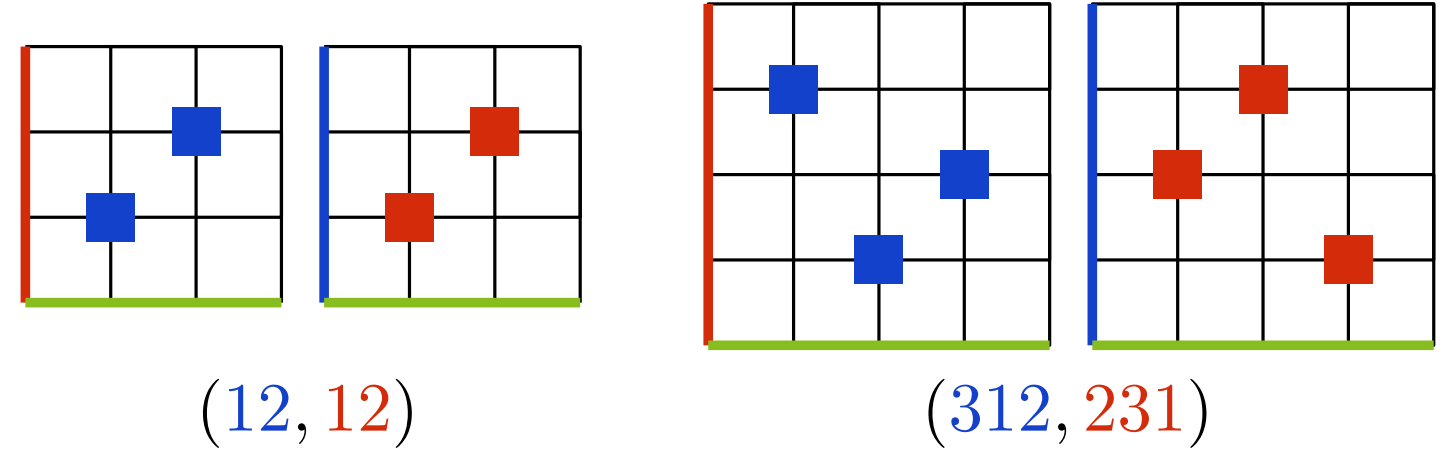
Patterns	TWE	Sequence	Comment
$(\textcolor{blue}{12}, \textcolor{red}{12})$	4	$1, 3, 17, 151, 1899, 31711, \dots$	weak-Bruhat intervals
$(\textcolor{blue}{12}, \textcolor{red}{12}), (\textcolor{blue}{12}, \textcolor{red}{21})$	6	$n! = 1, 2, 6, 24, 120 \dots$	$\sigma_1 \Rightarrow \sigma_2$
$(\textcolor{blue}{12}, \textcolor{red}{12}), (\textcolor{blue}{12}, \textcolor{red}{21}),$ $(\textcolor{blue}{21}, \textcolor{red}{12})$	4	$1, 1, 1, 1, 1, 1, \dots$	1 diagonal
$(\textcolor{blue}{12}, \textcolor{red}{12}), (\textcolor{blue}{12}, \textcolor{red}{21}),$ $(\textcolor{blue}{21}, \textcolor{red}{12}), (\textcolor{blue}{21}, \textcolor{red}{21})$	1	$1, 0, 0, 0, 0, 0, \dots$	
$(\textcolor{blue}{123}, \textcolor{red}{123})$	4	$1, 4, 35, 524, 11774, 366352, \dots$	<i>new</i>
$(\textcolor{blue}{123}, \textcolor{red}{132})$	24	$1, 4, 35, 524, 11768, 365558, \dots$	<i>new</i>
$(\textcolor{blue}{132}, \textcolor{red}{213})$	8	$1, 4, 35, 524, 11759, 364372, \dots$	<i>new</i>
$(\textcolor{blue}{12}, \textcolor{red}{12}), (\textcolor{blue}{132}, \textcolor{red}{312})$	48	$(n+1)^{n-1} = 1, 3, 16, 125, 1296 \dots$	[Atkinson et al. 93,95]
$(\textcolor{blue}{12}, \textcolor{red}{12}), (\textcolor{blue}{123}, \textcolor{red}{321})$	12	$1, 3, 16, 124, 1262, 15898, \dots$	distributive lattices inter.
$(\textcolor{blue}{12}, \textcolor{red}{12}), (\textcolor{blue}{231}, \textcolor{red}{312})$	8	$1, 3, 16, 122, 1188, 13844, \dots$	A295928?

[Bonichon & Morel '22]

Pattern avoidance classes

Patterns	TWE	Sequence	Comment
$(\textcolor{blue}{12}, \textcolor{red}{12})$	4	$1, 3, 17, 151, 1899, 31711, \dots$	weak-Bruhat intervals
$(\textcolor{blue}{12}, \textcolor{red}{12}), (\textcolor{blue}{12}, \textcolor{red}{21})$	6	$n! = 1, 2, 6, 24, 120 \dots$	$\sigma_1 \Rightarrow \sigma_2$
$(\textcolor{blue}{12}, \textcolor{red}{12}), (\textcolor{blue}{12}, \textcolor{red}{21}), (\textcolor{blue}{21}, \textcolor{red}{12})$	4	$1, 1, 1, 1, 1, 1, \dots$	1 diagonal
$(\textcolor{blue}{12}, \textcolor{red}{12}), (\textcolor{blue}{12}, \textcolor{red}{21}), (\textcolor{blue}{21}, \textcolor{red}{12}), (\textcolor{blue}{21}, \textcolor{red}{21})$	1	$1, 0, 0, 0, 0, 0, \dots$	
$(\textcolor{blue}{123}, \textcolor{red}{123})$	4	$1, 4, 35, 524, 11774, 366352, \dots$	<i>new</i>
$(\textcolor{blue}{123}, \textcolor{red}{132})$	24	$1, 4, 35, 524, 11768, 365558, \dots$	<i>new</i>
$(\textcolor{blue}{132}, \textcolor{red}{213})$	8	$1, 4, 35, 524, 11759, 364372, \dots$	<i>new</i>
$(\textcolor{blue}{12}, \textcolor{red}{12}), (\textcolor{blue}{132}, \textcolor{red}{312})$	48	$(n+1)^{n-1} = 1, 3, 16, 125, 1296 \dots$	[Atkinson et al. 93,95]
$(\textcolor{blue}{12}, \textcolor{red}{12}), (\textcolor{blue}{123}, \textcolor{red}{321})$	12	$1, 3, 16, 124, 1262, 15898, \dots$	distributive lattices inter.
$(\textcolor{blue}{12}, \textcolor{red}{12}), (\textcolor{blue}{231}, \textcolor{red}{312})$	8	$1, 3, 16, 122, 1188, 13844, \dots$	A295928?

[Bonichon & Morel '22]



Pattern avoidance classes

A295928 Number of triangular matrices $T(n,i,k)$, $k \leq i \leq n$, with entries "0" or "1" with the property that each triple $\{T(n,i,k), T(n,i,k+1), T(n,i-1,k)\}$ containing a single "0" can be successively replaced by $\{1, 1, 1\}$ until finally no "0" entry remains.

1, 3, 16, 122, 1188, 13844, 185448, 2781348, 45868268

([list](#); [graph](#); [refs](#); [listen](#); [history](#); [text](#); [internal format](#))

OFFSET 1,2

COMMENTS A triple $\{T(n,i,k), T(n,i,k+1), T(n,i-1,k)\}$ will be called a primitive triangle. It is easy to see that $b(n) = n(n-1)/2$ is the number of such triangles. At each step, exactly one primitive triangle is completed (replaced by $\{1, 1, 1\}$). So there are $b(n)$ "0"- and n "1"-terms. Thus the starting matrix has no complete primitive triangle. Furthermore, any triangular submatrix $T(m,i,k)$, $k \leq i \leq m < n$ cannot have more than m "1"-terms because otherwise it would have less "0"-terms than primitive triangles. The replacement of at least one "0"-term would complete more than one primitive triangle. This has been excluded.

So $T(n, i, k)$ is a special case of $U(n, i, k)$, described in [A101481](#): $a(n) < A101481(n+1)$.

A start matrix may serve as a pattern for a number wall used on worksheets for elementary mathematics, see link "Number walls". That is why I prefer the more descriptive name "fill matrix".

The algorithm for the sequence is rather slow because each start matrix is constructed separately. There exists a faster recursive algorithm which produces the same terms and therefore is likely to be correct, but it is based on a conjecture. For the theory of the recurrence, see "Recursive aspects of fill matrices". Probable extension $a(10)$ – $a(14)$: 821096828, 15804092592, 324709899276, 7081361097108, 163179784397820.

The number of fill matrices with n rows and all "1"- terms concentrated on the last two rows, is [A001960](#)(n). See link "Reconstruction of a sequence".

LINKS

[Table of \$n, a\(n\)\$ for \$n=1..9\$.](#)

Gerhard Kirchner, [Recursive aspects of fill matrices](#)

Gerhard Kirchner, [Number walls](#)

Gerhard Kirchner, [VB-program](#)

Gerhard Kirchner, [Reconstruction of a sequence](#)

Ville Salo, [Cutting Corners](#), arXiv:2002.08730 [math.DS], 2020.

Yuan Yao and Fedir Yudin, [Fine Mixed Subdivisions of a Dilated Triangle](#), arXiv:2402.13342 [math.CO], 2024.

EXAMPLE

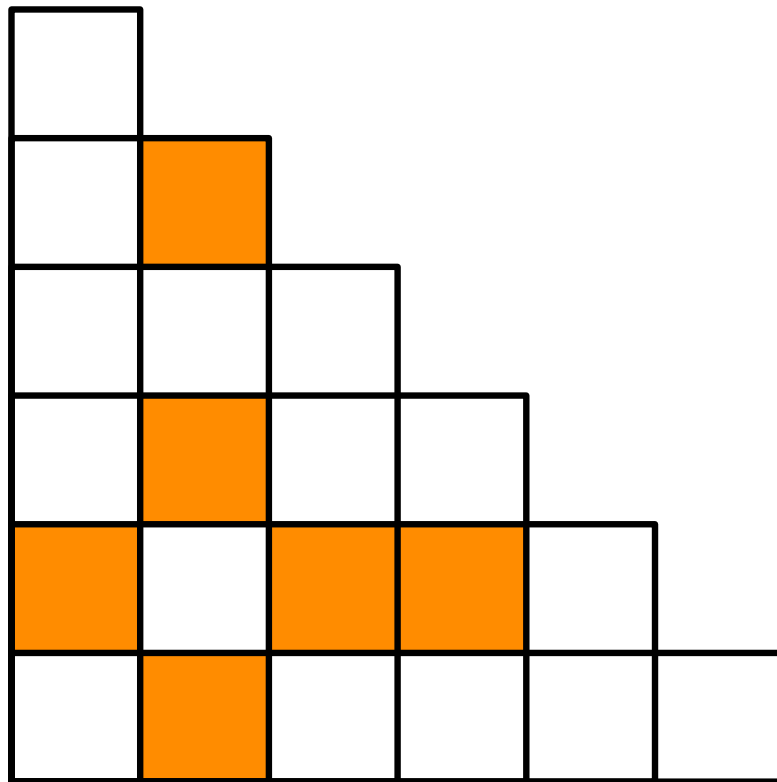
Example ($n=2$):
 $a(2)=3$

Example for completing a 3-matrix (3 bottom terms):

$$\begin{array}{cccc} 1 & & 1 & & 1 \\ 0 & 0 & \rightarrow & 1 & 0 \rightarrow & 1 & 1 \rightarrow & 1 & 1 \\ 1 & 1 & 0 & & 1 & 1 & 0 & & 1 & 1 & 1 \end{array}$$

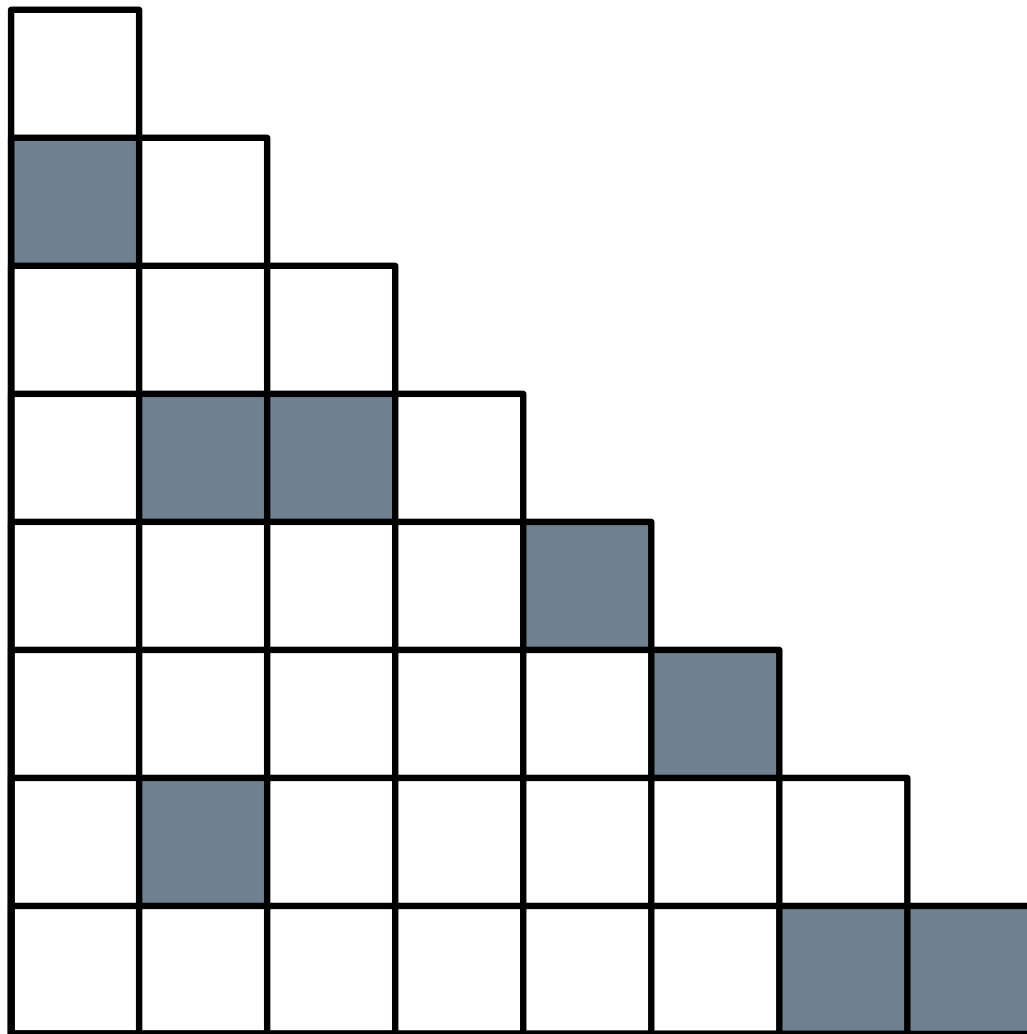
I- The objects

b) Triangle Bases



Filling configurations

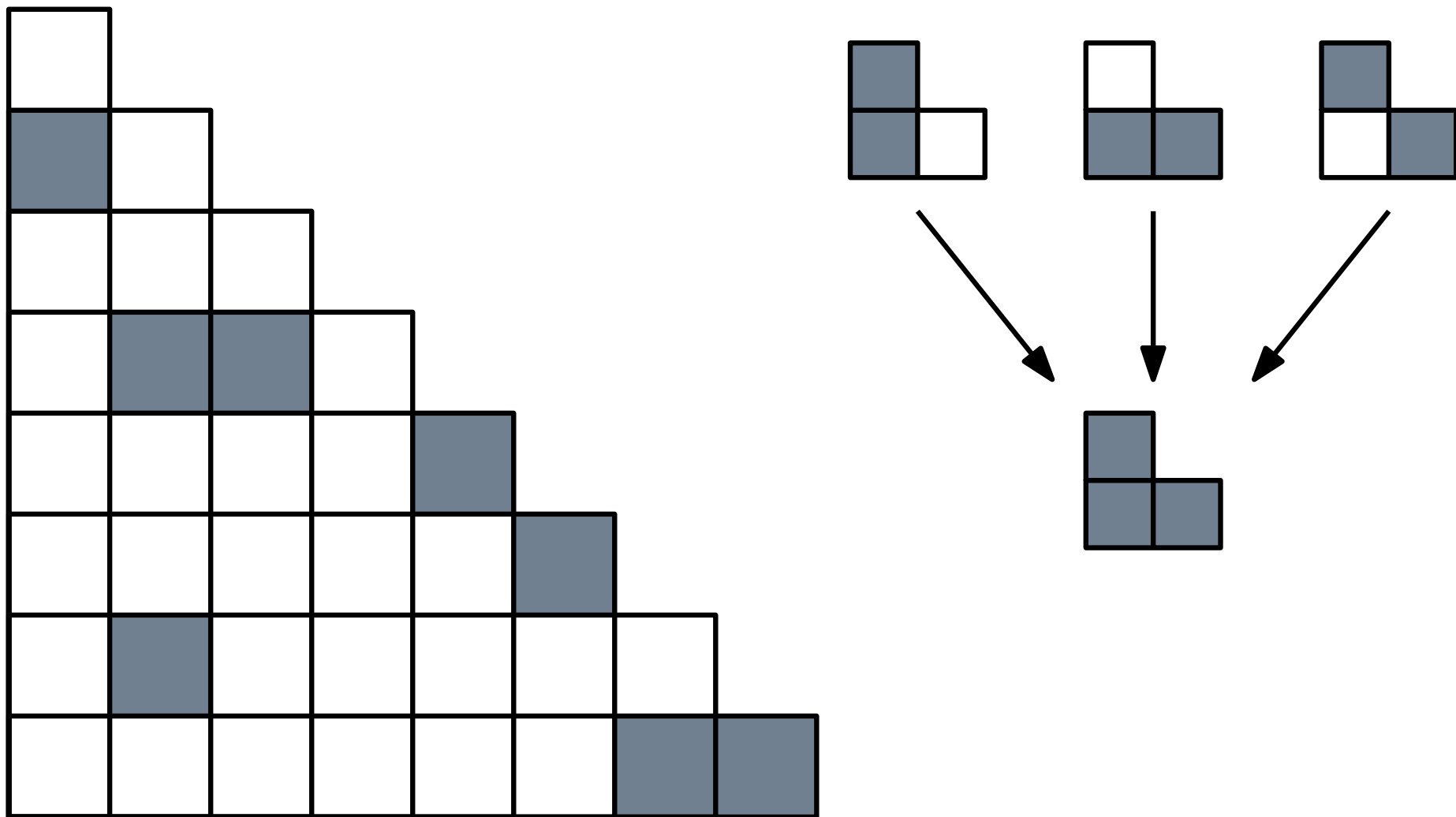
A **configuration** of size n is a set of n cells in the triangle T_n of size n .



Filling configurations

A **configuration** of size n is a set of n cells in the triangle T_n of size n .

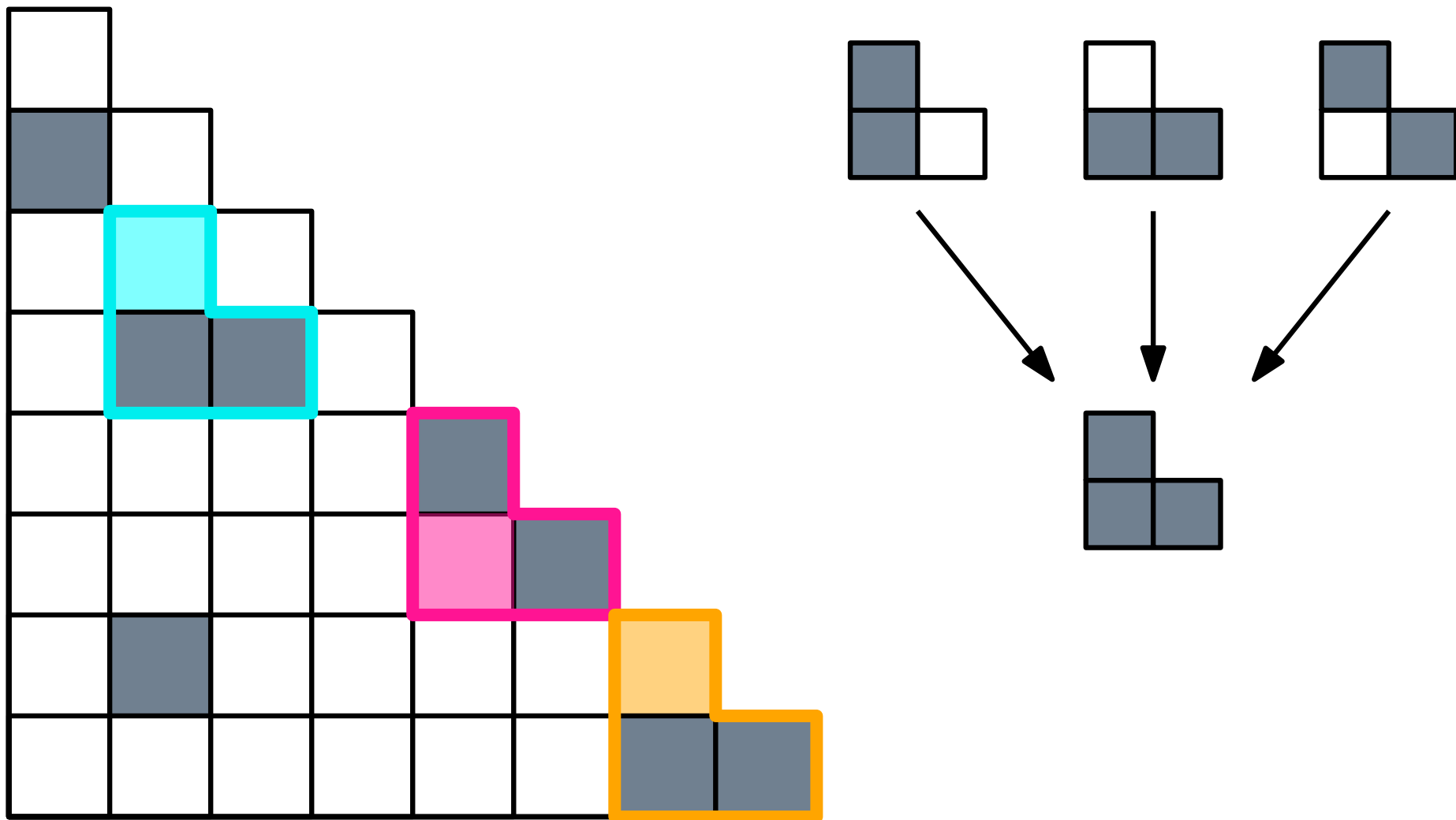
A **filling step** fills the empty cell of a triangle with exactly one empty cell.



Filling configurations

A **configuration** of size n is a set of n cells in the triangle T_n of size n .

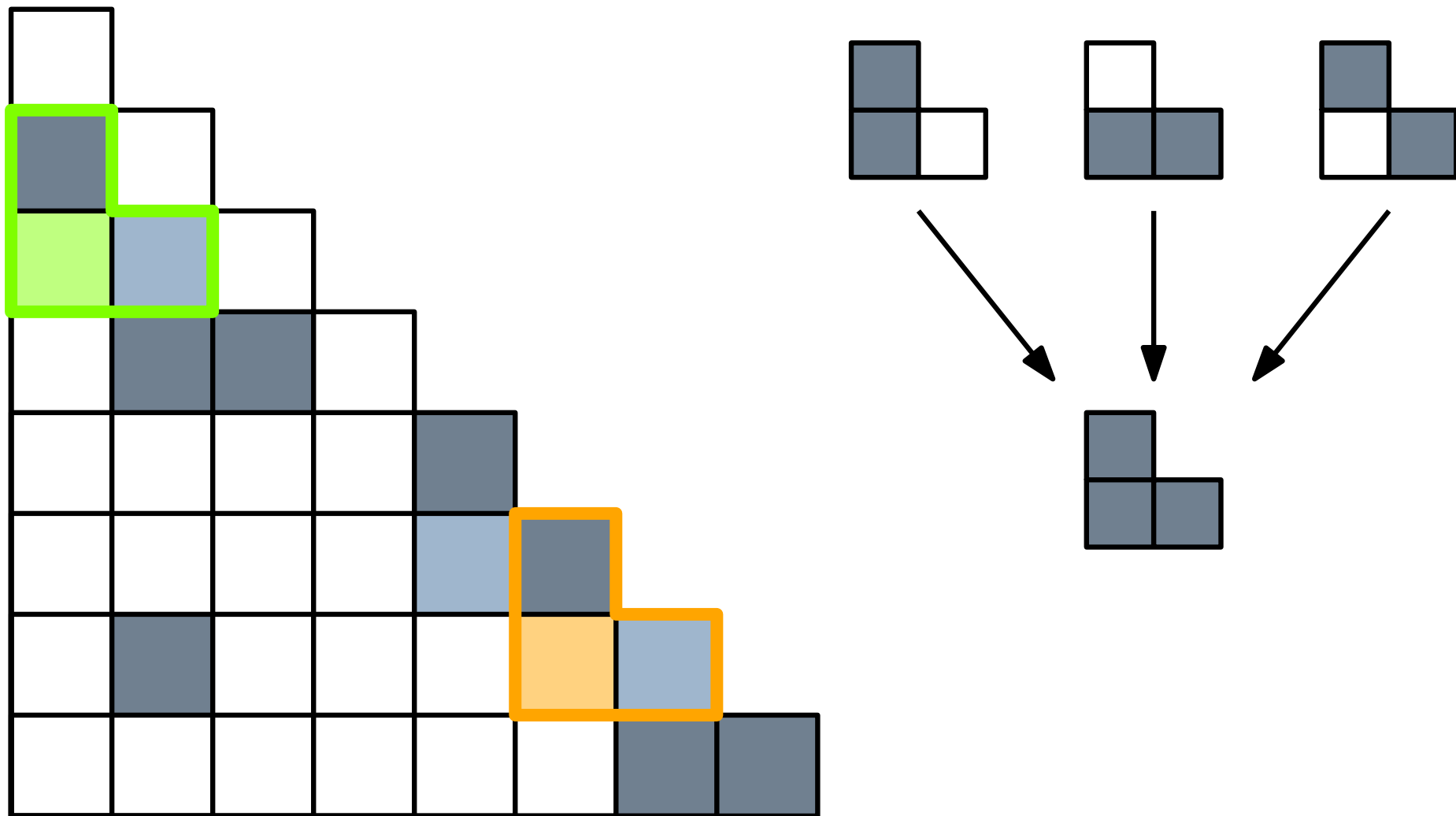
A **filling step** fills the empty cell of a triangle with exactly one empty cell.



Filling configurations

A **configuration** of size n is a set of n cells in the triangle T_n of size n .

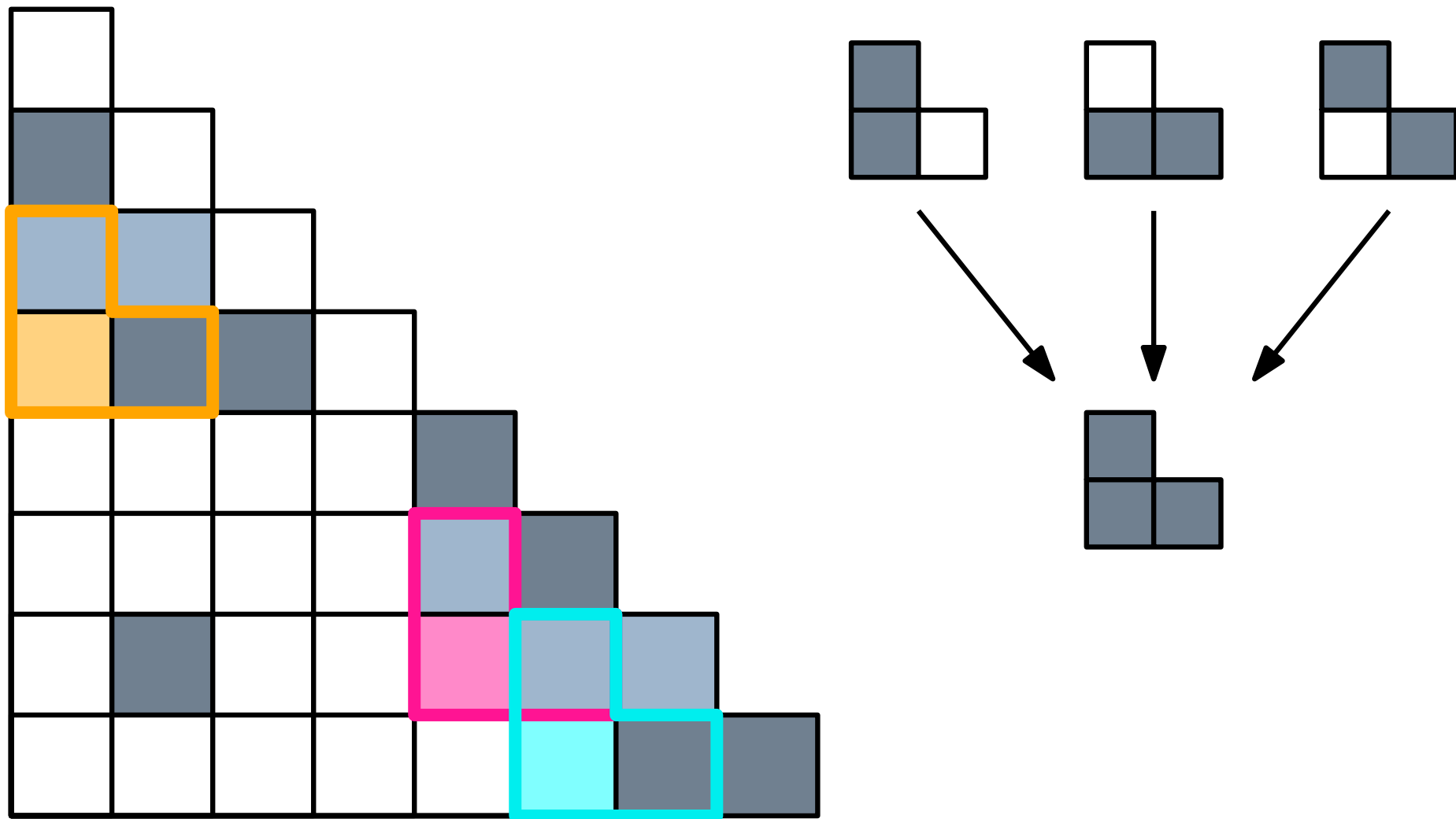
A **filling step** fills the empty cell of a triangle with exactly one empty cell.



Filling configurations

A **configuration** of size n is a set of n cells in the triangle T_n of size n .

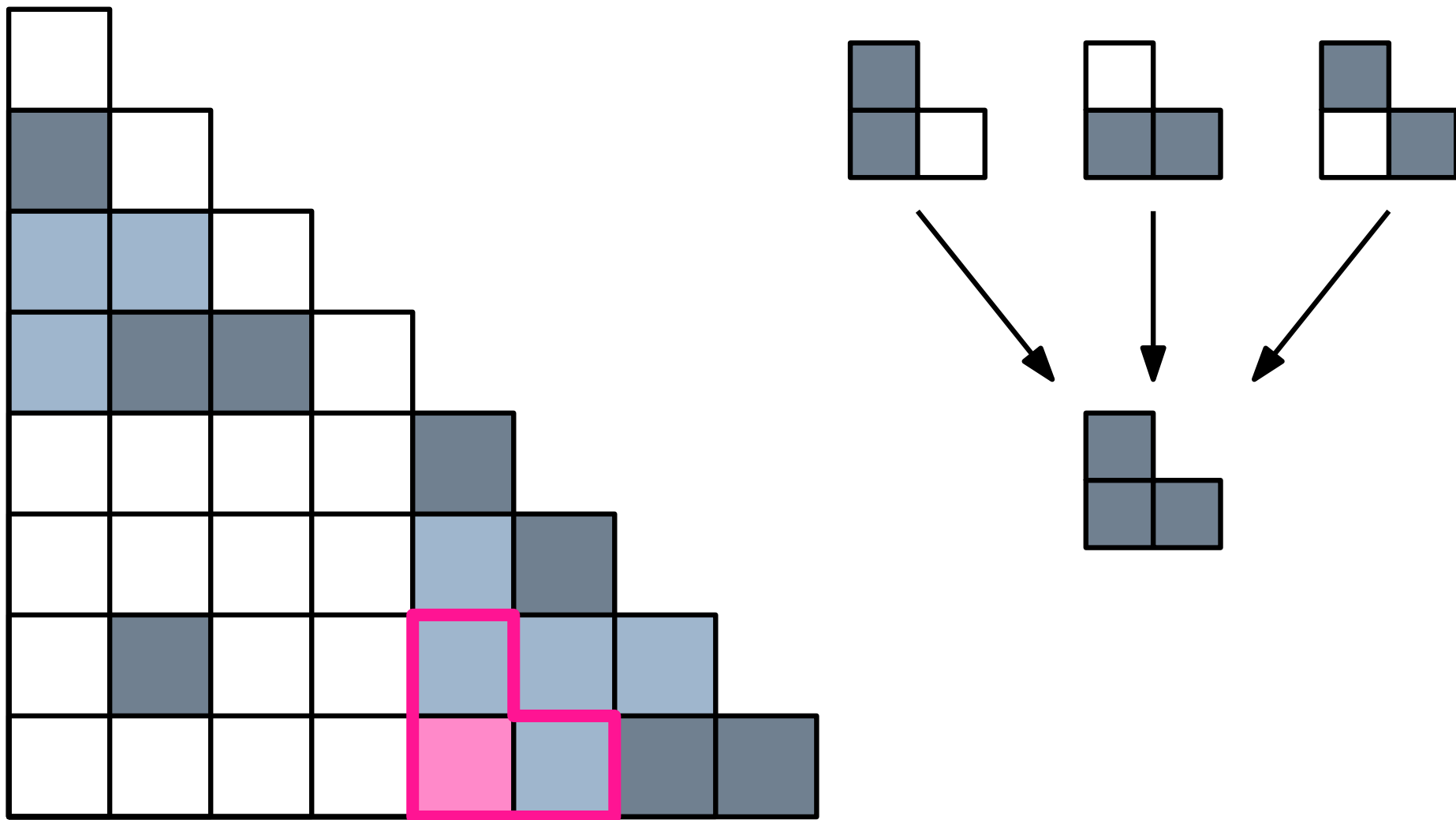
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Filling configurations

A **configuration** of size n is a set of n cells in the triangle T_n of size n .

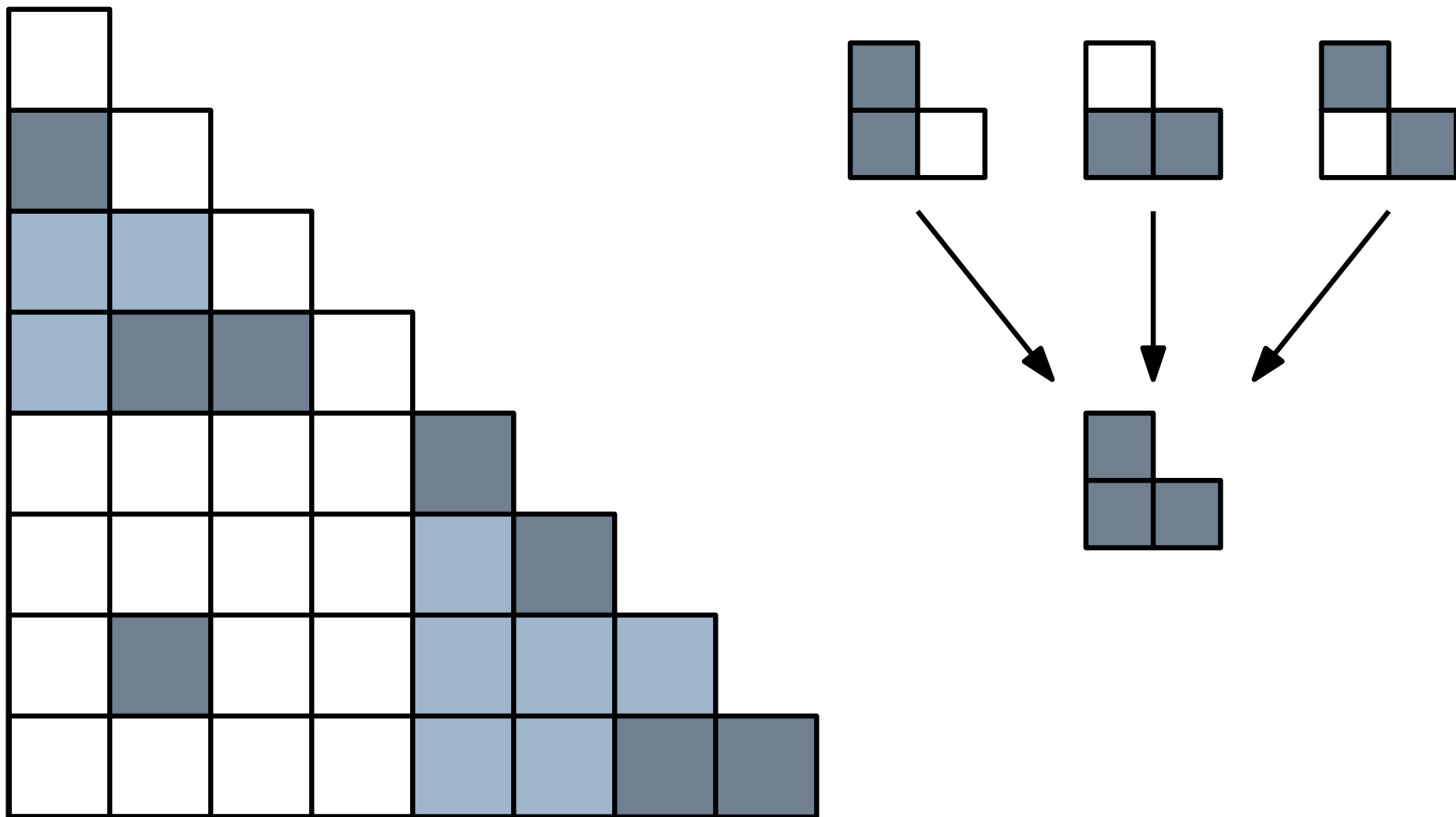
A **filling step** fills the empty cell of a triangle with exactly one empty cell.



Filling configurations

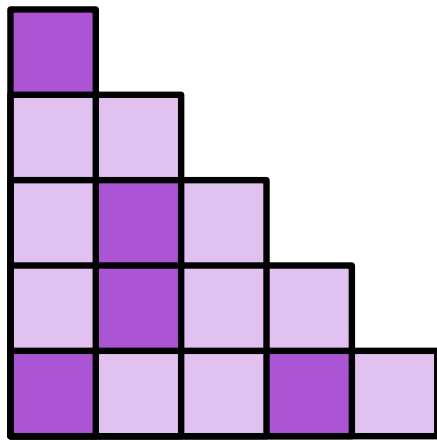
A **configuration** of size n is a set of n cells in the triangle T_n of size n .

A **filling step** fills the empty cell of a triangle with exactly one empty cell.

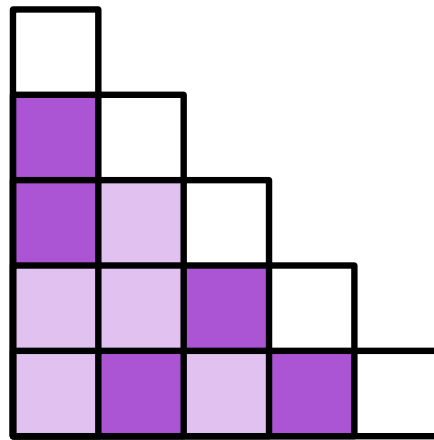


Triangle bases

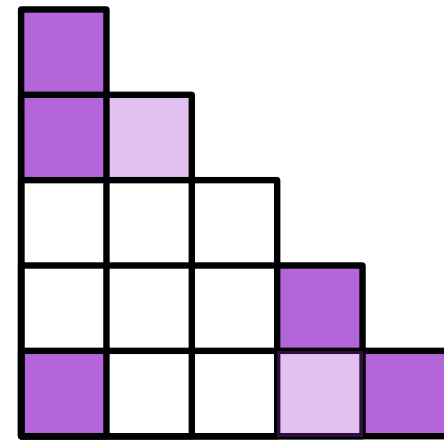
A **triangle basis** of size n is a configuration of n points that fills T_n .
Denote \mathcal{B}_n their set.



A basis.

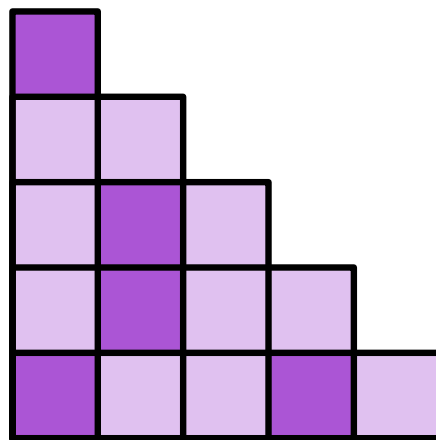


Not a basis.

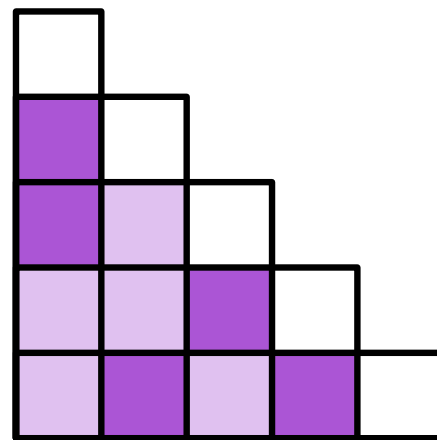


Triangle bases

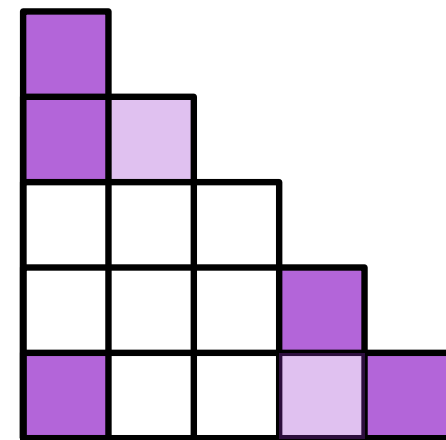
A **triangle basis** of size n is a configuration of n points that fills T_n .
Denote \mathcal{B}_n their set.



A basis.



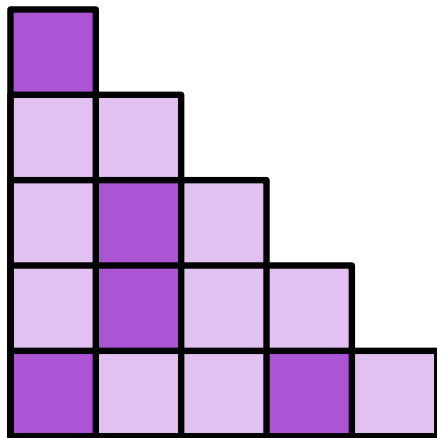
Not a basis.



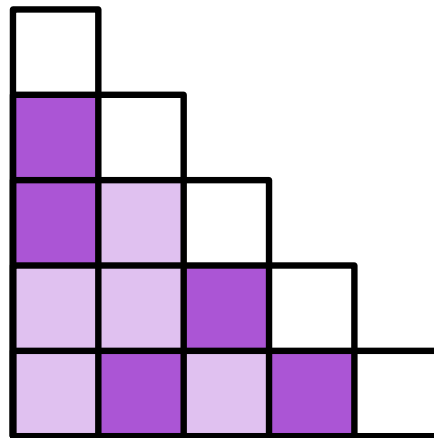
- Used to study “totally extremally permutive” subshifts, a generalization of bipermutive cellular automata [Salo ‘20].

Triangle bases

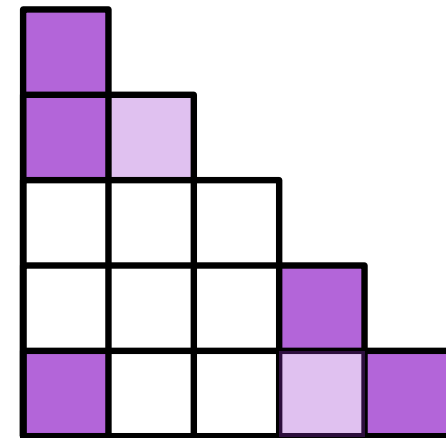
A **triangle basis** of size n is a configuration of n points that fills T_n .
Denote \mathcal{B}_n their set.



A basis.

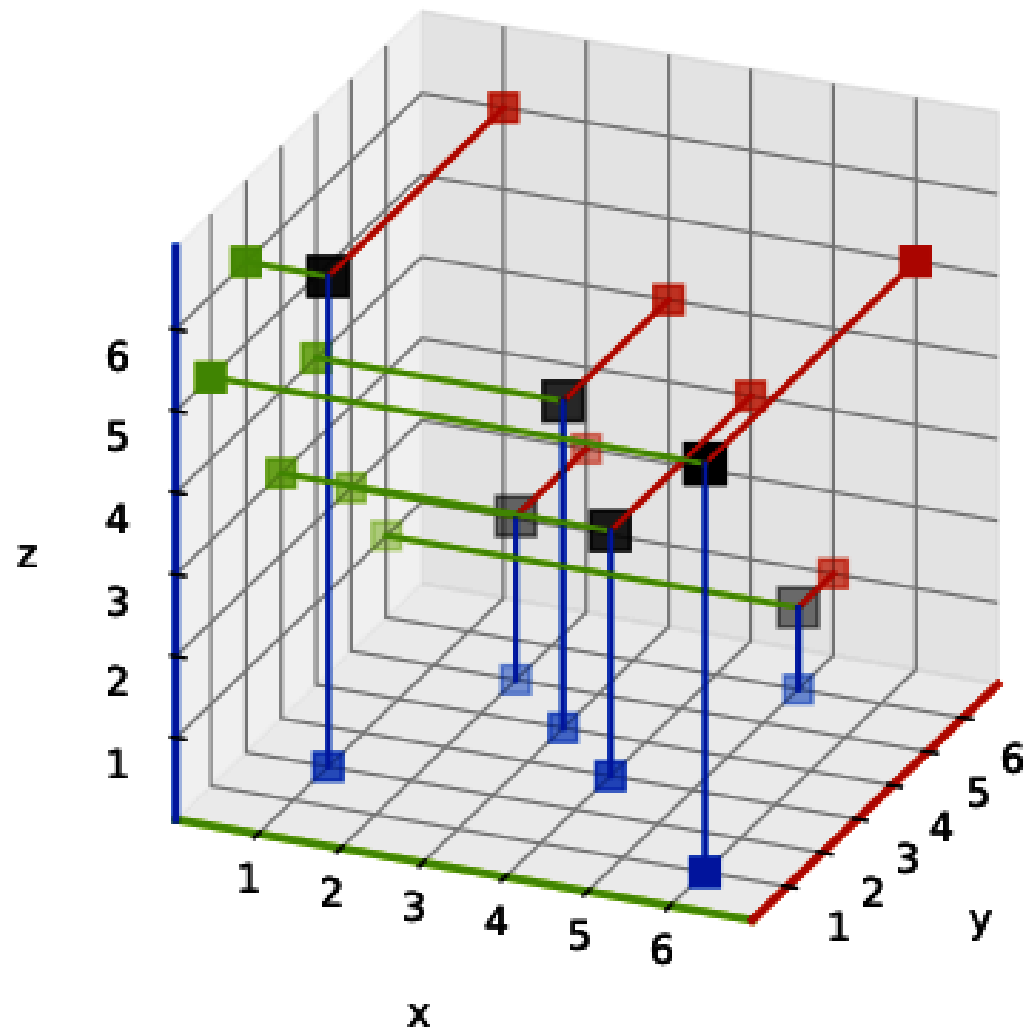


Not a basis.

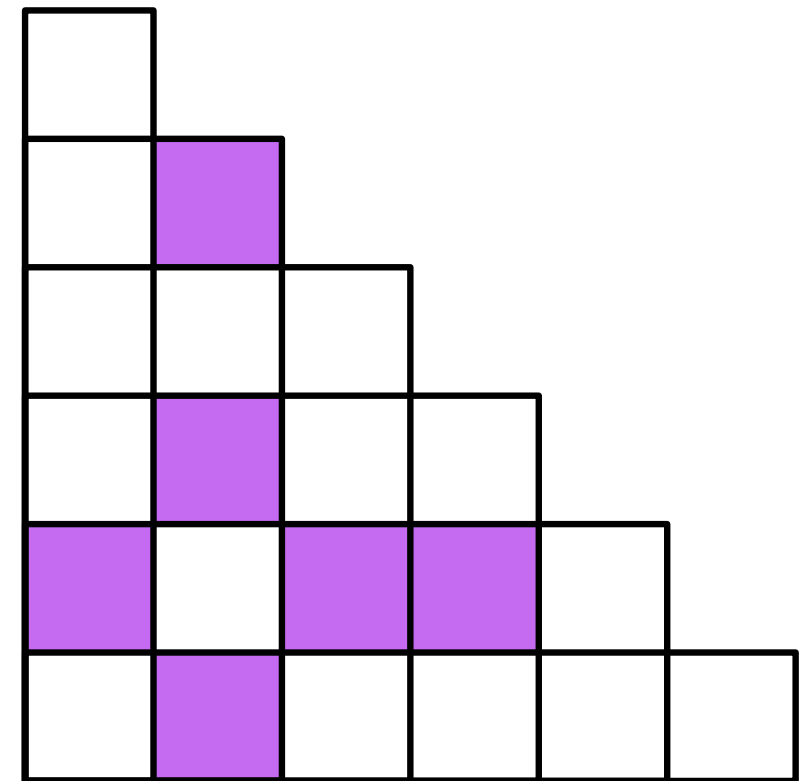


Theorem. [S. '25] For all n , the set of triangle bases of size n is in bijection with $Av_n((\textcolor{blue}{12}, \textcolor{red}{12}), (\textcolor{blue}{312}, \textcolor{red}{231}))$.

II- A bijection

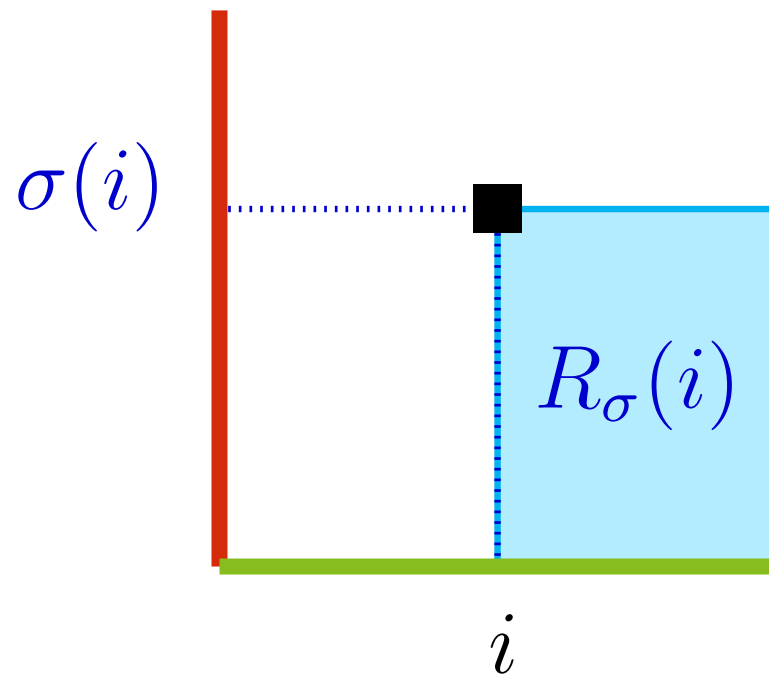


Γ



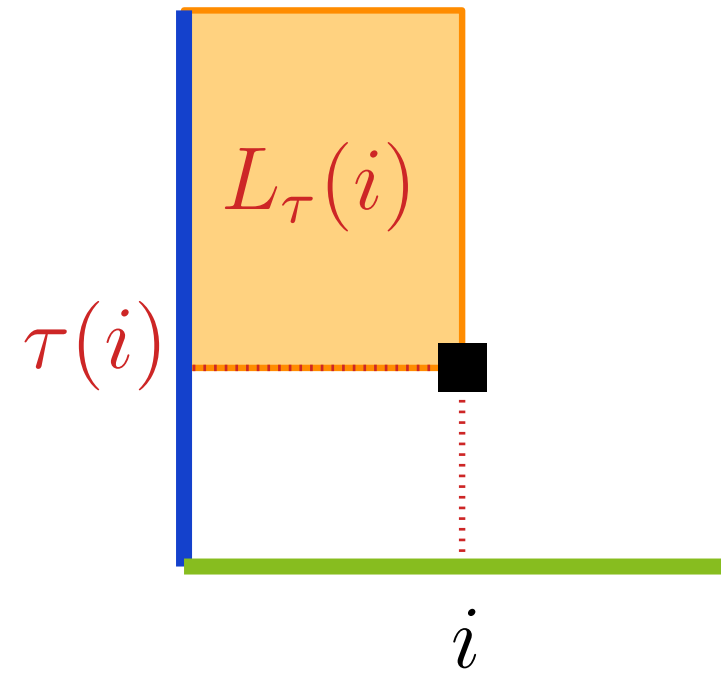
The bijection

An **inversion** of $\sigma \in \mathfrak{S}_n$ is $(i, j) \in \llbracket 1, n \rrbracket$ with $i < j$ and $\sigma(i) > \sigma(j)$.



Right inversion set at i

$$r_\sigma(i) = |R_\sigma(i)|$$



Left inversion set at i

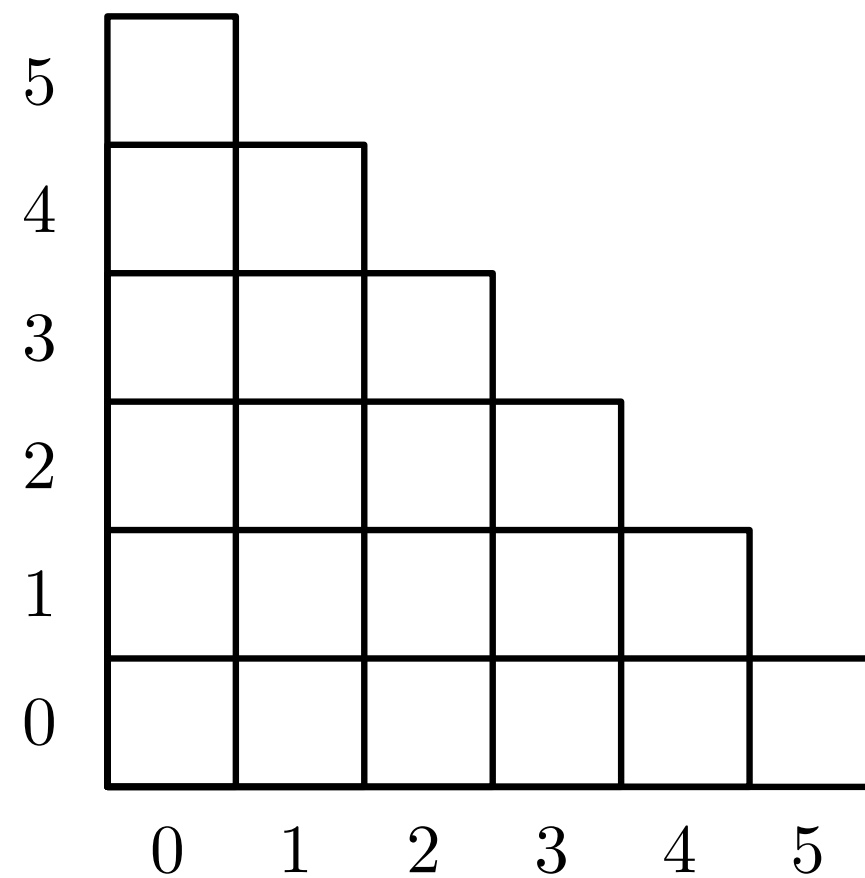
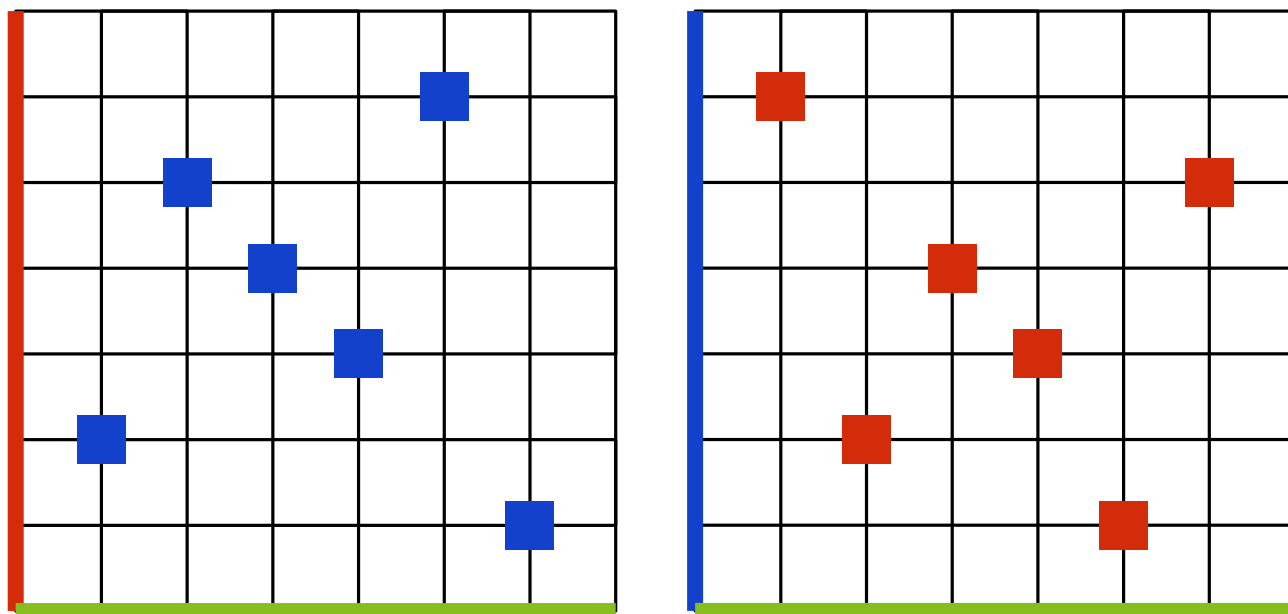
$$\ell_\tau(i) = |L_\tau(i)|$$

The bijection from 3-permutations to bases:

$$\Gamma : (\sigma, \tau) \mapsto \{(r_\sigma(i), \ell_\tau(i)) \mid i \in \llbracket 1, n \rrbracket\}$$

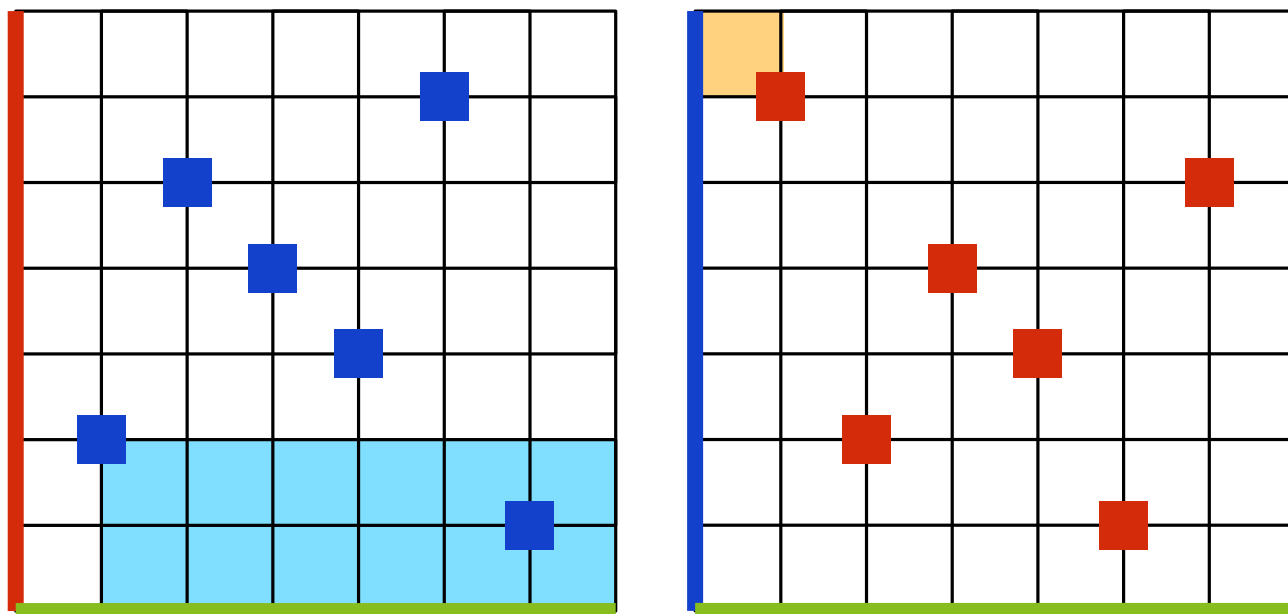
The bijection

$$\Gamma : (\sigma, \tau) \mapsto \{(r_\sigma(i), \ell_\tau(i)) \mid i \in \llbracket 1, n \rrbracket\}$$

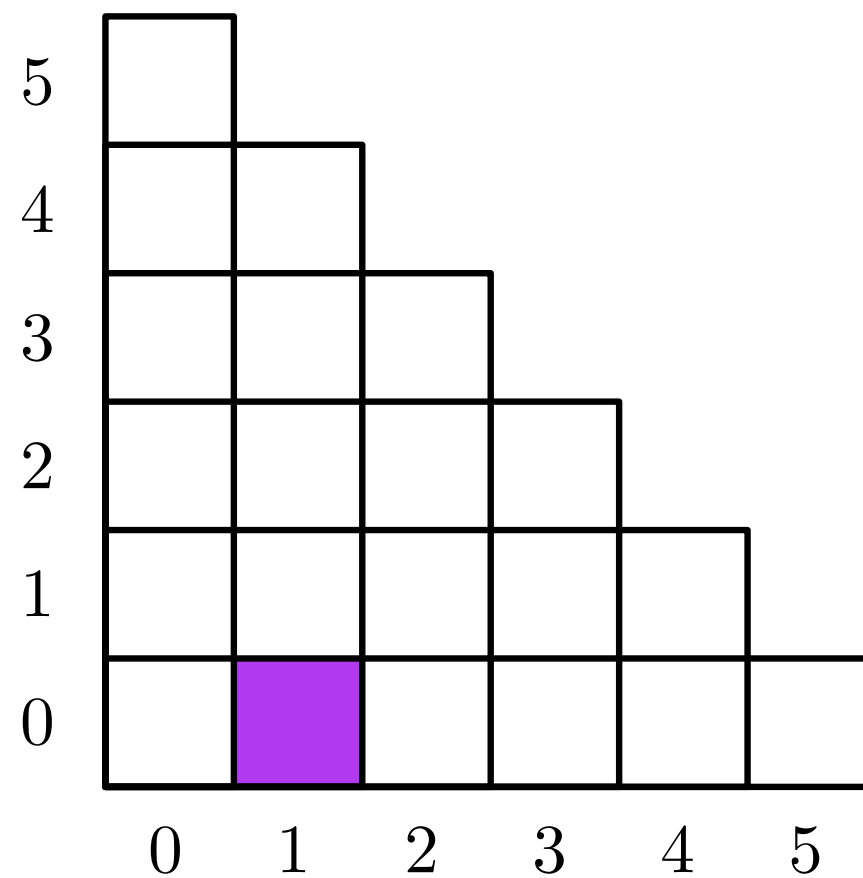


The bijection

$$\Gamma : (\sigma, \tau) \mapsto \{(r_\sigma(i), \ell_\tau(i)) \mid i \in \llbracket 1, n \rrbracket\}$$

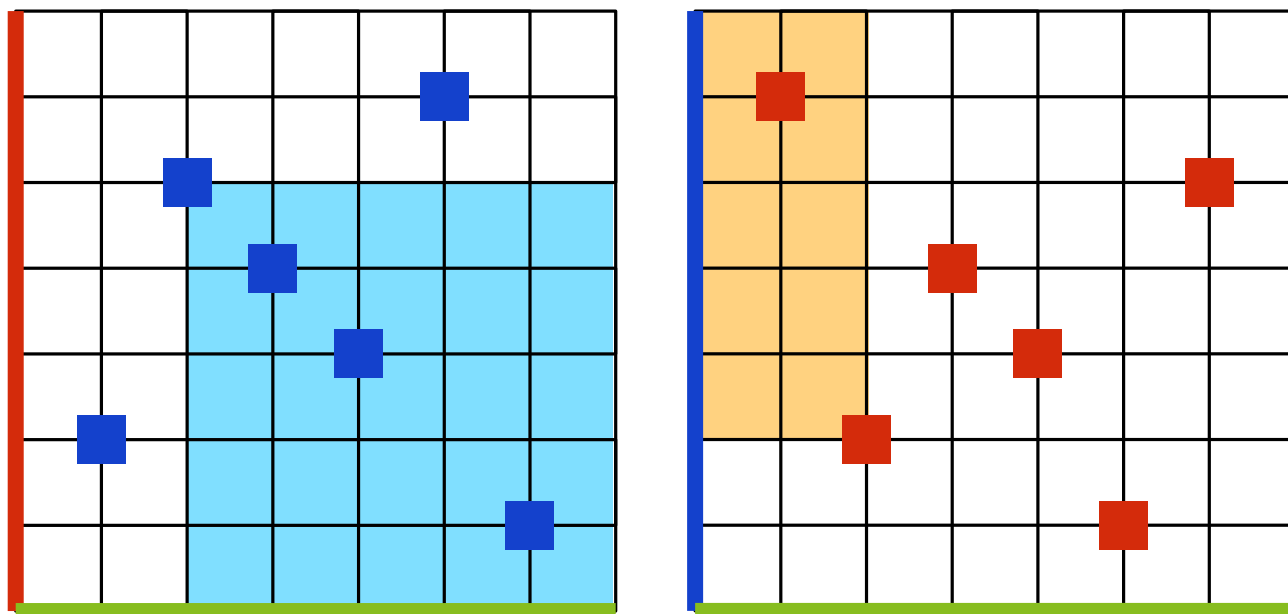


$$i = 1 \mapsto (1, 0)$$

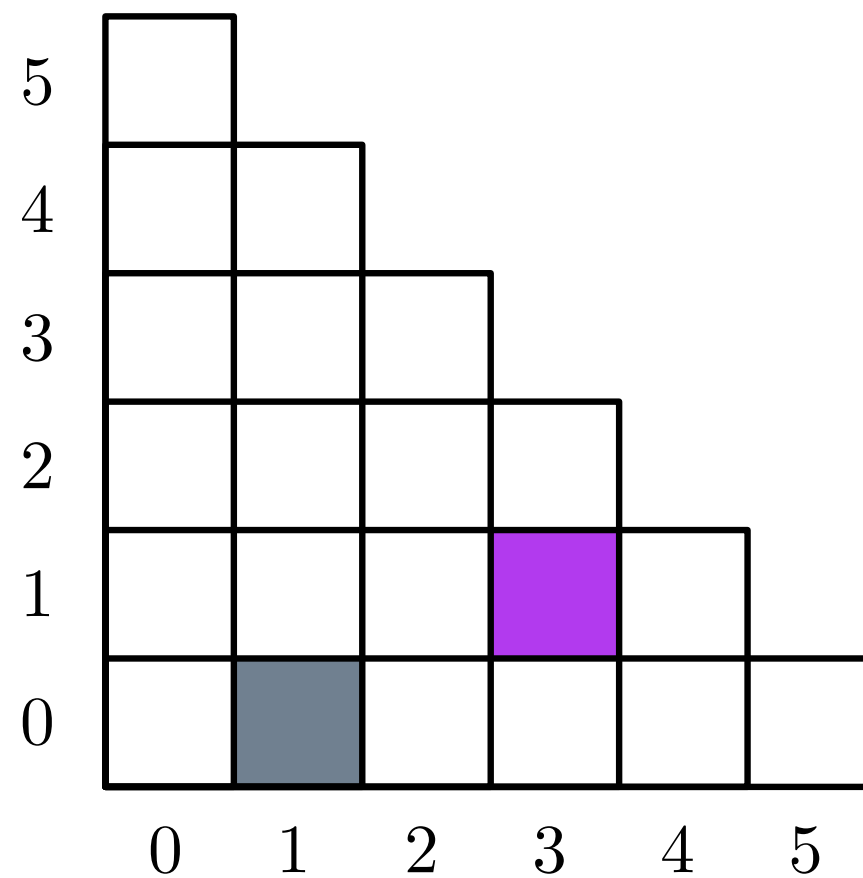


The bijection

$$\Gamma : (\sigma, \tau) \mapsto \{(r_\sigma(i), \ell_\tau(i)) \mid i \in \llbracket 1, n \rrbracket\}$$

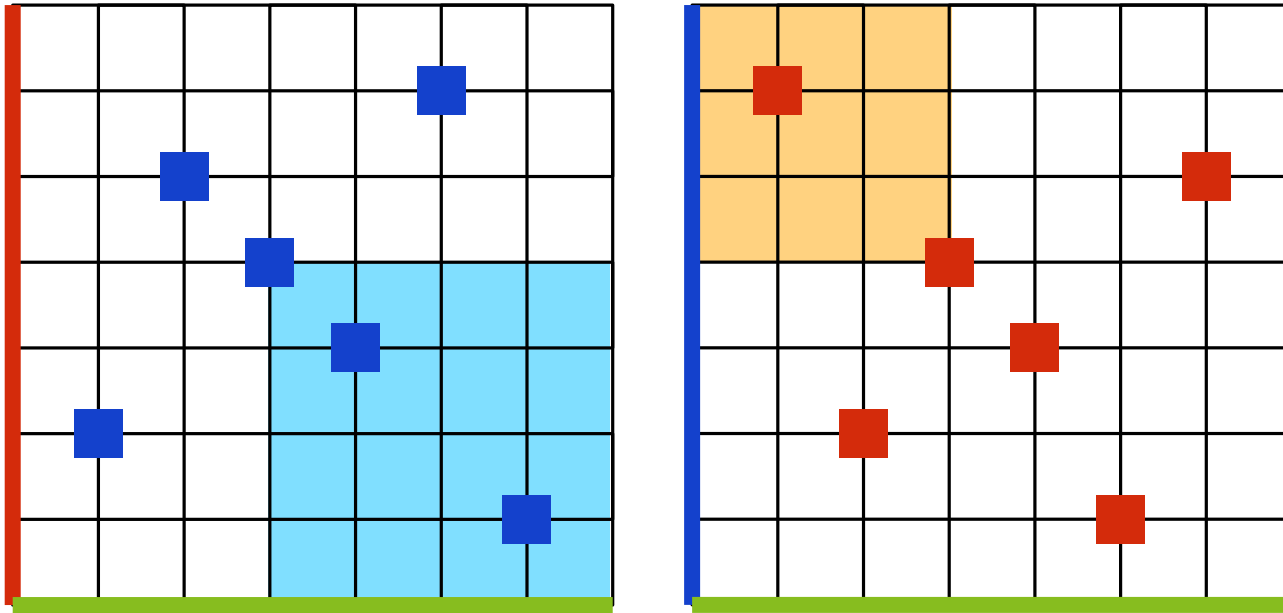


$$i = 2 \mapsto (3, 1)$$

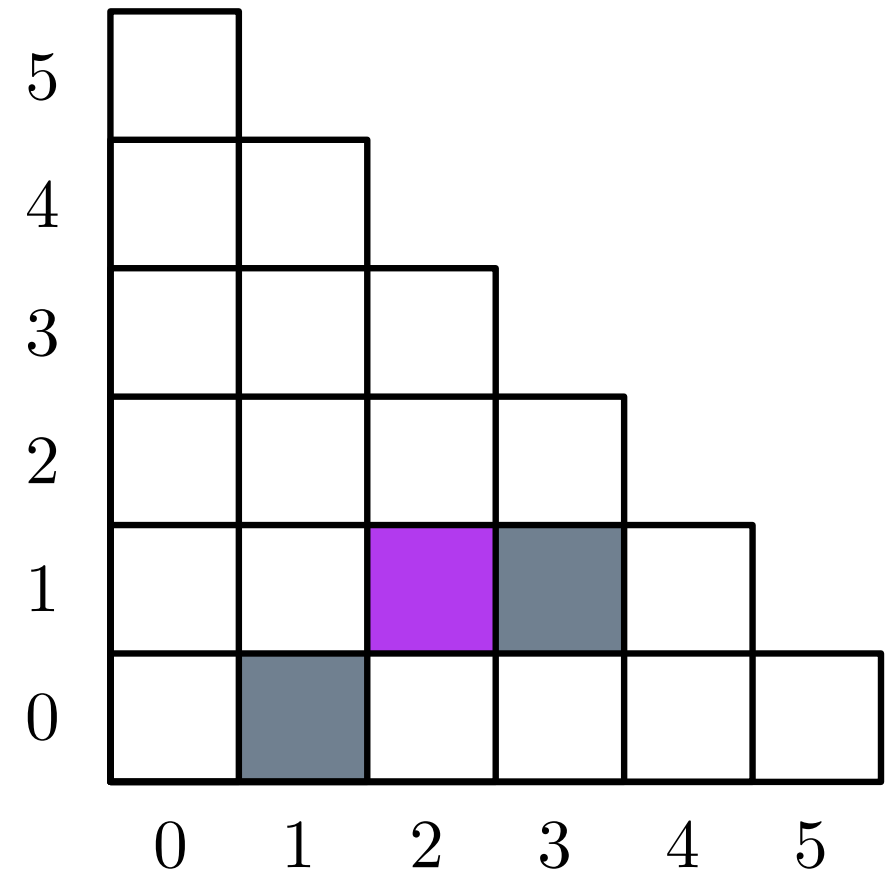


The bijection

$$\Gamma : (\sigma, \tau) \mapsto \{(r_\sigma(i), \ell_\tau(i)) \mid i \in \llbracket 1, n \rrbracket\}$$

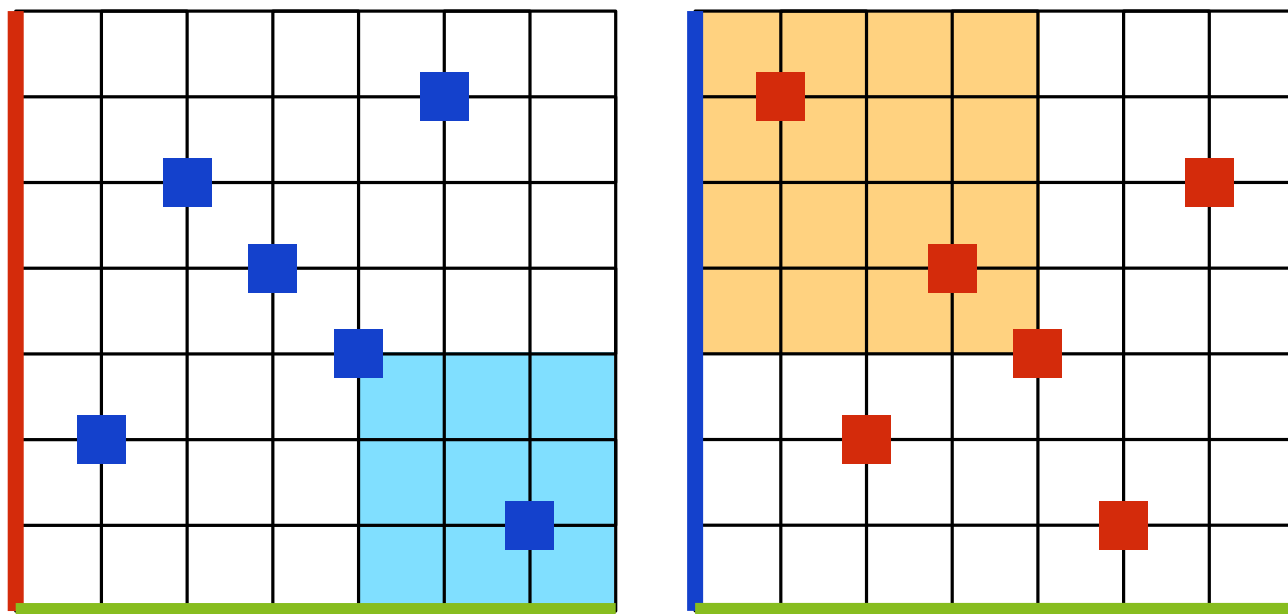


$$i = 3 \mapsto (2, 1)$$

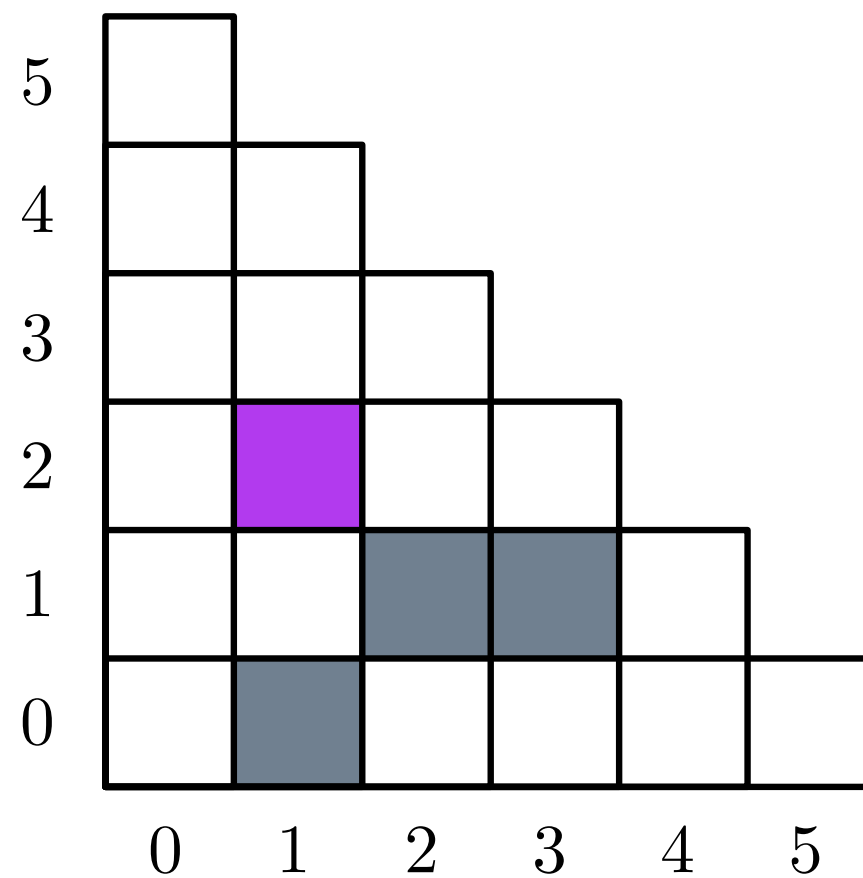


The bijection

$$\Gamma : (\sigma, \tau) \mapsto \{(r_\sigma(i), \ell_\tau(i)) \mid i \in \llbracket 1, n \rrbracket\}$$

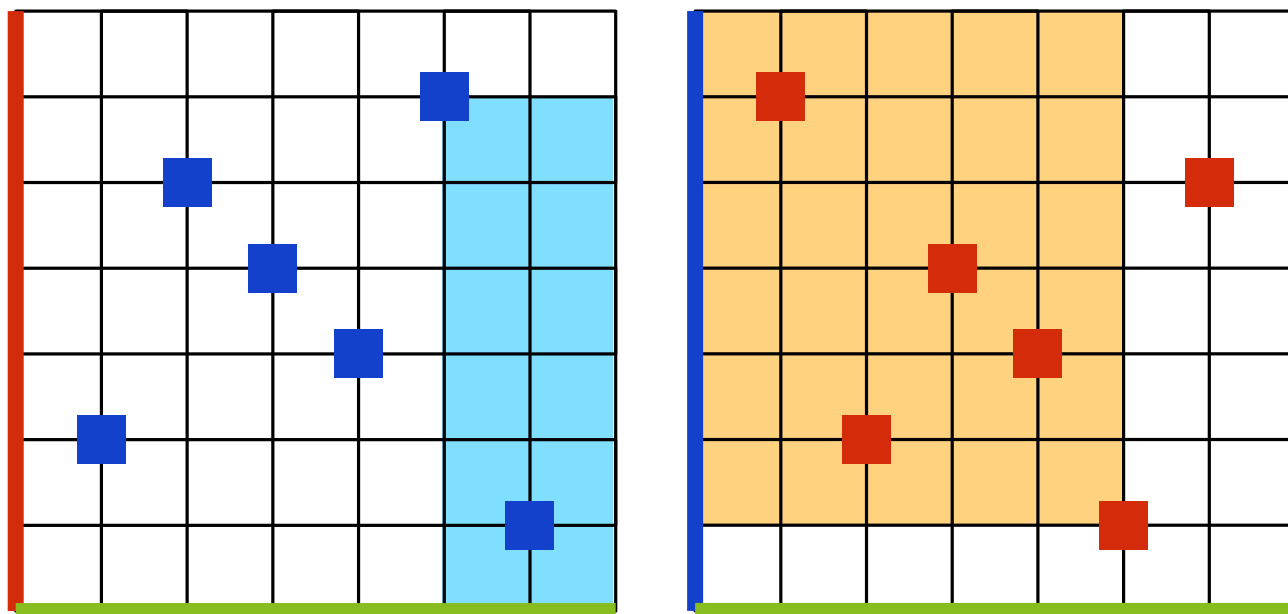


$$i = 4 \mapsto (1, 2)$$

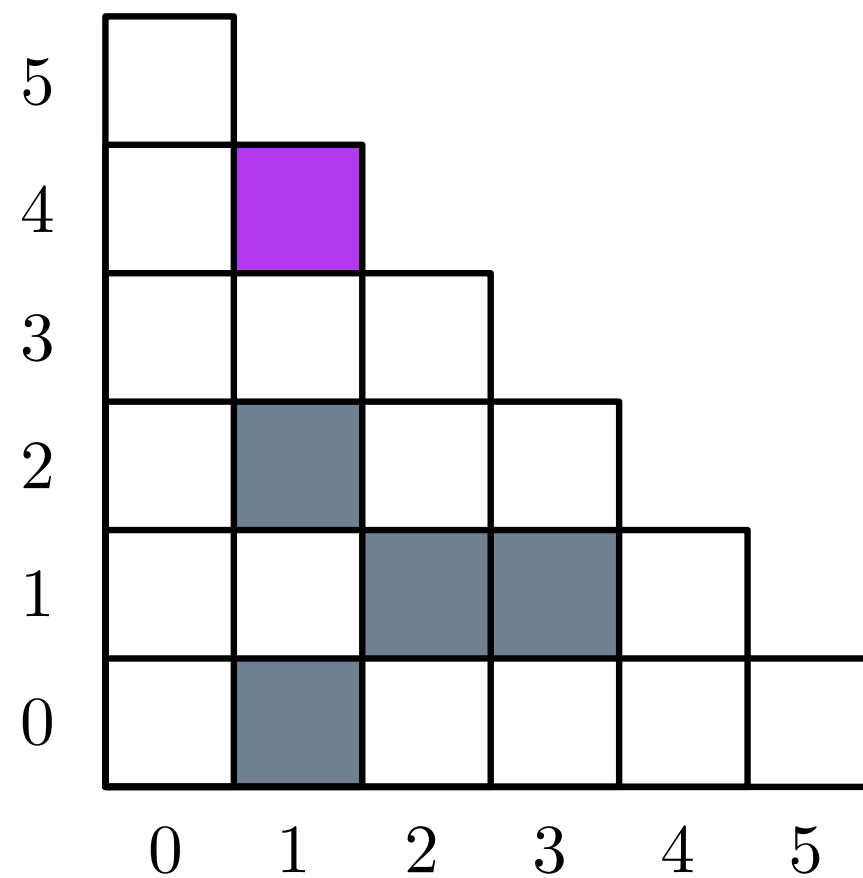


The bijection

$$\Gamma : (\sigma, \tau) \mapsto \{(r_\sigma(i), \ell_\tau(i)) \mid i \in \llbracket 1, n \rrbracket\}$$

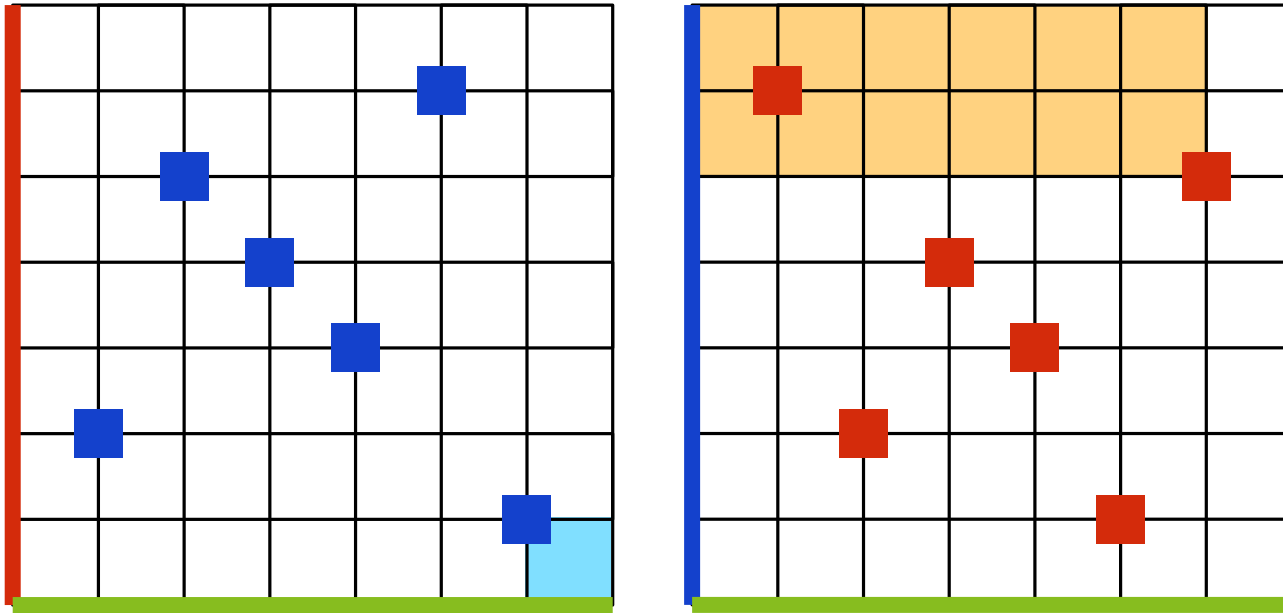


$$i = 5 \mapsto (1, 4)$$

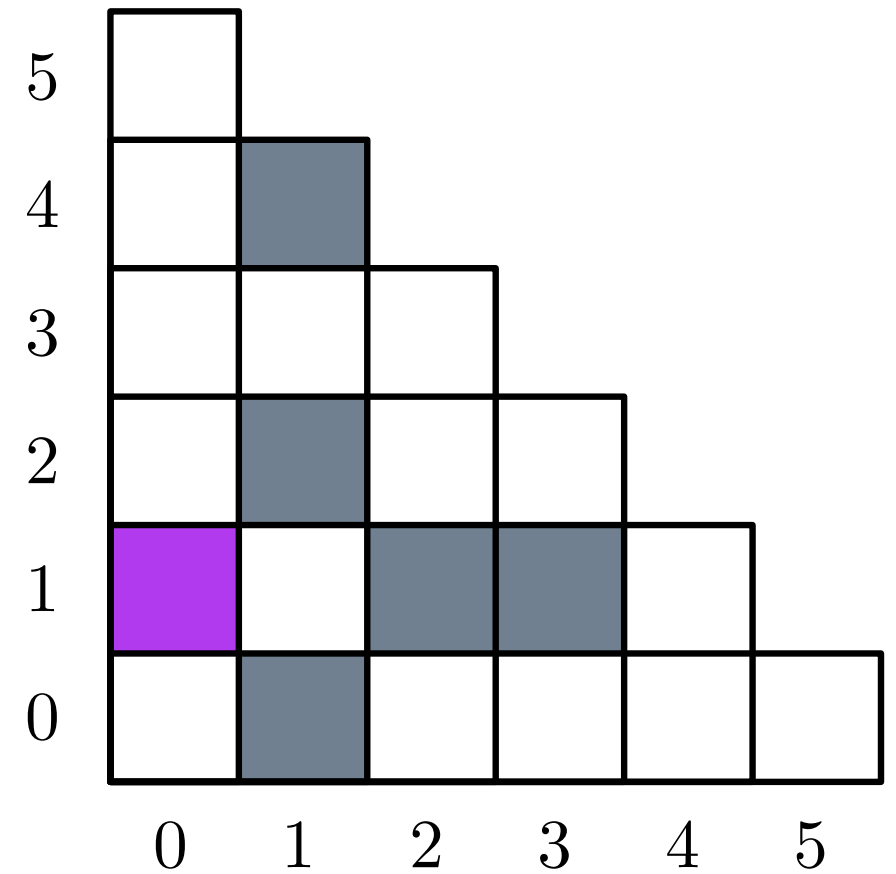


The bijection

$$\Gamma : (\sigma, \tau) \mapsto \{(r_\sigma(i), \ell_\tau(i)) \mid i \in \llbracket 1, n \rrbracket\}$$

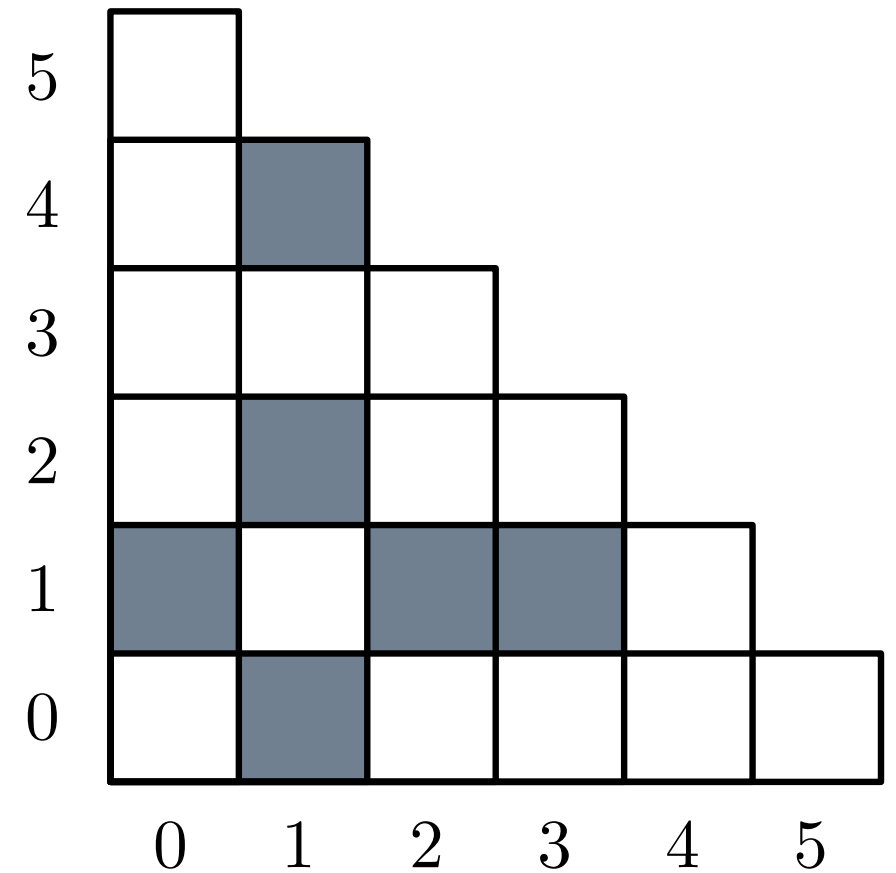
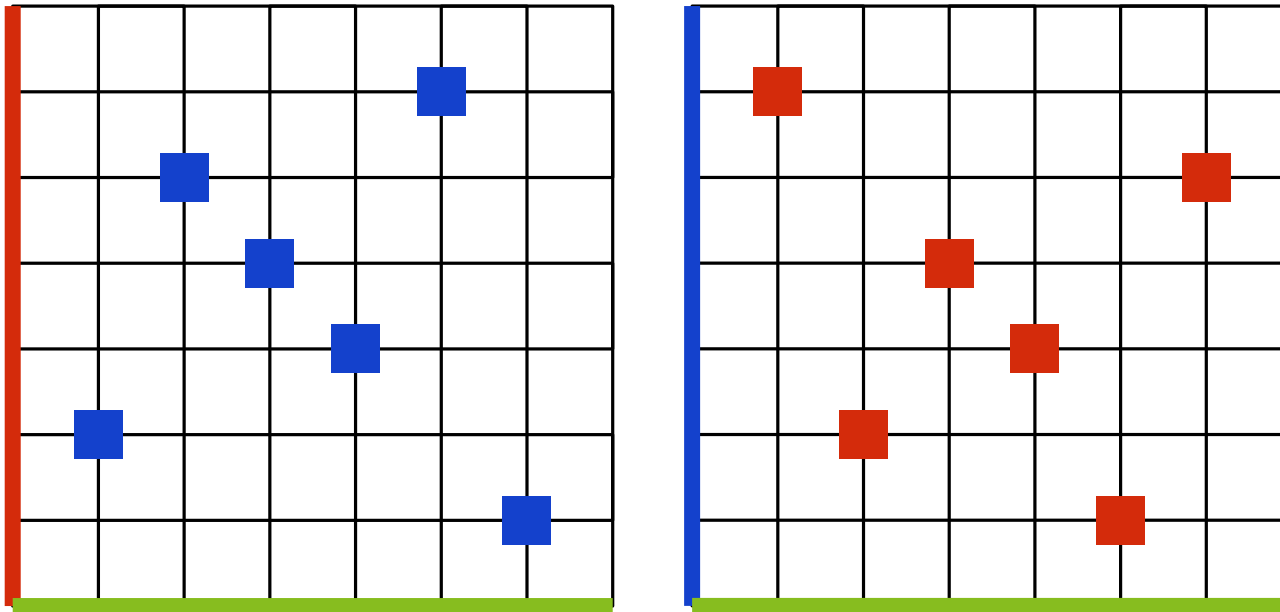


$$i = 6 \mapsto (0, 1)$$



The bijection

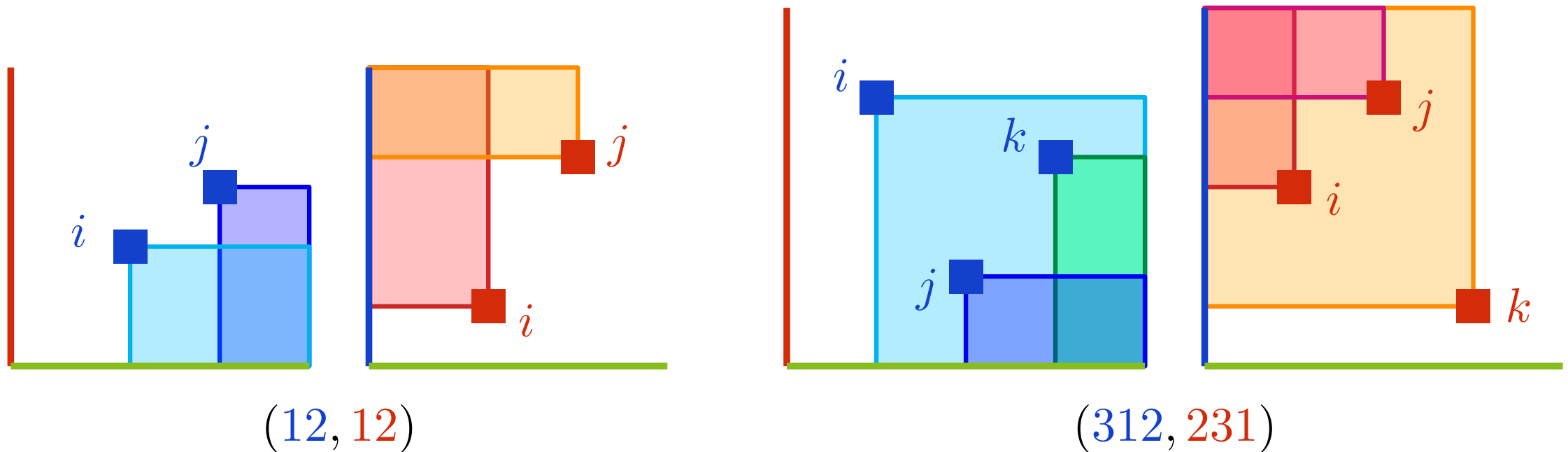
$$\Gamma : (\sigma, \tau) \mapsto \{(r_\sigma(i), \ell_\tau(i)) \mid i \in [1, n]\}$$



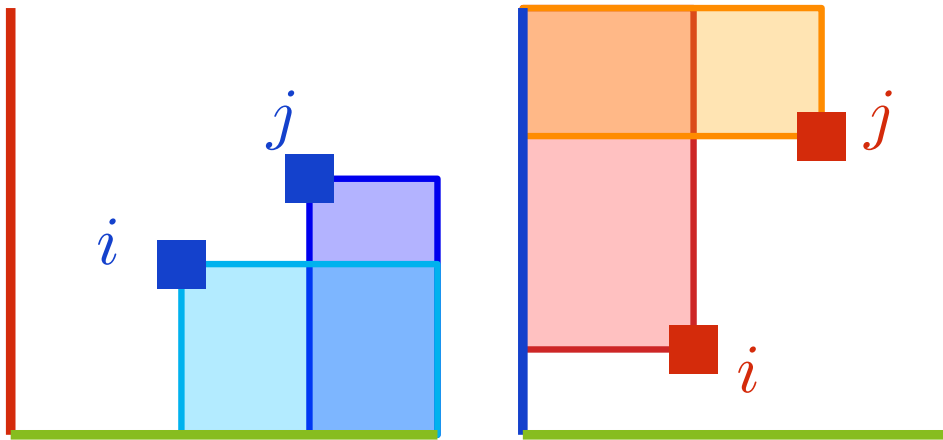
Theorem. [S. '25] For all n , Γ is a bijection between $Av_n((\textcolor{blue}{12}, \textcolor{red}{12}), (\textcolor{blue}{312}, \textcolor{red}{231}))$ and the triangle bases of size n .

Why does avoiding $(12, 12)$ and $(312, 231)$ lead to a triangle basis?

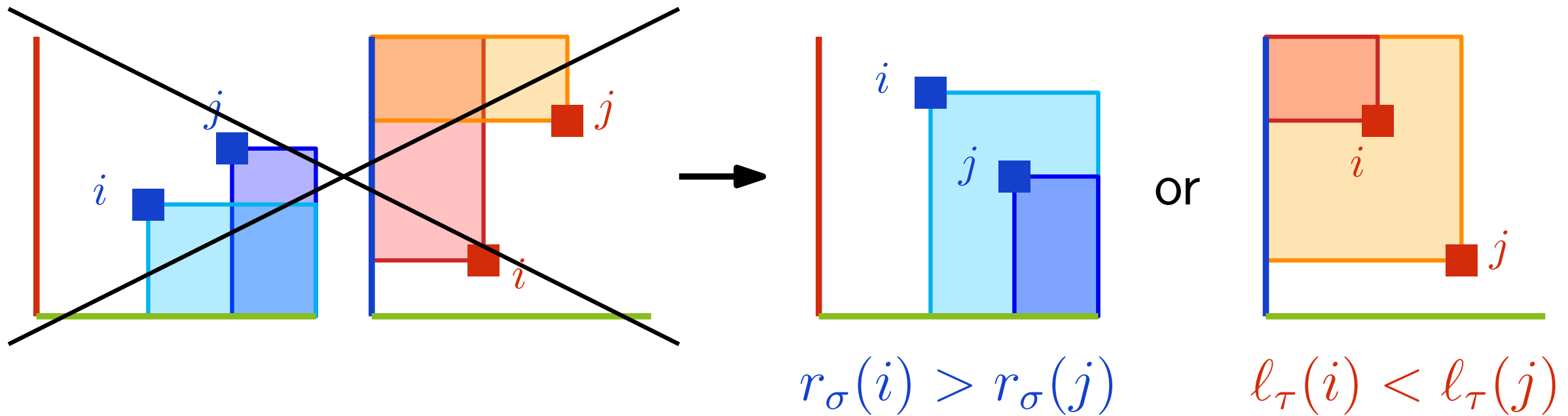
Intuition



Avoiding $(12, 12)$: no “points too close”

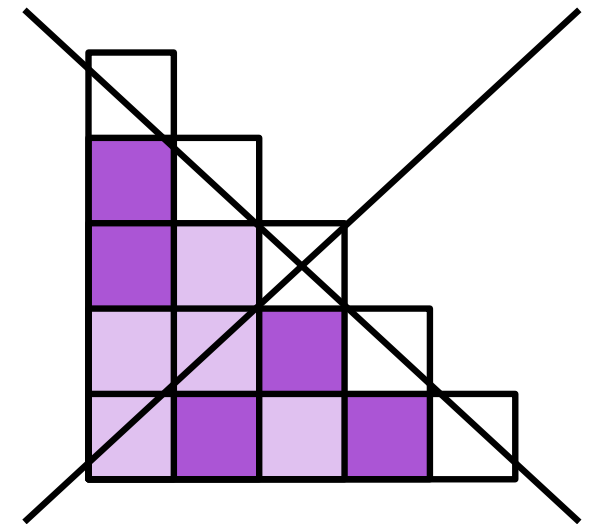


Avoiding $(12, 12)$: no “points too close”

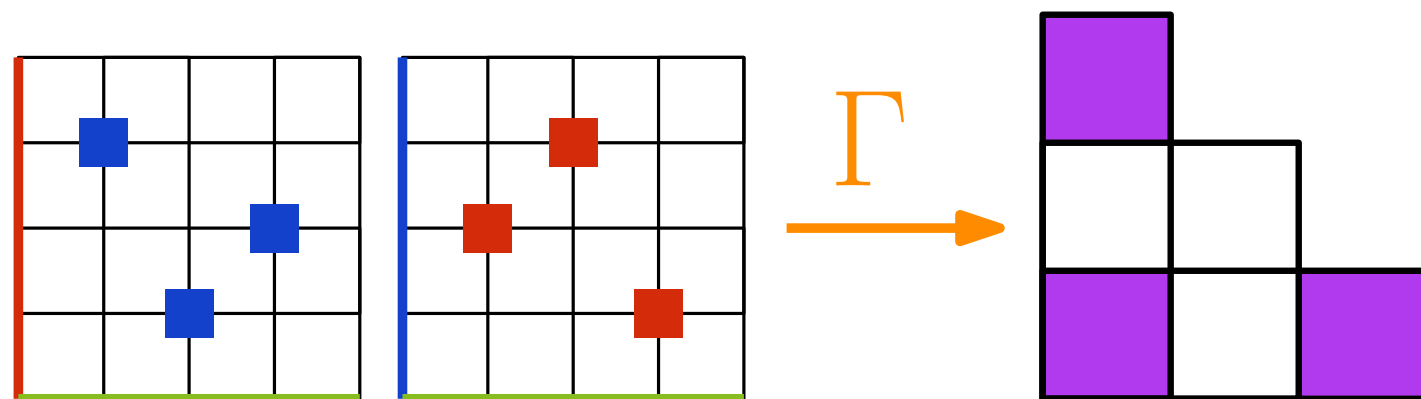


Consequence: If (σ, τ) avoids $(12, 12)$ then

- all points $(r_\sigma(i), l_\tau(i))$ are distinct
- the configuration is **sparse**: there is no triangle T of size k such that $|C \cap T| > k$.



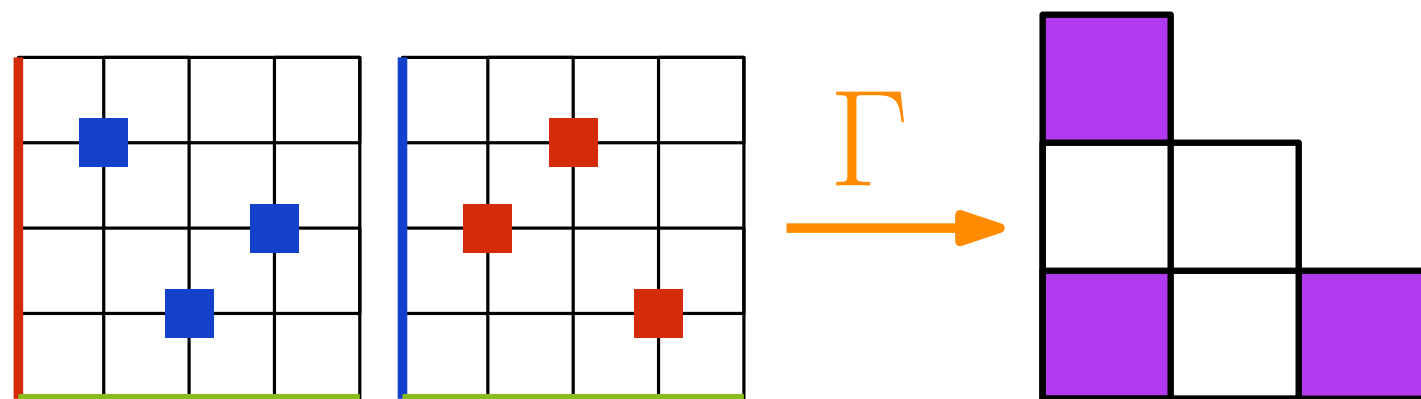
Avoiding $(312, 231)$: no “points too far”



the only sparse
configuration of size 3
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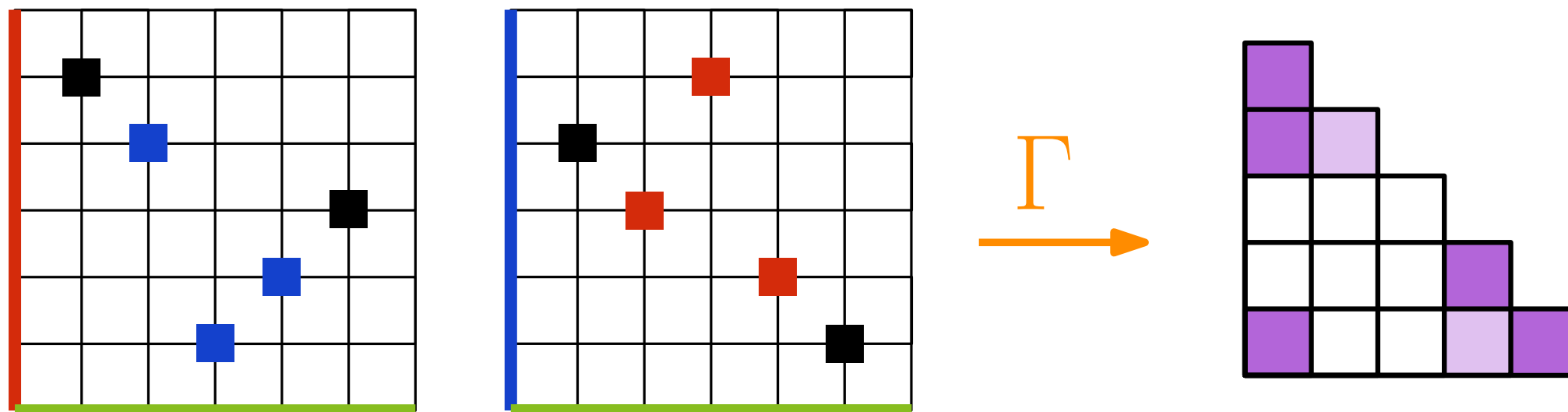
Intuition: Avoiding $(312, 231)$ prevents “gaps”.

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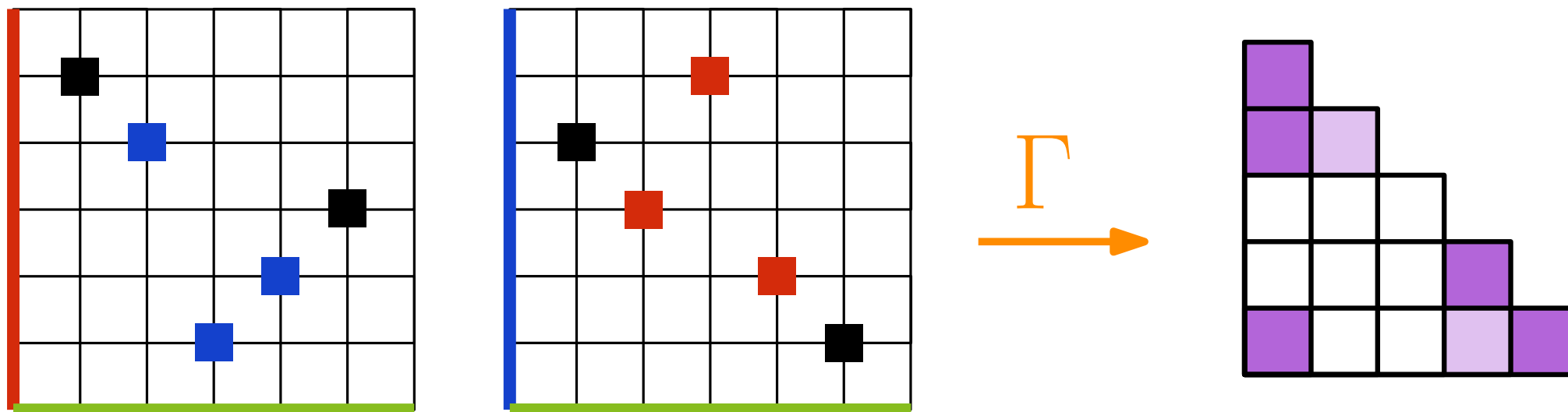
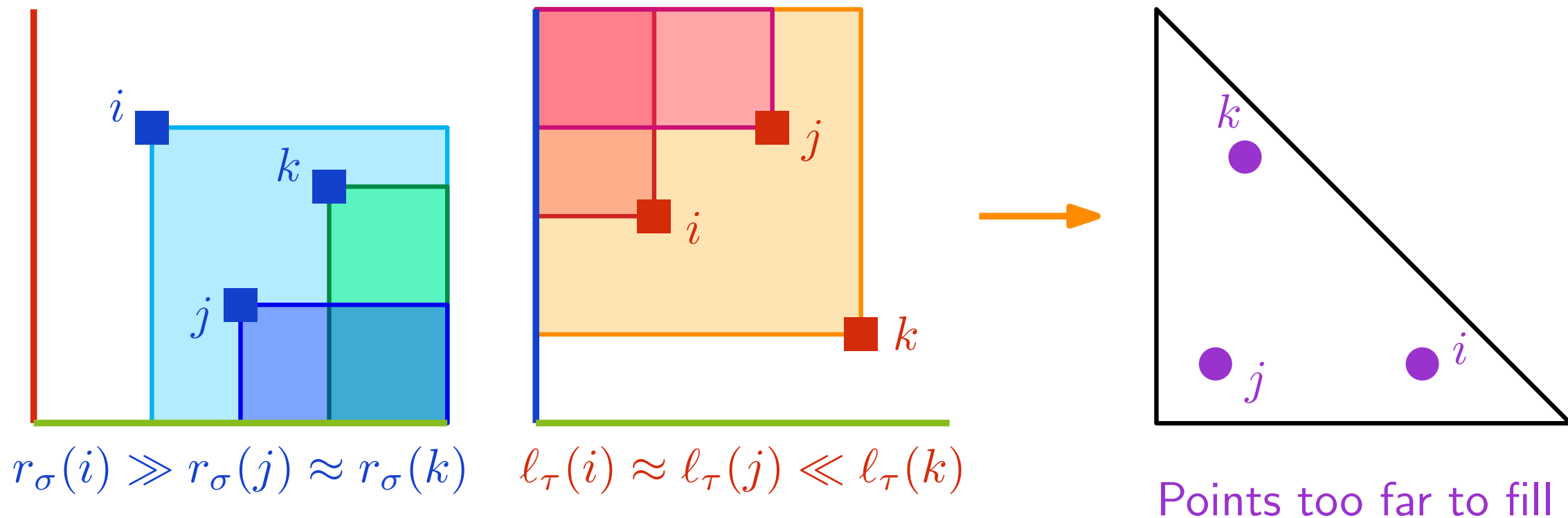


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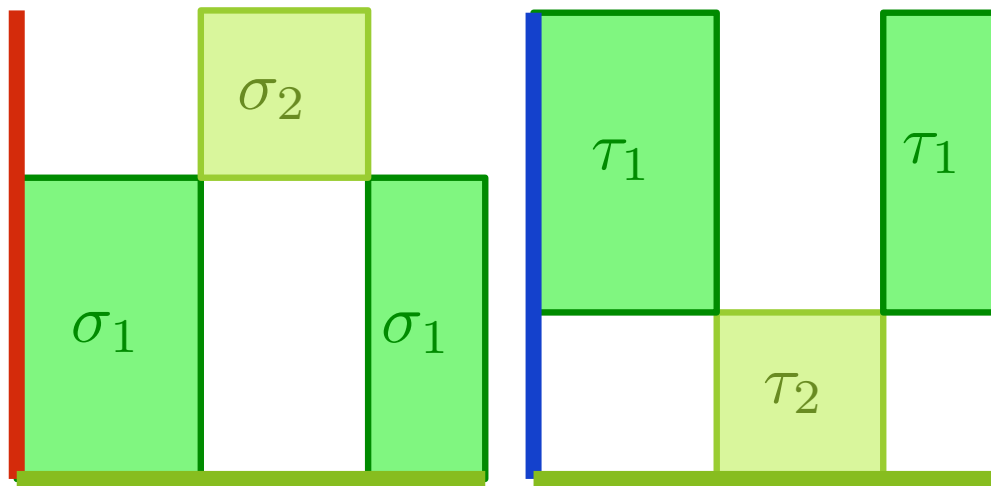


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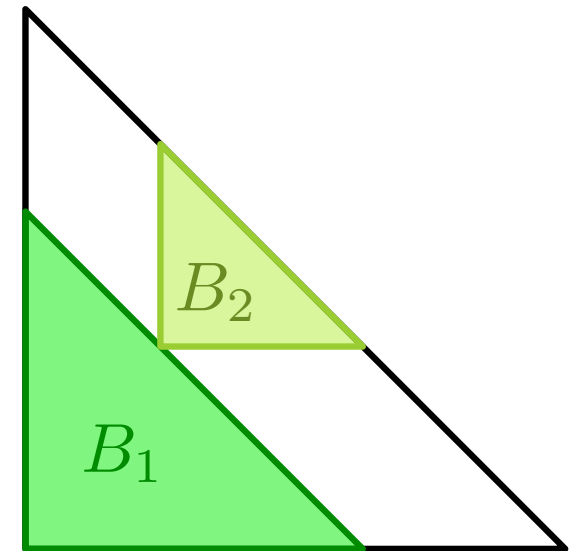


The key tool of the proof:

Isomorphic recursive decompositions



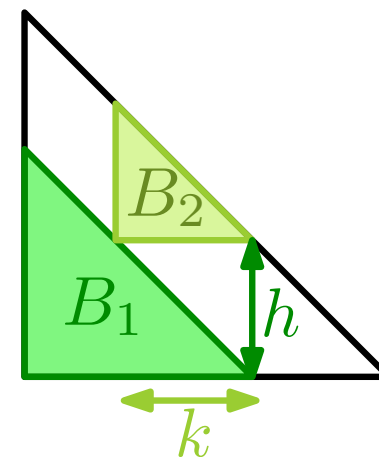
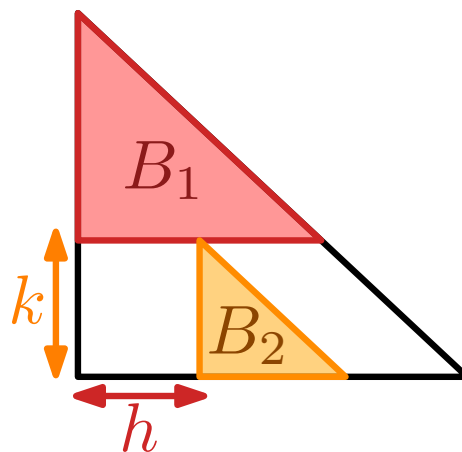
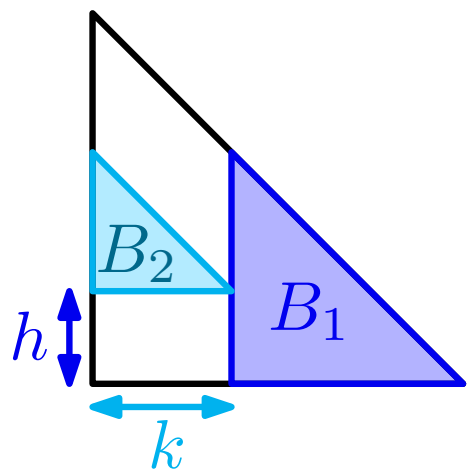
3-permutations



Bases

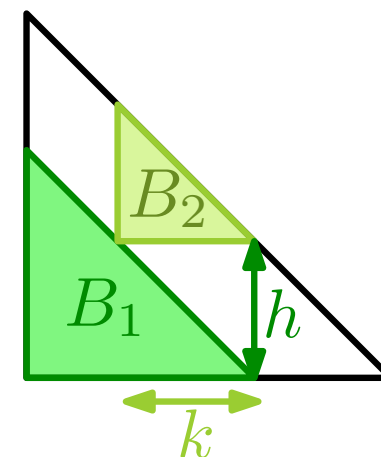
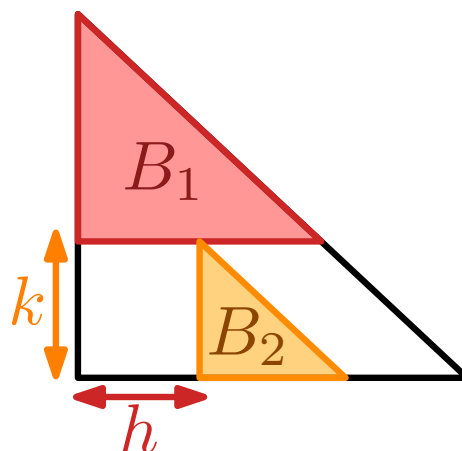
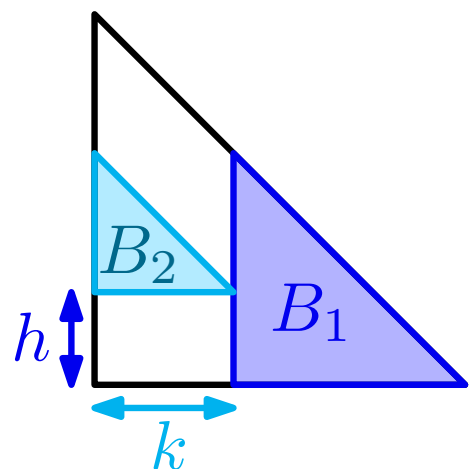
Isomorphic recursive decompositions

Lemma. [Salo, S. '22] Any basis of size $n \geq 2$ can be cut into two smaller bases in one of the 3 following ways.

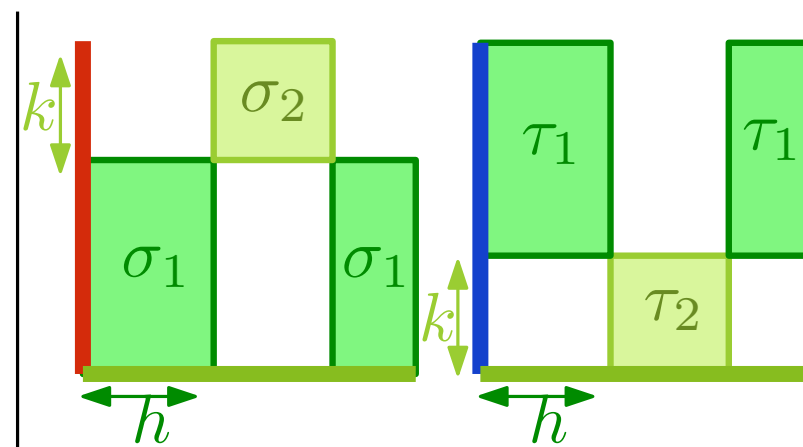
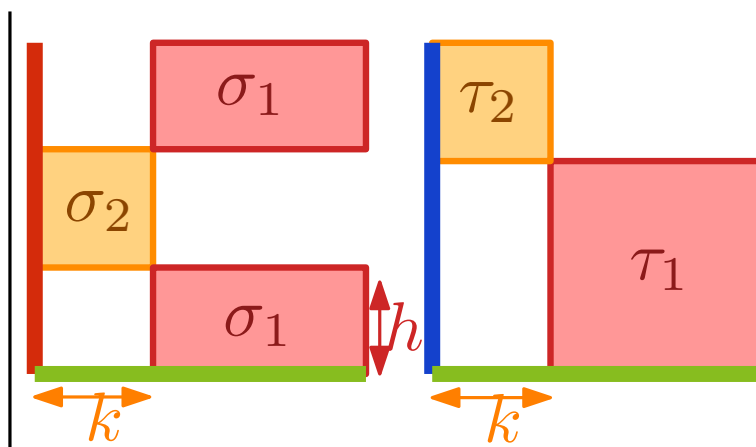
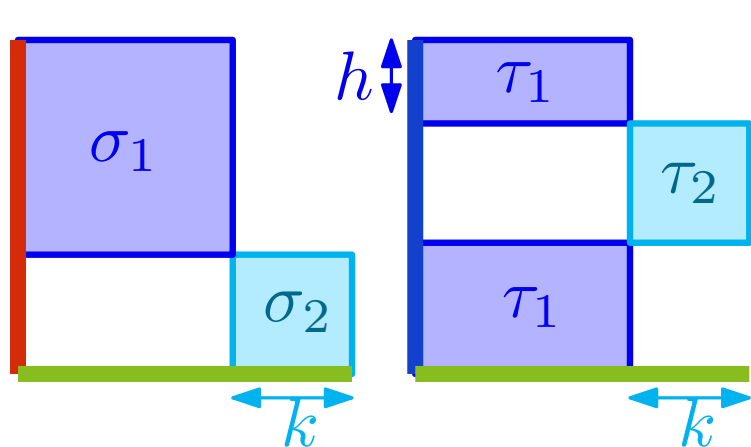


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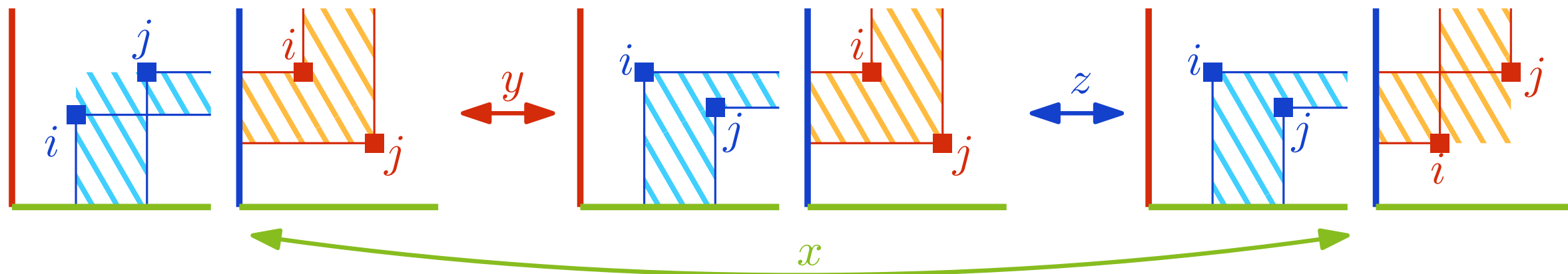
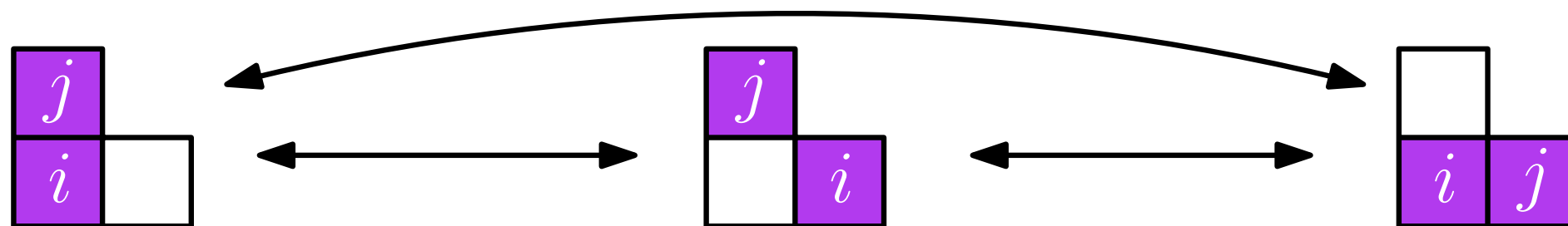
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- ▶ We can now prove everything by induction!
- $\Gamma(Av_n((12, 12), (312, 231))) \subset \mathcal{B}_n$,
- Γ is surjective,
- Γ is injective.

Nice properties and consequences

- Simple construction that transports symmetries.
- Links two objects that are understood very differently \implies tools transfer.
 - ▶ On bases: a canonical labelling on bases, maybe a characterisation by forbidden patterns.
 - ▶ On permutations: a dynamical system on 3-permutations (and others!) which could allow sampling.



What's next?

- No enumerative result.

► Best known bounds : $3n! \leq |\mathcal{B}_n| \leq c \left(\frac{e}{2}\right)^n n^{n-\frac{5}{2}}$ with $c > 0$.

$|Av_n((\textcolor{blue}{12}, \textcolor{red}{12}), (\textcolor{blue}{312}, \textcolor{red}{231}))| \leq |Av_n(\textcolor{blue}{12}, \textcolor{red}{12})| = \text{number of weak Bruhat intervals (unknown)}.$

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ありがとうございます！

Thank you!